# Numerical behavior of saddle point solvers 

Miro Rozložník<br>joint results with Pavel Jiránek and Valeria Simoncini

Institute of Computer Science, Czech Academy of Sciences,
Prague, Czech Republic

Applied Mathematics and Scientific Computing Seminar, Department of Mathematics, Temple University, Philadelphia, 20 February 2013

## Saddle point problems

We consider a saddle point problem with the symmetric $2 \times 2$ block form

$$
\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x}{y}=\binom{f}{0}
$$

- $A$ is a square $n \times n$ nonsingular (symmetric positive definite) matrix,
- $B$ is a rectangular $n \times m$ matrix of (full column) rank $m$.

Applications: mixed finite element approximations, weighted least squares, constrained optimization, computational fluid dynamics, electromagnetism etc. [Benzi, Golub and Liesen, 2005], [Elman, Silvester, Wathen, 2005]. For the updated list of applications leading to saddle point problems contact [Benzi].

## SOLUTION APPROACH

PRECONDITIONER ITERATIVE SOLVER

## Iterative solution of saddle point problems

1. segregated approach: outer iteration for solving the reduced Schur complement or null-space projected system;
2. coupled approach with block preconditioning: iteration scheme for solving the preconditioned system;
3. rounding errors in floating point arithmetic: numerical stability of the solver

Numerous solution schemes: inexact Uzawa algorithms, inexact null-space methods, inner-outer iteration methods, two-stage iteration processes, multilevel or multigrid methods, domain decomposition methods

Numerous preconditioning techniques and schemes: block diagonal preconditioners, block triangular preconditioners, constraint preconditioning, Hermitian/skew-Hermitian preconditioning and other splittings, combination preconditioning

Numerous iterative solvers: conjugate gradient (CG) method, MINRES, GMRES, flexible GMRES, GCR, BiCG, BiCGSTAB, ...

## Delay of convergence and limit on the final accuracy



Numerical experiments: a small model example

$$
\begin{gathered}
A=\operatorname{tridiag}(1,4,1) \in \mathbb{R}^{100 \times 100}, B=\operatorname{rand}(100,20), f=\operatorname{rand}(100,1), \\
\kappa(A)=\|A\| \cdot\left\|A^{-1}\right\|=5.9990 \cdot 0.4998 \approx 2.9983, \\
\kappa(B)=\|B\| \cdot\left\|B^{\dagger}\right\|=7.1695 \cdot 0.4603 \approx 3.3001 .
\end{gathered}
$$

## Schur complement reduction method

- Compute $y$ as a solution of the Schur complement system

$$
B^{T} A^{-1} B y=B^{T} A^{-1} f
$$

- compute $x$ as a solution of

$$
A x=f-B y
$$

- Segregated vs. coupled approach: $x_{k}$ and $y_{k}$ approximate solutions to $x$ and $y$, respectively.
- Inexact solution of systems with $A$ : every computed solution $\hat{u}$ of $A u=b$ is interpreted as an exact solution of a perturbed system

$$
(A+\Delta A) \hat{u}=b+\Delta b,\|\Delta A\| \leq \tau\|A\|,\|\Delta b\| \leq \tau\|b\|, \tau \kappa(A) \ll 1
$$

Iterative solution of the Schur complement system
choose $y_{0}$, solve $A x_{0}=f-B y_{0}$ compute $\alpha_{k}$ and $p_{k}^{(y)}$

$$
y_{k+1}=y_{k}+\alpha_{k} p_{k}^{(y)}
$$

$$
\text { solve } A p_{k}^{(x)}=-B p_{k}^{(y)}
$$

back-substitution:

