

# A new algorithm for computing quadrature-based bounds in conjugate gradients

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joint work with

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## Problem formulation

Consider a system

$$\mathbf{A}x = b$$

where  $\mathbf{A} \in \mathbb{R}^{n \times n}$  is **symmetric, positive definite**.

Without loss of generality,  $\|b\| = 1$ ,  $x_0 = 0$ .

# The conjugate gradient method

**input**  $\mathbf{A}, b$

$$r_0 = b, p_0 = r_0$$

**for**  $k = 1, 2, \dots$  **do**

$$\gamma_{k-1} = \frac{r_{k-1}^T r_{k-1}}{p_{k-1}^T \mathbf{A} p_{k-1}}$$

$$x_k = x_{k-1} + \gamma_{k-1} p_{k-1}$$

$$r_k = r_{k-1} - \gamma_{k-1} \mathbf{A} p_{k-1}$$

$$\delta_k = \frac{r_k^T r_k}{r_{k-1}^T r_{k-1}}$$

$$p_k = r_k + \delta_k p_{k-1}$$

**test quality of**  $x_k$

**end for**

$\mathbf{D}_k$

$$\begin{bmatrix} \gamma_0^{-1} & & & \\ & \ddots & & \\ & & \ddots & \\ & & & \gamma_{k-1}^{-1} \end{bmatrix}$$

$\mathbf{L}_k$

$$\begin{bmatrix} 1 & & & \\ \sqrt{\delta_1} & \ddots & & \\ & \ddots & \ddots & \\ & & \sqrt{\delta_{k-1}} & 1 \end{bmatrix}$$

# How to measure quality of an approximation in CG?

A practically relevant question

- **using residual information,**

- normwise backward error,
  - relative residual norm.

“Using of the residual vector  $r_k$  as a measure of the “goodness” of the estimate  $x_k$  is not reliable” [Hestenes & Stiefel 1952]

- **using error estimates,**

- **the A-norm of the error,**
  - the Euclidean norm of the error.

“The function  $(x - x_k, \mathbf{A}(x - x_k))$  can be used as a measure of the “goodness” of  $x_k$  as an estimate of  $x$ .” [Hestenes & Stiefel 1952]

The **A**-norm of the error plays an important role in stopping criteria [Deuflhard 1994], [Arioli 2004], [Jiránek, Strakoš, Vohralík 2006].

# The Lanczos algorithm

Let  $\mathbf{A}$  be symmetric, compute orthonormal basis of  $\mathcal{K}_k(\mathbf{A}, b)$

**input**  $\mathbf{A}, b$

$$v_1 = b/\|b\|, \delta_1 = 0$$

$$\beta_0 = 0, v_0 = 0$$

**for**  $k = 1, 2, \dots$  **do**

$$\alpha_k = v_k^T \mathbf{A} v_k$$

$$w = \mathbf{A} v_k - \alpha_k v_k - \beta_{k-1} v_{k-1}$$

$$\beta_k = \|w\|$$

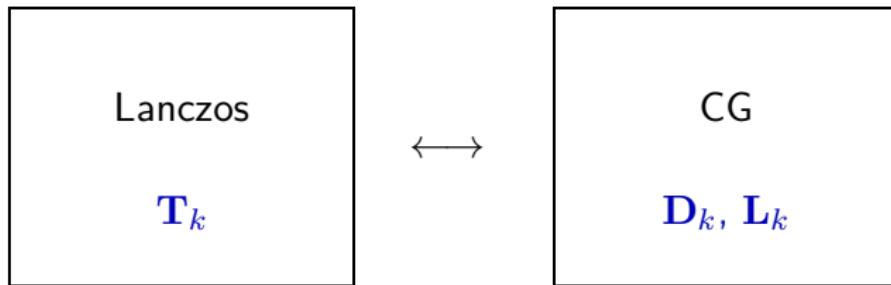
$$v_{k+1} = w/\beta_k$$

**end for**

$$\begin{bmatrix} & & & & \mathbf{T}_k \\ \alpha_1 & \beta_1 & & & \\ \beta_1 & \ddots & & & \\ & \ddots & \ddots & \beta_{k-1} & \\ & & \beta_{k-1} & \alpha_k & \end{bmatrix}$$

