# Flip processes Jan Hladký // Czech Academy

Frederik Garbe, J.H., Matas Šileikis, Fiona Skerman:

From flip processes to dynamical systems on graphons,

Ann. inst. Henri Poincare (B) Probab, 2023

Pedro Araújo, J. H., Eng Keat Hng, Matas Šileikis: Prominent examples of flip processes

arXiv: 2206.03884

Eng Keat Hng:

Characterization of flip process rules with the same trajectories

arXiv:2305.19925

## Erdos-Renyi random graph process

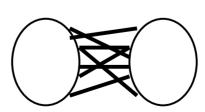
- Erdos-Renyi random graph process (n vertices)  $G_0, G_1, ..., G_{\binom{n}{2}}$ 
  - $-G_0$  is edgeless
  - $G_{r+1}$  is obtained from  $G_r$  by turning a randomly selected nonedge into an edge
- For  $r=\alpha n^2$ , the graph  $G_r$  is a.a.s. quasirandom of density  $2\alpha$ .

## Quasirandomness

- 1980's (Chung-Graham-Wilson, Szemeredi, ...)
- *Density* of a graph  $d=e(G)/\binom{n}{2}$
- A graph is  $\varepsilon$ -quasirandom if for each set of vertices U

$$\left| e(G[U]) - d\binom{|U|}{2} \right| < \varepsilon n^2$$

A nonquasirandom graph



## Triangle removal process

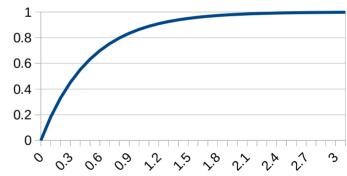
- Introduced by Bollobas-Erdos'90
- Start with G<sub>0</sub>=clique
- In step r, pick a random triangle of G<sub>r</sub> and delete it
- Bohman-Frieze-Lubetzky'15: Triangle removal process typically terminates when there are  $n^{3/2+o(1)}$  edges left.
  - Key in the proof: quasirandomness during the evolution

## Erdos-Renyi *flip* process

- Start with a graph  $G_0$  (for now the edgeless graph)
- In each step, "replace" a uniformly chosen pair with an edge
- Density computation for  $G_r$ ,  $r=\alpha n^2$ :

P[uv is an edge] = 1 - P[uv is not an edge]

...=1-
$$\left(1-\frac{1}{\binom{n}{2}}\right)^r \approx 1-\exp(-2r/n^2)=1-\exp(-2\alpha)$$



## Erdos-Renyi 50:50 flip process

- Start with a graph G<sub>0</sub> on n vertices
- In each step, "replace" a uniformly chosen pair with an edge or an non-edge (50:50)
- "Converges to quasirandom graph of density 0.5", after Cn<sup>2</sup> steps, C→∞

## Triangle removal *flip* process

- Start with a graph G<sub>o</sub> (for now the complete graph)
- In each step r pick three random vertices  $u_1$ ,  $u_2$ ,  $u_3$ ,
- If  $G_r[u_1, u_2, u_3]$  induces a triangle then remove it...

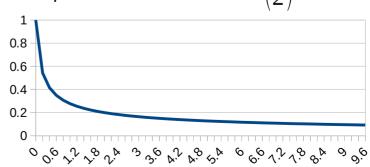
...otherwise  $G_{r+1} := G_r$ .

• Density computation:  $G_r$ ,  $r=\alpha n^2$ ,  $e(\alpha):=e(G_r)$ ,  $d(\alpha):=e(\alpha) / \binom{n}{2}$   $P[u_1u_2u_3 \text{ is a triangle}] \approx d(\alpha)^3$ 

$$e(\alpha + \epsilon) - e(\alpha) \approx -3d(\alpha)^3 \cdot \epsilon n^2$$

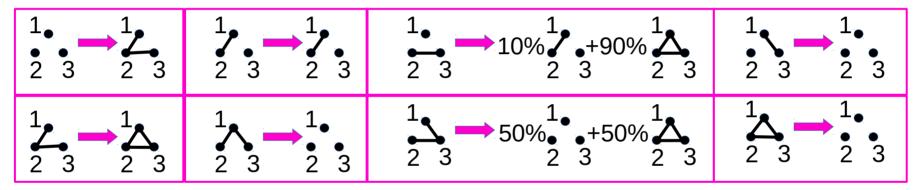
$$\frac{d(\alpha)}{\partial \alpha} = -6d(\alpha)^3 \quad \Longrightarrow \quad d(\alpha) = \frac{1}{\sqrt{1+12\alpha}}$$