A: $x_{k+1}=x_{k}+\alpha_{k} p_{k}^{(x)}$,
inner
B: solve $A x_{k+1}=f-B y_{k+1}, \quad$ iteration
C: solve $A u_{k}=f-A x_{k}-B y_{k+1}$,

$$
x_{k+1}=x_{k}+u_{k} .
$$


outer
iteration

$$
r_{k+1}^{(y)}=r_{k}^{(y)}-\alpha_{k} B^{T} p_{k}^{(x)}
$$

## Accuracy in the saddle point system

$$
\begin{aligned}
\left\|f-A x_{k}-B y_{k}\right\| & \leq \frac{O\left(\alpha_{1}\right) \kappa(A)}{1-\tau \kappa(A)}\left(\|f\|+\|B\| Y_{k}\right), \\
\left\|-B^{T} x_{k}-r_{k}^{(y)}\right\| & \leq \frac{O\left(\alpha_{2}\right) \kappa(A)}{1-\tau \kappa(A)}\left\|A^{-1}\right\|\|B\|\left(\|f\|+\|B\| Y_{k}\right), \\
Y_{k} & \equiv \max \left\{\left\|y_{i}\right\| \mid i=0,1, \ldots, k\right\} .
\end{aligned}
$$

| Back-substitution scheme |  | $\alpha_{1}$ | $\alpha_{2}$ |
| :--- | :--- | :---: | :---: |
| A: | Generic update | $\tau$ | $u$ |
|  | $x_{k+1}=x_{k}+\alpha_{k} p_{k}^{(x)}$ | $\tau$ | $\tau$ |
| B: | Direct substitution |  |  |
|  | $x_{k+1}=A^{-1}\left(f-B y_{k+1}\right)$ |  |  |
| C: | $\begin{array}{l}\text { Corrected dir. subst. }\end{array}$ |  |  |
|  | $x_{k+1}=x_{k}+A^{-1}\left(f-A x_{k}-B y_{k+1}\right)$ |  |  |$)$

$$
-B^{T} A^{-1} f+B^{T} A^{-1} B y_{k}=-B^{T} x_{k}-B^{T} A^{-1}\left(f-A x_{k}-B y_{k}\right)
$$

Generic update: $x_{k+1}=x_{k}+\alpha_{k} p_{k}^{(x)}$


Direct substitution: $x_{k+1}=A^{-1}\left(f-B y_{k+1}\right)$


Corrected direct substitution: $x_{k+1}=x_{k}+A^{-1}\left(f-A x_{k}-B y_{k+1}\right)$


## Null-space projection method

- compute $x \in N\left(B^{T}\right)$ as a solution of the projected system

$$
(I-\Pi) A(I-\Pi) x=(I-\Pi) f
$$

- compute $y$ as a solution of the least squares problem

$$
B y \approx f-A x
$$

$\Pi=B\left(B^{T} B\right)^{-1} B^{T}$ is the orthogonal projector onto $R(B)$.

- Schemes with the inexact solution of least squares with $B$. Every computed approximate solution $\bar{v}$ of a least squares problem $B v \approx c$ is interpreted as an exact solution of a perturbed least squares

$$
(B+\Delta B) \bar{v} \approx c+\Delta c,\|\Delta B\| \leq \tau\|B\|,\|\Delta c\| \leq \tau\|c\|, \tau \kappa(B) \ll 1
$$

## Null-space projection method

choose $x_{0}$, solve $B y_{0} \approx f-A x_{0}$ compute $\alpha_{k}$ and $p_{k}^{(x)} \in N\left(B^{T}\right)$
$x_{k+1}=x_{k}+\alpha_{k} p_{k}^{(x)}$
solve $B p_{k}^{(y)} \approx r_{k}^{(x)}-\alpha_{k} A p_{k}^{(x)}$ back-substitution:
A: $y_{k+1}=y_{k}+p_{k}^{(y)}$,
B: solve $B y_{k+1} \approx f-A x_{k+1}$,
$\left\{\begin{array}{l}\text { outer } \\ \text { iteration }\end{array}\right.$
$\left\{\begin{array}{l}\text { outer } \\ \text { iteration }\end{array}\right.$
C: solve $B v_{k} \approx f-A x_{k+1}-B y_{k}$,

$$
y_{k+1}=y_{k}+v_{k}
$$

inner
iteration
J
$r_{k+1}^{(x)}=r_{k}^{(x)}-\alpha_{k} A p_{k}^{(x)}-B p_{k}^{(y)}$