# CG versus Lanczos

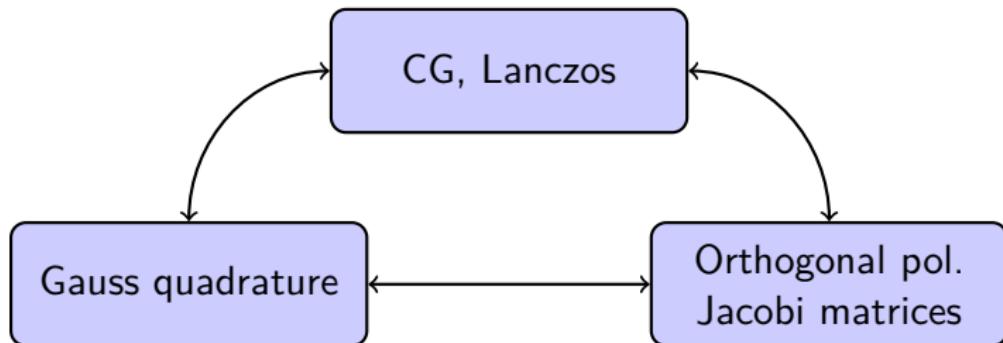
Let  $A$  be symmetric, positive definite



- Both algorithms generate an **orthogonal basis** of  $\mathcal{K}_k(\mathbf{A}, b)$ .
- Lanczos using a **three-term** recurrence  $\rightarrow \mathbf{T}_k$ .
- CG using a **coupled two-term** recurrence  $\rightarrow \mathbf{D}_k, \mathbf{L}_k$ .

$$\mathbf{T}_k = \mathbf{L}_k \mathbf{D}_k \mathbf{L}_k^T.$$

# CG, Lanczos and Gauss quadrature



At any iteration step  $k$ , CG (implicitly) determines **weights** and **nodes** of the  $k$ -point Gauss quadrature

$$\int_{\zeta}^{\xi} f(\lambda) d\omega(\lambda) = \sum_{i=1}^k \omega_i f(\theta_i) + \mathcal{R}_k[f].$$

Gauss quadrature for  $f(\lambda) \equiv \lambda^{-1}$

- Gauss quadrature

$$\int_{\zeta}^{\xi} \lambda^{-1} d\omega(\lambda) = \sum_{i=1}^k \frac{\omega_i}{\theta_i} + \mathcal{R}_k[\lambda^{-1}].$$

- Lanczos

$$\left( \mathbf{T}_n^{-1} \right)_{1,1} = \left( \mathbf{T}_k^{-1} \right)_{1,1} + \mathcal{R}_k[\lambda^{-1}].$$

- CG

$$\|x\|_{\mathbf{A}}^2 = \underbrace{\sum_{j=0}^{k-1} \gamma_j \|r_j\|^2}_{\tau_k} + \|x - x_k\|_{\mathbf{A}}^2.$$

**Important :**  $\mathcal{R}_k[\lambda^{-1}] > 0$ .

# Gauss-Radau quadrature for $f(\lambda) = \lambda^{-1}$

$\mu$  is prescribed

$$\int_{\zeta}^{\xi} f(\lambda) d\omega(\lambda) = \underbrace{\sum_{i=1}^k \tilde{\omega}_i f(\tilde{\theta}_i) + \tilde{\omega}_{k+1} f(\mu)}_{\left( \tilde{\mathbf{T}}_{k+1}^{-1} \right)_{1,1} \equiv \tilde{\tau}_{k+1}} + \mathcal{R}_k[f],$$

where

$$\tilde{\mathbf{T}}_{k+1} = \begin{bmatrix} \alpha_1 & \beta_1 & & & \\ \beta_1 & \ddots & \ddots & & \\ & \ddots & \ddots & \beta_{k-1} & \\ & & \beta_{k-1} & \alpha_k & \beta_k \\ & & & \beta_k & \tilde{\alpha}_{k+1} \end{bmatrix}$$

and  $\mu$  is an eigenvalue of  $\tilde{\mathbf{T}}_{k+1}$ .

