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## Correction of Coefficients of the Input-Output Model

An economic system described by the input-output model

$$(E - A_0) x = y \tag{0}$$

(see e.g. [1]) is fully determined by the matrix  $A_0 = (a_{ij}^0)_{n \times n}$ , but in practice it is very difficult to find out the exact values of its coefficients. Here we suppose that these values are unknown but  $a_{ij}^0 \in \langle a_{ij}, a_{ij} \rangle$ , where  $a_{ij}, a_{ij}$  are given real numbers,  $a_{ij} \leq a_{ij}$ , for i, j = 1, ..., n. Suppose it is known that a final demand y corresponds to a gross output x in the sense of model (0). The main problem of this paper is how to use this information for obtaining more precise estimations of  $a^o_{ij}$ , i.e. for the sharpening of intervals  $\langle \underline{a}_{ij}, \overline{a}_{ij} \rangle$ , i, j = 1, ..., n. The presented results are contained

Denote  $\underline{A} = (\underline{a}_{ij})_{n \times n}, \overline{A} = (a_{ij})_{n \times n}, \langle A, A \rangle = \{A \mid A \leq A \leq \overline{A}\}$  and let  $M = \{A \mid A \in \langle A, \overline{A} \rangle, (E - A)x = y\}$ . Now the problem under consideration can be formulated as the problem of computing numbers

$$\tilde{a}_{ii} = \max\{(A)_{ij} \mid A \in M\}, \quad \tilde{a}_{ij} = \min\{(A)_{ij} \mid A \in M\}, \quad i, j = 1, ..., n$$

Let  $\underline{A} = (\underline{a}_{ij})_{n \times n}$ ,  $\widetilde{A} = (\widetilde{a}_{ij})_{n \times n}$ . It is obvious that  $A_0 \in \langle \underline{A}, \widetilde{A} \rangle \in \langle A, \widetilde{A} \rangle$ . In the sequel we assume that the following two conditions are satisfied:

(i) 
$$x > 0$$
,

(ii) 
$$(E - \bar{A}) x \leq y \leq (E - \underline{A}) x$$
.

Condition (i) requires all the sectors of the system to be producing and (ii) holds if and only if  $M \neq \emptyset$ . The "if" part is clear, because if  $A \in M$ , then  $(E - A) x \leq (E - A) x = y \leq (E - A) x$ , the "only if" part is proved below.

Theorem 1: Let (i) and (ii) be satisfied. Then

$$\tilde{a}_{ij} = \min \left( \tilde{a}_{ij}; \frac{1}{x_j} \left( x_i - y_i - \sum_{k \neq j} a_{ik} x_k \right) \right) \tag{1}$$

$$\tilde{a}_{ij} = \min \left( \tilde{a}_{ij} : \frac{1}{x_j} (x_i - y_i - \sum_{k \neq j} a_{ik} x_k) \right)$$

$$\underline{a}_{ij} = \max \left( a_{ij} : \frac{1}{x_j} (x_i - y_i - \sum_{k \neq j} \tilde{a}_{ik} x_k) \right),$$

$$(1)$$

$$i, j = 1, \dots, n.$$

$$(2)$$

Proof: We shall prove the statement (1) only, because the proof of (2) is quite analogous. For  $l=1,\ldots,n$ 

put  $t_l = \frac{((E-A)x-y)_l}{(A-A)x_l}$  if  $((A-A)x_l)_l \neq 0$  and  $t_l = 0$  otherwise. It follows from (ii) that  $t_l \in (0,1)$  for each l.

Let  $a_{lk}^* - \underline{a_{lk}} + l_l(\overline{a_{lk}} - \underline{a_{lk}}), \ l, k = 1, \dots, n$ , and let  $A^* = (a_{lk}^*)_{n \times n}$ . Then  $A^* \in \langle \underline{A}, \overline{A} \rangle$  and  $(E - A^*)_{x} = y$ , so that  $A^* \in M$ . Hence  $M \neq \emptyset$ .