Accuracy in the saddle point system

$$
\begin{gathered}
\left\|f-A x_{k}-B y_{k}-r_{k}^{(x)}\right\| \leq \frac{O\left(\alpha_{3}\right) \kappa(B)}{1-\tau \kappa(B)}\left(\|f\|+\|A\| X_{k}\right) \\
\left\|-B^{T} x_{k}\right\| \leq \frac{O(\tau) \kappa(B)}{1-\tau \kappa(B)}\|B\| X_{k} \\
X_{k} \equiv \max \left\{\left\|x_{i}\right\| \mid i=0,1, \ldots, k\right\}
\end{gathered}
$$

| Back-substitution scheme |  | $\alpha_{3}$ |
| :--- | :--- | :---: |
| A: | Generic update <br> $y_{k+1}=y_{k}+p_{k}^{(y)}$ | $u$ |
| B: | Direct substitution <br> $y_{k+1}=B^{\dagger}\left(f-A x_{k+1}\right)$ | $\tau$ |
| C: | Corrected dir. subst. <br> $y_{k+1}=y_{k}+B^{\dagger}\left(f-A x_{k+1}-B y_{k}\right)$ | $u$ |



Generic update: $y_{k+1}=y_{k}+p_{k}^{(y)}$


Direct substitution: $y_{k+1}=B^{\dagger}\left(f-A x_{k+1}\right)$



Corrected direct substitution: $y_{k+1}=y_{k}+B^{\dagger}\left(f-A x_{k+1}-B y_{k}\right)$


## Stationary iterative methods

- $\mathcal{A}=\mathcal{A} x=b, \mathcal{M}-\mathcal{N}$
- A: $\mathcal{M} x_{k+1}=\mathcal{N} x_{k}+b$

B: $x_{k+1}=x_{k}+\mathcal{M}^{-1}\left(b-\mathcal{A} x_{k}\right)$

- Inexact solution of systems with $\mathcal{M}$ : every computed solution $\bar{y}$ of $\mathcal{M} y=z$ is interpreted as an exact solution of a perturbed system

$$
(\mathcal{M}+\Delta \mathcal{M}) \bar{y}=z, \quad\|\Delta \mathcal{M}\| \leq \tau\|\mathcal{M}\|, \quad \tau k(\mathcal{M}) \ll 1
$$

Accuracy of the computed approximate solution

$$
\begin{aligned}
& \text { A } \mathcal{M} x_{k+1}=\mathcal{N} x_{k}+b \\
& \qquad \frac{\left\|\hat{x}_{k+1}-x\right\|}{\|x\|} \leq \tau \frac{\left\|\left|\mathcal{M}^{-1}\right||\mathcal{M}||x|\right\|}{\|x\|} \\
& \text { B } \quad x_{k+1}=x_{k}+\mathcal{M}^{-1}\left(b-\mathcal{A} x_{k}\right) \\
& \\
& \quad \frac{\left\|\hat{x}_{k+1}-x\right\|}{\|x\|} \leq O(u) \frac{\left\|\left|\mathcal{M}^{-1}\right|(|\mathcal{M}|+|\mathcal{N}|)|x|\right\|}{\|x\|}
\end{aligned}
$$

new_value=old_value+small_correction

Two-stage iterative methods

$$
\begin{gathered}
\mathcal{M}_{1} x_{k+1 / 2}=\mathcal{N}_{1} x_{k}+b, \quad \mathcal{A}=\mathcal{M}_{1}-\mathcal{N}_{1} \\
\mathcal{M}_{2} x_{k+1}=\mathcal{N}_{2} x_{k+1 / 2}+b, \quad \mathcal{A}=\mathcal{M}_{2}-\mathcal{N}_{2} \\
x_{k+1 / 2}=x_{k}+\mathcal{M}_{1}^{-1}\left(b-\mathcal{A} x_{k}\right) \\
x_{k+1}=x_{k+1 / 2}+\mathcal{M}_{2}^{-1}\left(b-\mathcal{A} x_{k+1 / 2}\right) \\
\Leftrightarrow \\
x_{k+1}=x_{k}+\left(\mathcal{M}_{1}^{-1}+\mathcal{M}_{2}^{-1}-\mathcal{M}_{2}^{-1} \mathcal{A} \mathcal{M}_{1}^{-1}\right)\left(b-\mathcal{A} x_{k}\right) \\
=x_{k}+\left(\mathcal{I}+\mathcal{M}_{2}^{-1} \mathcal{N}_{1}\right) \mathcal{M}_{1}^{-1}\left(b-\mathcal{A} x_{k}\right) \\
=x_{k}+\mathcal{M}_{2}^{-1}\left(\mathcal{I}+\mathcal{N}_{2} \mathcal{M}_{1}^{-1}\right)\left(b-\mathcal{A} x_{k}\right) \\
\frac{\left\|x_{k+1}-x\right\|}{\|x\|} \leq O(u) \frac{\left\|\left|\mathcal{M}^{-1}\right|(|\mathcal{M}|+|\mathcal{N}|)|x|\right\|}{\|x\|}
\end{gathered}
$$