**Important:** if  $0 < \mu \leq \lambda_{\min}$ , then  $\mathcal{R}_k[\lambda^{-1}] < 0$ .

# Idea of estimating the $\mathbf{A}$ -norm of the error

[Golub & Strakoš 1994], [Golub & Meurant 1994, 1997]

Consider two quadrature rules at steps  $k$  and  $k+d$ ,  $d > 0$ ,

$$\begin{aligned}\|x\|_{\mathbf{A}}^2 &= \tau_k + \|x - x_k\|_A^2, \\ \|x\|_{\mathbf{A}}^2 &= \hat{\tau}_{k+d} + \hat{\mathcal{R}}_{k+d}.\end{aligned}$$

Then

$$\|x - x_k\|_{\mathbf{A}}^2 = \hat{\tau}_{k+d} - \tau_k + \hat{\mathcal{R}}_{k+d}.$$

**Gauss** quadrature:  $\hat{\mathcal{R}}_{k+d} > 0 \rightarrow \text{lower bound}$ ,

**Radau** quadrature:  $\hat{\mathcal{R}}_{k+d} < 0 \rightarrow \text{upper bound}$ .

How to compute efficiently

$$\hat{\tau}_{k+d} - \tau_k ?$$

## How to compute efficiently $\hat{\tau}_{k+d} - \tau_k$ ?

$$\|x - x_k\|_{\mathbf{A}}^2 = \hat{\tau}_{k+d} - \tau_k + \hat{\mathcal{R}}_{k+d}.$$

For **numerical reasons**, it is not convenient to compute  $\tau_k$  and  $\hat{\tau}_{k+d}$  explicitly. Instead,

$$\begin{aligned}\hat{\tau}_{k+d} - \tau_k &= \sum_{j=k}^{k+d-2} (\tau_{j+1} - \tau_j) + (\hat{\tau}_{j+d} - \tau_{j+d-1}) \\ &\equiv \sum_{j=k}^{k+d-2} \Delta_j + \hat{\Delta}_{k+d-1},\end{aligned}$$

and update the  $\Delta_j$ 's **without subtractions**. Recall  $\tau_j = (\mathbf{T}_j^{-1})_{1,1}$ .

# Golub and Meurant approach

[Golub & Meurant 1994, 1997]: Use **tridiagonal matrices**



and compute  $\Delta$ 's using **updating** strategies,  
**no need to store** tridiagonal matrices.

Use the formulas

$$\|x - x_k\|_{\mathbf{A}}^2 = \sum_{j=k}^{k+d-1} \Delta_j + \|x - x_{k+d}\|_{\mathbf{A}}^2,$$

$$\|x - x_k\|_{\mathbf{A}}^2 = \sum_{j=k}^{k+d-2} \Delta_j + \Delta_{k+d-1}^{(\mu)} + \mathcal{R}_{k+d}^{(R)}.$$

# CGQL (Conjugate Gradients and Quadrature via Lanczos)

**input**  $\mathbf{A}$ ,  $b$ ,  $x_0$ ,  $\mu$

$$r_0 = b - \mathbf{A}x_0, p_0 = r_0$$

$$\delta_0 = 0, \gamma_{-1} = 1, c_1 = 1, \beta_0 = 0, d_0 = 1, \tilde{\alpha}_1^{(\mu)} = \mu,$$

**for**  $k = 1, \dots$ , until convergence **do**

**CG-iteration** ( $k$ )

$$\alpha_k = \frac{1}{\gamma_{k-1}} + \frac{\delta_{k-1}}{\gamma_{k-2}}, \quad \beta_k^2 = \frac{\delta_k}{\gamma_{k-1}^2}$$

$$d_k = \alpha_k - \frac{\beta_{k-1}^2}{d_{k-1}}, \quad \Delta_{k-1} = \|r_0\|^2 \frac{c_k^2}{d_k},$$

$$\tilde{\alpha}_{k+1}^{(\mu)} = \mu + \frac{\beta_k^2}{\alpha_k - \tilde{\alpha}_k^{(\mu)}},$$

$$\Delta_k^{(\mu)} = \|r_0\|^2 \frac{\beta_k^2 c_k^2}{d_k (\tilde{\alpha}_{k+1}^{(\mu)} d_k - \beta_k^2)}, \quad c_{k+1}^2 = \frac{\beta_k^2 c_k^2}{d_k^2}$$

**Estimates**( $k, d$ )

**end for**

# Our approach

[Meurant & T. 2013]: Update  $LDL^T$  decompositions of  $\mathbf{T}_k$  and  $\tilde{\mathbf{T}}_k$



- We use **tridiagonal matrices** only **implicitly**.
- We get **very simple formulas** for updating  $\Delta_{k-1}$  and  $\Delta_k^{(\mu)}$ .
- In [Meurant & T. 2013], this idea is used also for **other types of quadratures** (Gauss-Lobatto, Anti-Gauss).