Let 
$$i, j \in \{1, \dots, n\}$$
. Put  $\alpha_{ij} = \frac{1}{x_j}(x_i - y_i - \sum_{k \neq j} \underline{a}_{ik}x_k)$ . Let  $A = (a_{ik})_{n \leq n} \in M$ ; then  $(E - A) x = y$ , thus

$$a_{ij} = \frac{1}{x_i} (x_i - y_t - \sum\limits_{k \neq j} a_{ik} x_k) \le \alpha_{ij}$$
. Since  $a_{ij} \le \bar{a}_{ij}$ , it implies that  $a_{ij} \le \min{(\bar{a}_{ij}; \alpha_{ij})}$ , hence also

$$\tilde{a}_{ij} = \max \{(A)_{ij} \mid A \in M\} \leq \min (\tilde{a}_{ij}; \alpha_{ij}).$$

To complete the proof, it will suffice to find a matrix  $A \in M$  such that  $(A)_{ij} = \min (\bar{a}_{ij}; \alpha_{ij})$ . We distinguish two

- 1) Let  $\alpha_{ij} \leq \bar{a}_{ij}$ . It follows from (ii) that  $\underline{a}_{ij} \leq \alpha_{ij}$ , so that  $\alpha_{ij} \in \langle \underline{a}_{ij}, \bar{a}_{ij} \rangle$ . Put  $a_{ij}^1 = \alpha_{ij}$ ,  $a_{ik}^1 = \underline{a}_{ik}$  for  $k \neq j$ ,  $a_{ik}^1 = a_{ik}^*$  for  $l \neq i$ ,  $k = 1, \ldots, n$ , and let  $A^1 = (a_{ik}^1)_{n \times n}$ . Then  $A^1 \in M$  and  $(A^1)_{ij} = \alpha_{ij} = \min{(\bar{a}_{ij}; \alpha_{ij})}$ .
  - 2) Let  $\tilde{a}_{ij} < \alpha_{ij}$ . For  $t \in \langle 0, 1 \rangle$  define the real function

$$\varphi_{ij}(t) = \bar{a}_{ij}x_j + \sum_{k+j} (a_{ik} + t(\bar{a}_{ik} - \underline{a}_{ik})) x_k$$
.

Then  $\varphi_{ij}(0) < x_i - y_i$ ,  $\varphi_{ij}(1) = (\bar{A}x)_i \ge x_i - y_i$ , thus there exists a  $\tau_{ij} \in (0, 1)$  such that  $\varphi_{ij}(\tau_{ij}) = x_i - y_i$ . Now put  $a_{ij}^2 = \bar{a}_{ij}$ ,  $a_{ik}^2 = a_{ik} + \tau_{ij}(\bar{a}_{ik} - a_{ik})$  for  $k \ne j$ ,  $a_{ik}^2 = a_{ik}^*$  for  $l \ne i$ ,  $k = 1, \ldots, n$ ,  $A^2 = (a_{ik}^2)_{n \times n}$ . Then  $A^2 \in M$  and  $(A^2)_{ij} = \bar{a}_{ij} = \min(\bar{a}_{ij}; \alpha_{ij})$ , Q. E. D.

We say the ij-th coefficient (of the matrix  $A_0$ ) is corrected if  $\underline{a}_{ij} < \underline{a}_{ij}$  or  $\tilde{a}_{ij} < a_{ij}$ , i.e. if  $\tilde{a}_{ij} = \underline{a}_{ij} < a_{ij} = \underline{a}_{ij}$ . Taking  $p_{ij} = (a_{ij} - \underline{a}_{ij}) \, x_j$ ,  $i, j = 1, \ldots, n$ , it follows from (1) and (2) that

$$\tilde{a}_{ij} < \bar{a}_{ij}$$
 if and only if  $p_{ij} > ((E - \underline{A}) x - y))_i$ , (3)

$$\underline{a}_{ij} < \underline{a}_{ij} \quad \text{if and only if} \quad p_{ij} > (y - (E - \overline{A}) x)_i,$$
 (4)

so that the ij-th coefficient is corrected if and only if

$$p_{ij} > \min\left(\left(\left(E - \underline{A}\right)x - y\right)_i, \quad (y - \left(E - A\right)x)_i\right). \tag{5}$$

Theorem 2: Let (i), (ii) hold. Then

a) if  $\tilde{a}_{ij} < \bar{a}_{ij}$  for some i, j, then  $a_{ik} = a_{ik}$  for all  $k \neq j$ ,

b) if  $\underline{a}_{ij} < \underline{a}_{ij}$  for some i, j, then  $\tilde{a}_{ik} = a_{ik}$  for all  $k \neq j$ .