## Preconditioning of saddle point problems

$\mathcal{A}$ symmetric indefinite, $\mathcal{P}$ positive definite

$$
\begin{gathered}
\mathcal{A}=\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right) \approx \mathcal{P}=\mathcal{R}^{T} \mathcal{R} \\
\left(\mathcal{R}^{-T} \mathcal{A} \mathcal{R}^{-1}\right) \mathcal{R}\binom{x}{y}=\mathcal{R}^{-T}\binom{f}{0}
\end{gathered}
$$

$$
\mathcal{R}^{-T} \mathcal{A} \mathcal{R}^{-1} \text { is symmetric indefinite! }
$$

## Symmetric indefinite or nonsymmetric preconditioner

$\mathcal{P}$ symmetric indefinite or nonsymmetric

$$
\begin{aligned}
& \mathcal{P}^{-1} \mathcal{A}\binom{x}{y}=\mathcal{P}^{-1}\binom{f}{0} \\
& \left(\mathcal{A P}^{-1}\right) \mathcal{P}\binom{x}{y}=\binom{f}{0}
\end{aligned}
$$

$\mathcal{P}^{-1} \mathcal{A}$ and $\mathcal{A} \mathcal{P}^{-1}$ are nonsymmetric!

## Schur complement approach with indefinite preconditioner

$$
\begin{aligned}
\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x}{y} & =\binom{f}{0}, \quad \mathcal{P}=\left(\begin{array}{cc}
A & B \\
B^{T} & B^{T} A^{-1} B-I
\end{array}\right) \\
\mathcal{A P}^{-1} & =\left(\begin{array}{cc}
I & 0 \\
(I-S) B^{T} A^{-1} & S
\end{array}\right)
\end{aligned}
$$

$S=B^{T} A^{-1} B, \mathcal{A} \mathcal{P}^{-1}$ nonsymmetric but diagonalizable and it has a 'nice' spectrum!

$$
\sigma\left(\mathcal{A} \mathcal{P}^{-1}\right) \subset\{1\} \cup \sigma\left(B^{T} A^{-1} B^{T}\right)
$$

[Durazzi, Ruggiero 2003], [Fortin, El-Maliki, 2009?]

Krylov method with the preconditioner: basic properties

$$
\begin{gathered}
\binom{x_{0}}{y_{0}}, r_{0}=\binom{0}{s_{0}}, e_{k+1}=\binom{x-x_{k+1}}{y-y_{k+1}} \\
r_{k+1}=\binom{f}{0}-\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x_{k+1}}{y_{k+1}} \\
r_{0}=\binom{0}{s_{0}} \Rightarrow r_{k+1}=\binom{0}{s_{k+1}} \\
\Rightarrow A x_{k+1}+B y_{k+1}=f
\end{gathered}
$$

Preconditioned CG method: saddle point problem and indefinite preconditioner

$$
r_{k+1}^{T} \mathcal{P}^{-1} r_{j}=0, j=0, \ldots, k
$$

$y_{k+1}$ is an iterate from CG applied to the Schur complement system

$$
B^{T} A^{-1} B y=B^{T} A^{-1} f!
$$

satisfying

$$
\begin{gathered}
\left\|y-y_{k+1}\right\|_{B^{T} A^{-1} B}= \\
\min _{u \in x_{0}+K_{k+1}\left(B^{T} A^{-1} B, B^{T} A^{-1} f\right)}\|y-u\|_{B^{T} A^{-1} B}
\end{gathered}
$$