# CGQ (Conjugate Gradients and Quadrature)

[Meurant & T. 2013]

**input**  $\mathbf{A}$ ,  $b$ ,  $x_0$ ,  $\mu$ ,

$$r_0 = b - \mathbf{A}x_0, p_0 = r_0$$

$$\Delta_0^{(\mu)} = \frac{\|r_0\|^2}{\mu},$$

**for**  $k = 1, \dots$ , until convergence **do**

**CG-iteration( $k$ )**

$$\Delta_{k-1} = \gamma_{k-1} \|r_{k-1}\|^2,$$

$$\Delta_k^{(\mu)} = \frac{\|r_k\|^2 (\Delta_{k-1}^{(\mu)} - \Delta_{k-1})}{\mu (\Delta_{k-1}^{(\mu)} - \Delta_{k-1}) + \|r_k\|^2}$$

**Estimates( $k, d$ )**

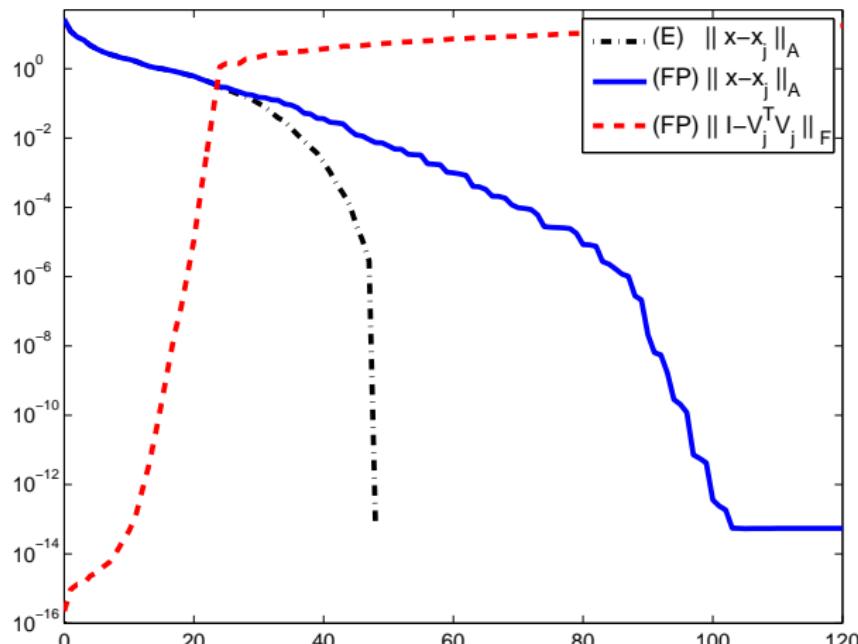
**end for**

## Conclusions (theoretical part)

- **Simple formulas** for computing **bounds** on  $\|x - x_k\|_{\mathbf{A}}$ .
- Almost **for free**.
- Work well also with **preconditioning**.
- Behaviour in **finite precision arithmetic?**

# CG in finite precision arithmetic

Orthogonality is lost, convergence is delayed!



Identities need not hold in finite precision arithmetic!