Separable first-order ODE



## Flip process of order k (here, k=3)

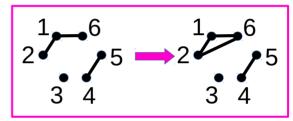
. Rule  $\mathcal{R}$ 



- Start with a (large) graph G<sub>0</sub>
- Step  $G_r \Rightarrow G_{r+1}$ : Sample k vertices and replace the induced graph according to  $\mathcal{R}$

## More examples of flip processes

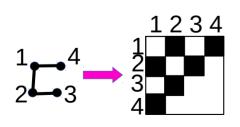
- Ignorant flip process
- F-Removal flip process
- Complement flip process
- Component completion flip process
- The stirring flip process
- The extremist flip process

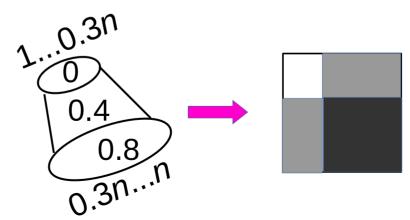


Component completion

## Graphons (limits of dense graphs)

- Borgs-Chayes-Lovasz-Sos-Szegedy-Vesztergombi 2004
- Useful framework for extremal and probabilistic questions
- Graphon is a symmetric function W:[0,1]<sup>2</sup> → [0,1]
- Cut norm measures how similar two graphons are





## Trajectories

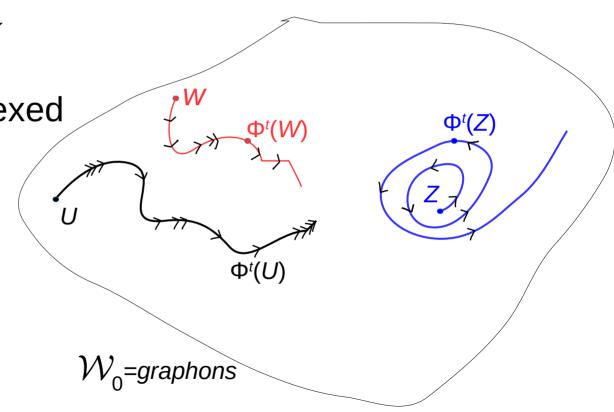
• Fixed rule  $\mathcal{R}$  of order k

We construct time-indexed

trajectories

$$\Phi: \mathcal{W}_0 \times [0, \infty) \to \mathcal{W}_0$$

Construction later



### Transference theorem

**Theorem 5.1.** For every  $k \in \mathbb{N}$  there is a constant C > 0 so that the following holds. Given a rule  $\mathcal{R}$  of order k and a graph G on the vertex set [n], let  $(G_i)_{i\geq 0}$  be the flip process starting with  $G_0 = G$ .

For any T > 0 and  $\varepsilon > 0$  have

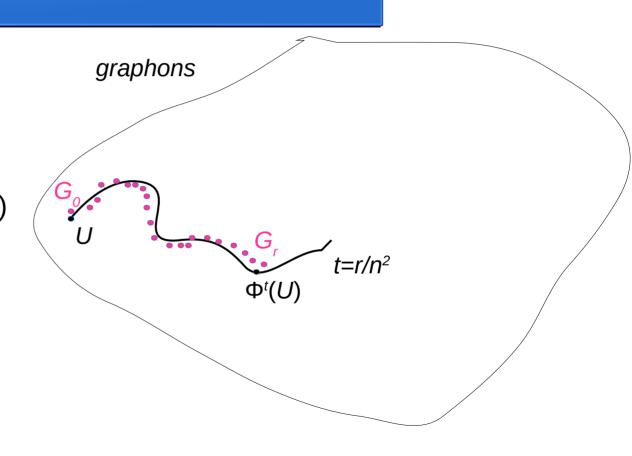
$$\max_{i \in (0, Tn^2] \cap \mathbb{Z}} \left\| W_{G_i} - \Phi^{i/n^2} W_G \right\|_{\square} < \varepsilon$$