## Preconditioned CG algorithm

$$
\begin{aligned}
& \binom{x_{0}}{y_{0}}, r_{0}=b-\mathcal{A}\binom{x_{0}}{y_{0}}=\binom{0}{s_{0}} \\
& \binom{p_{0}^{(x)}}{p_{0}^{(y)}}=\mathcal{P}^{-1} r_{0}=\mathcal{P}^{-1}\binom{0}{s_{0}} \\
& k=0,1, \ldots \\
& \alpha_{k}=\left(\binom{0}{s_{k}}, \mathcal{P}^{-1}\binom{0}{s_{k}}\right) /\left(\mathcal{A}\binom{p_{k}^{(x)}}{p_{k}^{(y)}},\binom{p_{k}^{(x)}}{p_{k}^{(y)}}\right) \\
& \binom{x_{k+1}}{y_{k+1}}=\binom{x_{k}}{y_{k}}+\alpha_{k}\binom{p_{k}^{(x)}}{p_{k}^{(y)}} \\
& r_{k+1}=r_{k}-\alpha_{k} \mathcal{A}\binom{p_{k}^{(x)}}{p_{k}^{(y)}}=\binom{0}{s_{k+1}} \\
& \beta_{k}=\left(\binom{0}{s_{k+1}}, \mathcal{P}^{-1}\binom{0}{s_{k+1}}\right) /\left(\binom{0}{\left.s_{k}, z_{k}, p_{k}\right)}, \mathcal{P}^{-1}\binom{0}{s_{k}}\right) \\
& \binom{p_{k+1}^{(x)}}{p_{k+1}^{(y)}}=\mathcal{P}^{-1}\binom{0}{s_{k+1}}+\beta_{k}\binom{p_{k}^{(x)}}{p_{k}^{(y)}}=\binom{-A^{-1} B p_{k+1}^{(y)}}{p_{k+1}^{(y)}}
\end{aligned}
$$

Generic update: $x_{k+1}=x_{k}+\alpha_{k} p_{k}^{(x)}$ with $p_{k}^{(x)}=-A^{-1} B p_{k}^{(y)}$



## Saddle point problem and indefinite constraint preconditioner

$$
\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x}{y}=\binom{f}{0}, \quad \mathcal{P}=\left(\begin{array}{cc}
I & B \\
B^{T} & 0
\end{array}\right)
$$

$$
\mathcal{A} \mathcal{P}^{-1}=\left(\begin{array}{cc}
A(I-\Pi)+\Pi & (A-I) B\left(B^{T} B\right)^{-1} \\
0 & I
\end{array}\right)
$$

$$
\Pi=B\left(B^{T} B\right)^{-1} B^{T}-\text { orth. projector onto } \operatorname{span}(B)
$$

[Lukšan, Vlček, 1998], [Gould, Keller, Wathen 2000]
[Perugia, Simoncini, Arioli, 1999], [R, Simoncini, 2002]

Indefinite constraint preconditioner: spectral properties

## $\mathcal{A} \mathcal{P}^{-1}$ nonsymmetric and non-diagonalizable! but it has a 'nice' spectrum:

$$
\begin{aligned}
\sigma\left(\mathcal{A P}^{-1}\right) & \subset\{1\} \cup \sigma(A(I-\Pi)+\Pi) \\
& \subset\{1\} \cup \sigma((I-\Pi) A(I-\Pi))-\{0\}
\end{aligned}
$$

and only 2 by 2 Jordan blocks!
[Lukšan, Vlček 1998], [Gould, Wathen, Keller, 1999], [Perugia, Simoncini 1999]

Krylov method with the constraint preconditioner: basic properties

$$
\begin{gathered}
\binom{x_{0}}{y_{0}}, r_{0}=\binom{s_{0}}{0}, e_{k+1}=\binom{x-x_{k+1}}{y-y_{k+1}} \\
r_{k+1}=\binom{f}{0}-\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x_{k+1}}{y_{k+1}} \\
r_{0}=\binom{s_{0}}{0} \Rightarrow r_{k+1}=\binom{s_{k+1}}{0} \\
\Rightarrow B^{T}\left(x-x_{k+1}\right)=0 \\
\end{gathered} \begin{aligned}
& \Rightarrow x_{k+1} \in N u l l\left(B^{T}\right)!
\end{aligned}
$$