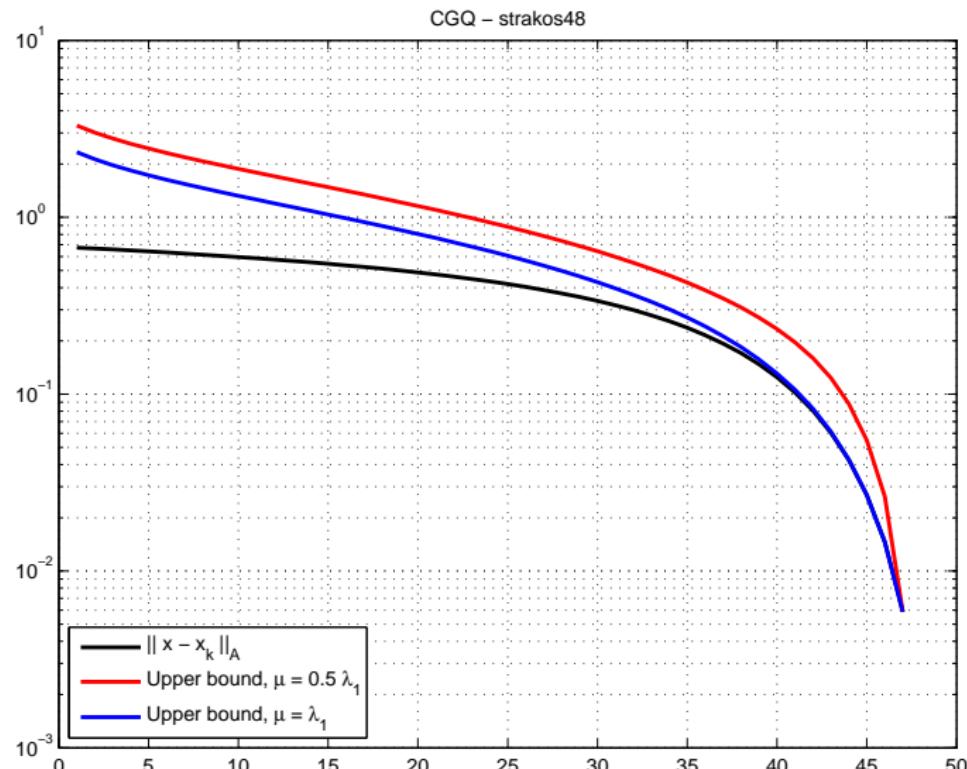
# Bounds in finite precision arithmetic

- **Observation:** CGQL and CGQ give the **same results** (up to a small inaccuracy).
- Do the bounds **correspond to**  $\|x - x_k\|_{\mathbf{A}}$ ?
- **Gauss quadrature** lower bound → **yes** [Strakoš & T. 2002].
- What about the **Gauss-Radau** upper bound?

$$\begin{aligned}\|x - x_k\|_{\mathbf{A}}^2 &= \Delta_k^{(\mu)} + \mathcal{R}_{k+1}^{(R)}, \\ \|x - x_k\|_{\mathbf{A}} &\leq \sqrt{\Delta_k^{(\mu)}}.\end{aligned}$$

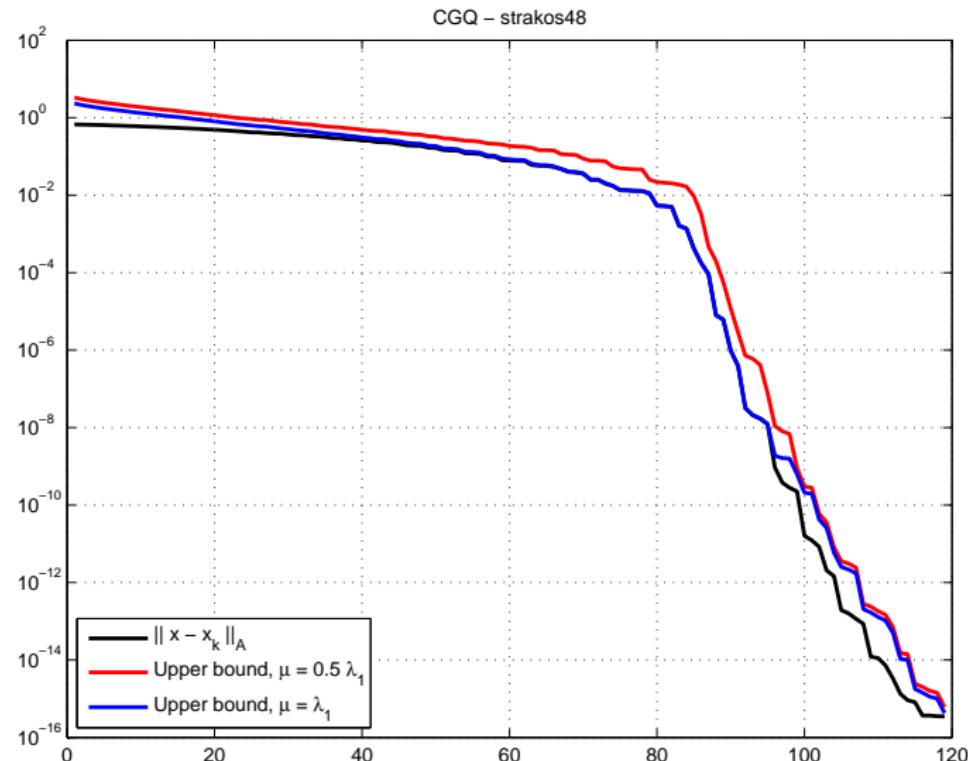
# Gauss-Radau upper bound, exact arithmetic