 $with\ probability\ at\ least$ 

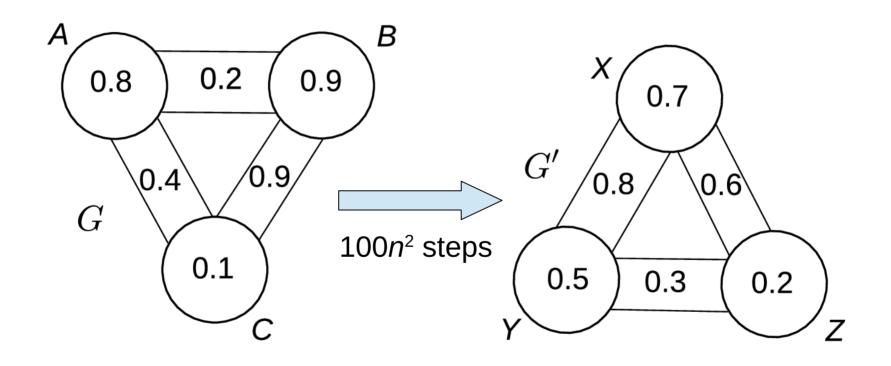
(42) 
$$1 - \frac{CTe^{CT}}{\varepsilon} \exp\left((2\ln 2)n - \frac{C\varepsilon^3 n^2}{e^{CT}}\right).$$

#### Transference theorem

Given  $\mathcal{R}$  and corresponding trajectories  $\Phi: \mathcal{W}_0 \times [0, \infty) \to \mathcal{W}_0$ , whenever a large n-vertex  $G_0$  is close to U (in cut norm) then w.h.p.  $G_r$  is close to  $\Phi^t(U)$  for  $t:=r/n^2$ .

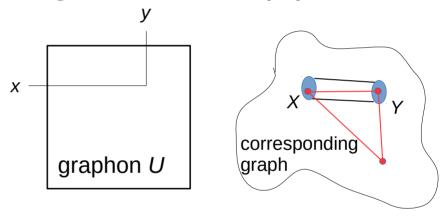


## Cut norm, not cut distance



## Constructing trajectories I

- In this example, consider the Triangle removal flip process
- $(\Phi^{\varepsilon}(U)-U)(x,y)$ correspondence with a graph  $|X|=|Y|=\gamma n$  and  $\varepsilon n^2$  steps  $U(x,y)=e(X,Y)/(\gamma n)^2$



• Number of removed edges between X and Y in  $\varepsilon n^2$  steps:

$$\epsilon n^2 \cdot \gamma^2 \cdot t_{xy}^{"}(K_{3}, U)$$

Density change at (x,y):  $-\epsilon \cdot t_{xy}^{\cdot \cdot}(K_{3},U)$ 

$$t_{xy}^{"}(K_{3},W) = \int_{z} W(x,y)W(x,z)W(y,z)$$

## Constructing trajectories II

• Velocity field  $V: \mathcal{W}_0 \to \mathcal{W}$  (signed graphons)

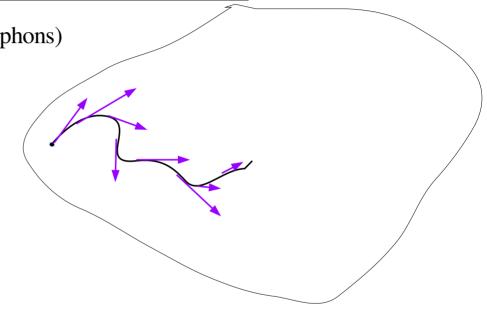
$$V(W) = \lim_{\epsilon \to 0} \frac{\Phi^{\epsilon}(W) - W}{\epsilon}$$