## Preconditioned CG method: error norm

$$
\begin{gathered}
r_{k+1}^{T} \mathcal{P}^{-1} r_{j}=0, j=0, \ldots, k \\
x_{k+1} \text { is an iterate from CG applied to } \\
(I-\Pi) A(I-\Pi) x=(I-\Pi) f! \\
\text { satisfying } \\
\left\|x-x_{k+1}\right\|_{A}=\min _{u \in x_{0}+\operatorname{span}\left\{(I-\Pi) s_{j}\right\}}\|x-u\|_{A}
\end{gathered}
$$

[Lukšan, Vlček 1998], [Gould, Wathen, Keller, 1999]

## Preconditioned CG algorithm

$$
\begin{aligned}
& \binom{x_{0}}{y_{0}}, r_{0}=b-\mathcal{A}\binom{x_{0}}{y_{0}}=\binom{s_{0}}{0} \\
& \binom{p_{0}^{(x)}}{p_{0}^{(y)}}=\mathcal{P}^{-1} r_{0}=\mathcal{P}^{-1}\binom{s_{0}}{0} \\
& k=0,1, \ldots \\
& \alpha_{k}=\left(\binom{s_{k}}{0}, \mathcal{P}^{-1}\binom{s_{k}}{0}\right) /\left(\mathcal{A}\binom{p_{k}^{(x)}}{p_{k}^{(y)}},\binom{p_{k}^{(x)}}{p_{k}^{(y)}}\right) \\
& \alpha_{k}=\left(r_{k}, z_{k}\right) /\left(\mathcal{A} p_{k}, p_{k}\right) \\
& \binom{x_{k+1}}{y_{k+1}}=\binom{x_{k}}{y_{k}}+\alpha_{k}\binom{p_{k}^{(x)}}{p_{k}^{(y)}} \\
& r_{k+1}=r_{k}-\alpha_{k} \mathcal{A}\binom{p_{k}^{(x)}}{p_{k}^{(y)}}=\binom{s_{k+1}}{0} \\
& \beta_{k}=\left(\binom{s_{k+1}}{0}, \mathcal{P}^{-1}\binom{s_{k+1}}{0}\right) /\left(\binom{s_{k}}{0}, \mathcal{P}^{-1}\binom{s_{k}}{0}\right) \\
& \beta_{k}=\left(r_{k+1}, z_{k+1}\right) /\left(r_{k}, z_{k}\right) \\
& \binom{p_{k+1}^{(x)}}{p_{k+1}^{(y)}}=\mathcal{P}^{-1}\binom{s_{k+1}}{0}+\beta_{k}\binom{p_{k}^{(x)}}{p_{k}^{(y)}}
\end{aligned}
$$

## Preconditioned CG method: residual norm

$$
\left\|x_{k+1}-x\right\| \rightarrow 0
$$

but in general

$$
y_{k+1} \nrightarrow y
$$

which is reflected in

$$
\left\|r_{k+1}\right\|=\left\|\binom{s_{k+1}}{0}\right\| \nrightarrow 0!
$$

but under appropriate scaling yes!

## Preconditioned CG method: residual norm

$$
\begin{gathered}
x_{k+1} \rightarrow x \\
x-x_{k+1}=\phi_{k+1}((I-\Pi) A(I-\Pi))\left(x-x_{0}\right) \\
s_{k+1}=\phi_{k+1}(A(I-\Pi)+\Pi) s_{0} \\
\sigma((I-\Pi) A(I-\Pi)) \sim \sigma(A(I-\Pi)+\Pi) ?
\end{gathered}
$$

$$
\{1\} \in \sigma((I-\Pi) \alpha A(I-\Pi))-\{0\}
$$

$$
\Rightarrow\left\|r_{k+1}\right\|=\left\|\binom{s_{k+1}}{0}\right\| \rightarrow 0!
$$

How to avoid misconvergence?