Strakoš matrix,  $n = 48$ ,  $\lambda_1 = 0.1$ ,  $\lambda_n = 1000$ ,  $\rho = 0.9$ ,  $d = 1$



# Gauss-Radau upper bound, finite precision arithmetic

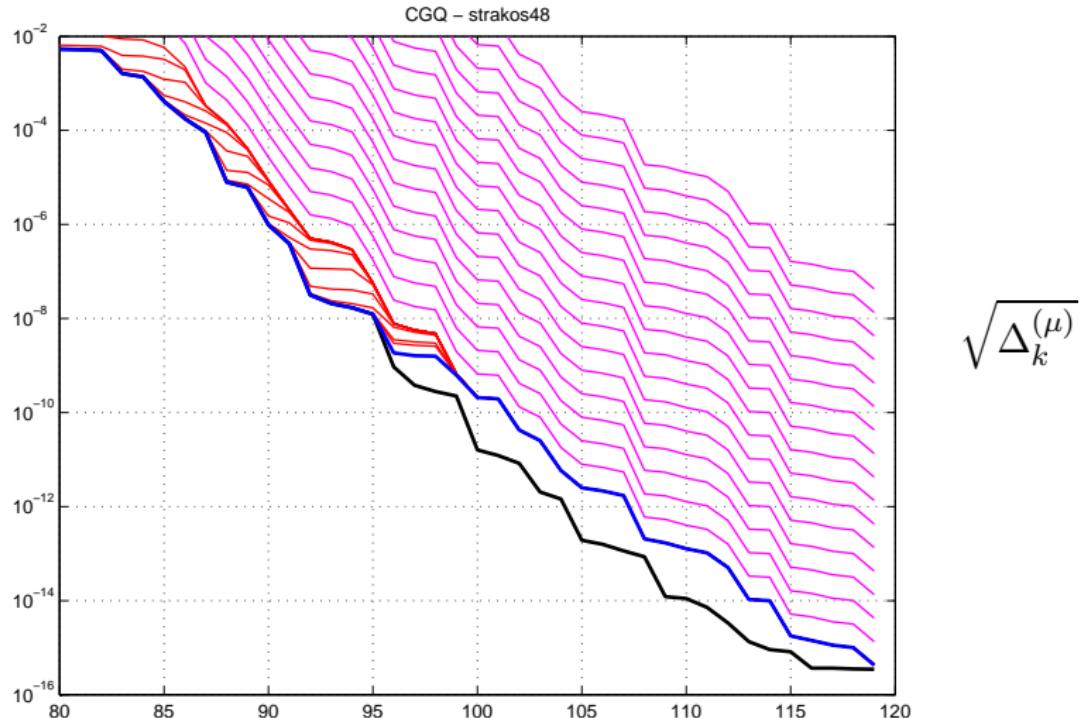
Strakoš matrix,  $n = 48$ ,  $\lambda_1 = 0.1$ ,  $\lambda_n = 1000$ ,  $\rho = 0.9$ ,  $d = 1$



# Gauss-Radau upper bound, finite precision arithmetic

Strakoš matrix,  $n = 48$ ,  $\lambda_1 = 0.1$ ,  $\lambda_n = 1000$ ,  $\rho = 0.9$ ,  $d = 1$