Integral equation

$$\Phi^T W = -W + \oint_{t=0}^T V(\Phi^t W)$$

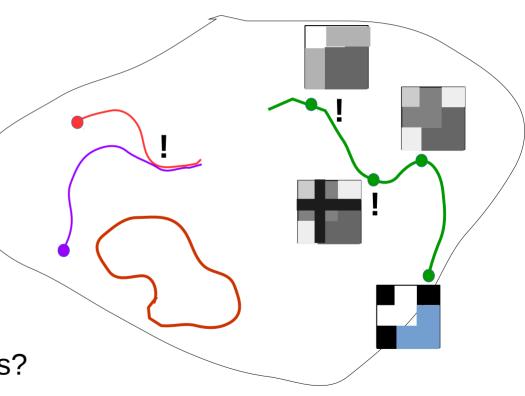


$$V(W)(x,y) = -W(x,y) \int_{z} W(x,z)W(y,z)$$

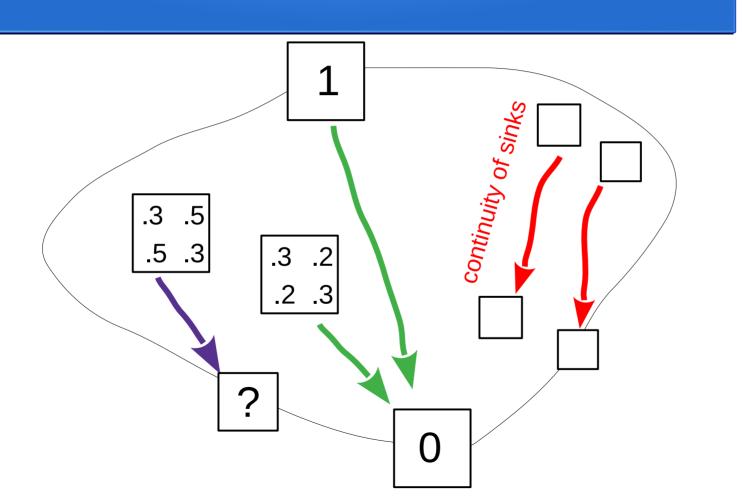


## What is all this good for?

- No confluences
- Going back in time
- Block structure preservation
- Limits  $t \rightarrow \infty$ :
  - Stable and unstable fixed points (often constants)
  - Periodic trajectory
  - Really complicated trajectories?
- Speed of convergence



## On the triangle removal process



## Behavior of individual trajectories

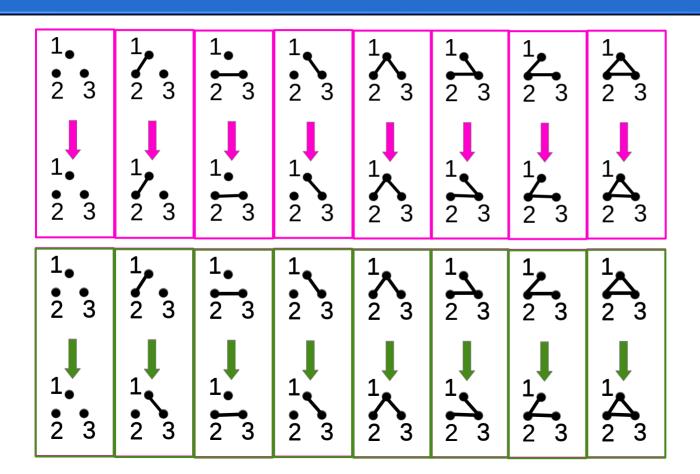
Fix a rule  $\mathcal{R}$ . X is a graphon.

- "Typically", trajectory ( $\Phi^t(X)$ : t) converges to a graphon.
- We have an example of a periodic nonconstant trajectory,  $\Phi^{t}(X)=\Phi^{t+7}(X)$
- Does there exist a really complicated trajectory?

$$S:=\{\Phi^{t}(X): t\in [0,+\infty)\}$$

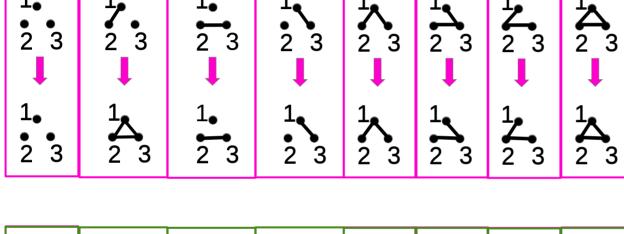
- → Can the set S be totally unbounded?
- $\rightarrow$  Can we have limsup  $\Phi^t(X)=1$  and liminf  $\Phi^t(X)=0$ ?

## Uniqueness (Eng Keat Hng)



Labels matter!

## Uniqueness (Eng Keat Hng)



#### Theorem:

If there is not an obvious reason for two rules of the same order to have the same trajectories, then (some) trajectories will be different.