- Scaling by a constant $\alpha>0$ such that

$$
\begin{gathered}
\{1\} \in \operatorname{conv}(\sigma((I-\Pi) \alpha A(I-\Pi))-\{0\}) \\
\left(\begin{array}{cc}
A & B \\
B^{T} & 0
\end{array}\right)\binom{x}{y}=\binom{f}{0} \Longleftrightarrow\left(\begin{array}{cc}
\alpha A & B \\
B^{T} & 0
\end{array}\right)\binom{x}{\alpha y}=\binom{\alpha f}{0} \\
v: \quad\|(I-\Pi) v\| \neq 0, \quad \alpha=\frac{1}{((I-\Pi) v, A(I-\Pi) v)}!
\end{gathered}
$$

- Scaling by a diagonal $A \rightarrow(\operatorname{diag}(A))^{-1 / 2} A(\operatorname{diag}(A))^{-1 / 2}$ often gives what we want!
- Different direction vector $p_{k}^{(y)}$ so that $\left\|r_{k+1}\right\|=\left\|s_{k+1}\right\|$ is locally minimized!

$$
y_{k+1}=y_{k}+\left(B^{T} B\right)^{-1} B^{T} s_{k}
$$

[Braess, Deuflhard,Lipikov 1999], [Hribar, Gould, Nocedal, 1999], [Jiránek, R, 2008]

Numerical experiments: a small model example

$$
\begin{aligned}
& A=\operatorname{tridiag}(1,4,1) \in \mathrm{R}^{25,25}, B=\operatorname{rand}(25,5) \in \mathrm{R}^{25,5} \\
& f=\operatorname{rand}(25,1) \in \mathrm{R}^{25} \\
& \sigma(A) \subset[2.0146,5.9854] \\
& \alpha=1 / \tau \quad \sigma\left(\left(\begin{array}{cc}
\alpha A & B \\
B^{T} & 0
\end{array}\right)\left(\begin{array}{cc}
I & B \\
B^{T} & 0
\end{array}\right)^{-1}\right) \\
& 1 / 100 \quad[0.0207,0.0586] \cup\{1\} \\
& 1 / 10 \quad[0.2067,0.5856] \cup\{1\} \\
& 1 / 4 \quad[0.5170,1.4641] \\
& 1 \quad\{1\} \cup[2.0678,5.8563] \\
& 4 \quad\{1\} \cup[8.2712,23.4252]
\end{aligned}
$$




PROBLEM

SOLUTION APPROACH

PRECONDITIONER
MAXIMUM ATTAINABLE ACCURACY

ITERATIVE SOLVER

## Conclusions: segregated solution approach

- The accuracy measured by the residuals of the saddle point problem depends on the choice of the back-substitution scheme [Jiránek, R, 2008]. The schemes with (generic or corrected substitution) updates deliver approximate solutions which satisfy either the first or second block equation to working accuracy.
- Care must be taken when solving nonsymmetric systems [Jiránek, R, 2008], all bounds of the limiting accuracy depend on the maximum norm of computed iterates, cf. [Greenbaum 1994,1997], [Sleijpen, et al. 1994].



## Conclusions: coupled approach with indefinite preconditioner

- Short-term recurrence methods are applicable for saddle point problems with indefinite preconditioning at a cost comparable to that of symmetric solvers. There is a tight connection between the simplified $\mathrm{Bi}-\mathrm{CG}$ algorithm and the classical CG.
- The convergence of CG applied to saddle point problem with indefinite preconditioner for all right-hand side vectors is not guaranteed. For a particular set of right-hand sides the convergence can be achieved by the appropriate scaling of the saddle point problem.
- Since the maximum attainable accuracy depends heavily on the size of computed residuals, a good scaling of the problems leads to approximate solutions satisfying both two block equations to the working accuracy.


## Thank you for your attention.

```
http://www.cs.cas.cz/~miro
```

M. Rozložník and V. Simoncini, Krylov subspace methods for saddle point problems with indefinite preconditioning, SIAM J. Matrix Anal. Appl., 24 (2002), pp. 368-391.
P. Jiránek and M. Rozložník. Limiting accuracy of segregated solution methods for nonsymmetric saddle point problems. J. Comput. Appl. Math. 215 (2008), pp. 28-37.
P. Jiránek and M. Rozložník. Maximum attainable accuracy of inexact saddle point solvers. SIAM J. Matrix Anal. Appl., 29(4):1297-1321, 2008.