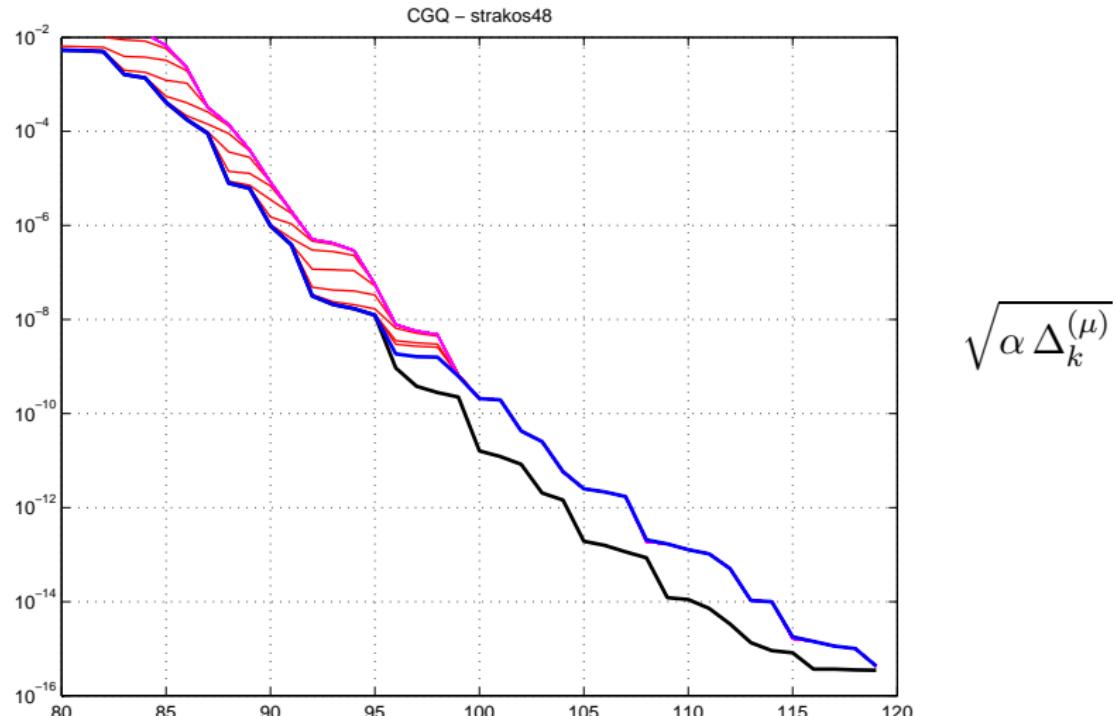
$$\boxed{\mu = \alpha \lambda_1, \alpha \in (0, 1], \alpha \approx 1 \text{ (red)}, \alpha \approx 0 \text{ (magenta)}}$$



# Gauss-Radau upper bound, finite precision arithmetic

Strakoš matrix,  $n = 48$ ,  $\lambda_1 = 0.1$ ,  $\lambda_n = 1000$ ,  $\rho = 0.9$ ,  $d = 1$

$$\boxed{\mu = \alpha \lambda_1, \alpha \in (0, 1], \text{ red} \approx 1, \text{ magenta} \approx 0}$$



# Conclusions (numerical observation)

## Gauss-Radau upper bound

- It seems that  $\sqrt{\varepsilon}$  is a **limiting level** for the accuracy of the **Gauss-Radau** upper bound.
- We **cannot avoid subtractions** in computing this bound.  
If  $\mu \approx \lambda_1$ , then  **$T_k - \mu I$  may be ill conditioned**.
- Simple formulas  $\rightarrow$  **investigation** of numerical behaviour.
- **Understanding** can help
  - in suggesting **another approach**,
  - in **improving Gauss quadrature** lower bound  
(adaptive choice of  $d$ ).

## Related papers

- G. Meurant and P. Tichý, [On computing quadrature-based bounds for the  $\mathbf{A}$ -norm of the error in CG, *Numer. Algorithms*, 62 (2013), pp. 163–191.]
- G. H. Golub and G. Meurant, [Matrices, moments and quadrature with applications, Princeton University Press, USA, 2010.]
- Z. Strakoš and P. Tichý, [On error estimation in CG and why it works in finite precision computations, *ETNA*, 13 (2002), pp. 56–80.]
- G. H. Golub and G. Meurant, [Matrices, moments and quadrature. II. *BIT*, 37 (1997), pp. 687–705.]
- G. H. Golub and Z. Strakoš, [Estimates in quadratic formulas, *Numer. Algorithms*, 8 (1994), pp. 241–268.]

**Thank you for your attention!**