Proof: Suppose there exist  $i, j, k, j \neq k$ , such that  $\tilde{a}_{ij} < a_{ij}, a_{ik} < a_{ik}$ . Using (3) and (4), we obtain

$$p_{ij}+p_{ik}>((\bar{A}-\underline{A})x)_i=\sum_i p_{ii}\geq p_{ij}+p_{ik}$$
,

which is a contradiction.

Theorem 3: Let (i), (ii) be satisfied. Then for i, j, k = 1, ..., n,

$$0 < p_{ij} \leq p_{ik} \Rightarrow \frac{\tilde{a}_{ik} - a_{ik}}{a_{ik} - a_{ik}} \leq \frac{\tilde{a}_{ij} - a_{ij}}{a_{ij} - a_{ij}}.$$
 (6)

Proof: We may suppose that  $k \neq j$  and  $\tilde{a}_{ij} = \underline{a}_{ij} < a_{ij} = \underline{a}_{ij}$ , because the other cases are trivial. Then two cases are possible:

1)  $\tilde{a}_{ij} < \tilde{a}_{ij}$ . Then  $p_{ik} \ge p_{ij} > ((E-A)x - y)_i$ , thus  $\tilde{a}_{ik} < a_{ik}$  and according to Theorem 2,  $\tilde{a}_{ij} = \underline{a}_{ij}$ ,  $\underline{a}_{ik} = \underline{a}_{ik}$ . Now we have  $(\tilde{a}_{ij} - \underline{a}_{ij}) x_j = ((E-A)x - y)_i = (\tilde{a}_{ik} - \underline{a}_{ik}) x_k$ , so that

$$\frac{\tilde{a}_{ik} - a_{ik}}{\tilde{a}_{ik} - \tilde{a}_{ik}} = \frac{((E - A)x - y)_i}{p_{ik}} \le \frac{((E - A)x - y)_i}{p_{ij}} = \frac{\tilde{a}_{ij} - a_{ij}}{\tilde{a}_{ij} - a_{ij}}.$$

2)  $\underline{a}_{ij} < \underline{a}_{ij} < \underline{a}_{ij}$ . In the similar manner we obtain  $(\tilde{a}_{ij} - \underline{a}_{ij})x_j = (y - (E - \bar{A})x)_t = (\tilde{a}_{ik} - \underline{a}_{ik})x_k$  and it again

Note that the *ij*-th coefficient is corrected if and only if  $\frac{\tilde{a}_{ij}-a_{ij}}{\bar{a}_{ij}-a_{ij}} < 1$ . Thus Theorem 3 implies that if  $0 < p_{ij} \le p_{ik}$  and the ij-th coefficient is corrected, then the ik-th one is also corrected. This result can be used also in the following way: let

$$p_{iji} = \max p_{ij}, \quad i = 1, \ldots, n;$$

then if the  $ij_i$ -th coefficient is not corrected, then no coefficient of the i-th row is corrected. Note also that if  $0 < p_{ij} = p_{ik}$ , then  $\frac{\tilde{a}_{ij} - \tilde{a}_{ij}}{\tilde{a}_{ij} - \tilde{a}_{ij}} = \frac{\tilde{a}_{ik} - \tilde{a}_{ik}}{\tilde{a}_{ik} - \tilde{a}_{ik}}$ .

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## **Multicommodity Location Problems**

Recently A. Warszawski has developed a procedure for the multicommodity location problem with no capacity  $constraints \ on \ the \ supply \ sources. \ The \ purpose \ \hat{of} \ this \ contribution \ is \ to \ describe \ how \ this \ procedure \ can \ be \ extended$ to the problems with nontrivial capacity constraints.

The single-commodity problem involves locating of one or more supply sources within a set of given possible sites so as to minimize the sum of the setup costs and the transportation costs to a set of destinations with given demands. Denoting by m the number of possible sites, n the number of destinations,  $g_i$  the setup costs associated with placing the supply source at site i,  $A_i$  the capacity of the supply source if placed at site i,  $b_i$  the demand at