

Limits of dense graph sequences

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Lovász, Szegedy *JCTB*'06 (Fulkerson Prize'12)

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idea: convergence notion for sequences of finite graphs
compactification of the space of finite graphs \Rightarrow
... *graphons* symmetric Lebesgue-m. functions $\Omega^2 \rightarrow [0, 1]$

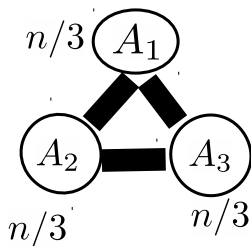
Why? same story as with \mathbb{Q} vs \mathbb{R} : only the latter allows
reasonable e.g. variational and integral calculus
for example $\operatorname{argmin}(x^3 - 2x)$

Limits of dense graph sequences: an abstract approach

F is a “fixed graph” of order k , G is “large” of order n

We define **subgraph density** $t(F, G)$:

$$t(F, G) := \frac{\# \text{ copies of } F \text{ in } G}{\binom{n}{k}} = \mathbf{P}[G[\text{random } k\text{-set}] \cong F]$$



$$k = 2: t(\bullet\bullet, G) = \frac{1}{3} \quad t(\bullet\text{---}\bullet, G) = \frac{2}{3}$$

$$k = 3: t(\bullet\bullet\bullet, G) = \frac{1}{9} \quad t(\bullet\text{---}\bullet\bullet, G) = 0$$

$$t(\bullet\text{---}\bullet\text{---}\bullet, G) = \frac{2}{3} \quad t(\triangle, G) = \frac{2}{9}$$

$$k = 4: t(\boxtimes, G) = 0 \quad \dots$$

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A sequence of graphs G_1, G_2, \dots **converges** if for each F , the sequence $t(F, G_1), t(F, G_2), \dots$ converges.

We get a **limit object** Ψ , $t(F, \Psi) = \lim_n t(F, G_n)$.

Why dense graph sequences?

If the proportion of edges $\searrow 0$ ($\lim \frac{e(G_n)}{n^2} = 0$) we get a trivial limit.

That is, the theory is void for trees, planar graphs, ...

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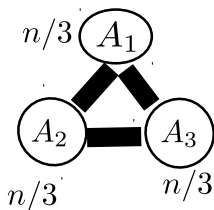
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Razborov'07 **flag algebras** (next slide)

Extremal graph theory and Razborov's flag algebras I



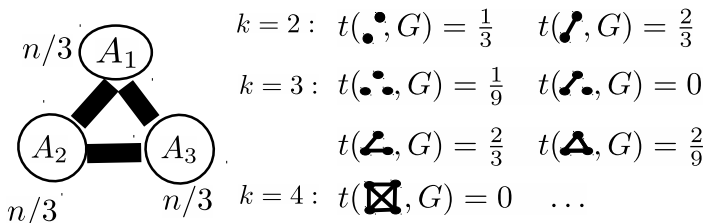
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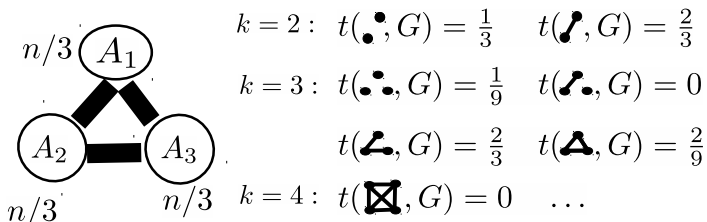
Extremal graph theory and Razborov's flag algebras I



Theorem (\approx Turán 1941) For each $\epsilon > 0$ there exists $\delta > 0$: If an n -vertex graph has more than $(\frac{2}{3} + \epsilon) \binom{n}{2}$ edges then it contains $> \delta$ -proportion of \boxtimes 's.

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Razborov: But lets rather study these relations on the limit space!

Approach to proving Turán: Suppose the theorem is false.

G_1, G_2, \dots all contain $(\frac{2}{3} + \epsilon)$ -proportion of edges but proportion of \boxtimes 's tends to 0. Pass to a subsequential limit Ψ . $t(\bullet, \Psi) \geq \frac{2}{3} + \epsilon$ and $t(\boxtimes, \Psi) = 0$. Derive a contradiction.

Extremal graph theory and Razborov's flag algebras II

Razborov provides tools for deriving relations that hold on the limit object (Cauchy-Schwarz calculus, variational calculus, ...)

First application: Razborov'08 (**AMS Robbins Prize'12**) solves the **triangle density problem** of Lovász and Simonovits 1983:

Suppose a graph has a given proportion of edges. What proportion of triangles can it have?

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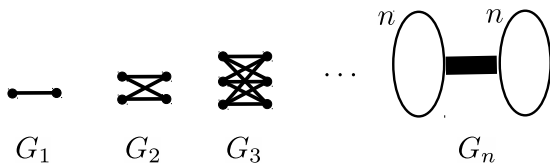
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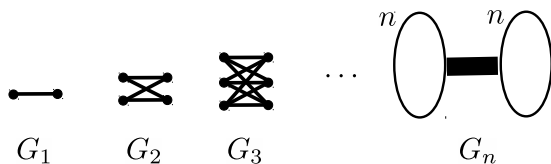
- ▶ H.-Kráľ'-Norine'09: [Caccetta-Häggkvist conj.](#) (progress)
- ▶ H.-Hatami-Kráľ'-Norine-Razborov'11 [conjecture of Erdős 1984](#)
- ▶ HHKNR'11 [conjecture of Jagger-Štovíček-Thomason 1996](#)
- ▶ ... and many more

Hatami-Norine *J.AMS'11* deciding whether an inequality between subgraph densities holds for all graph limits is undecidable

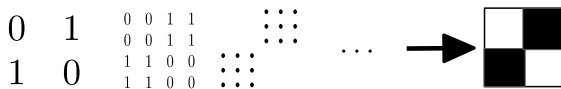
Graphons I



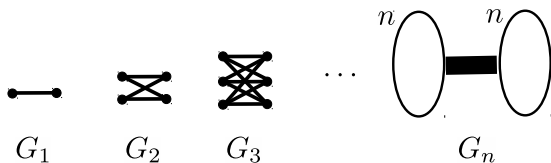
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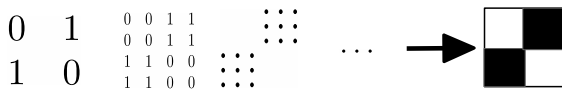
Represent these graphs by their adjacency matrices:



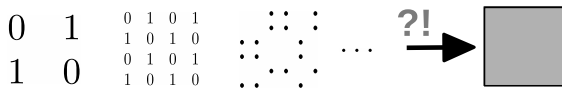
Graphons I



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... works if you do things the right way. But, ...



In general Szemerédi's Regularity Lemma can be used to determine "the right way" of ordering the vertices.

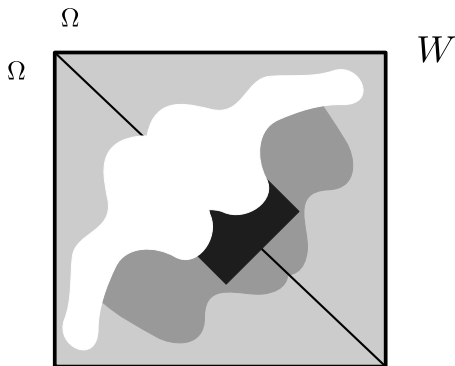
Graphons II

A **graphon** is a symmetric Lebesgue-m. function. $W : \Omega^2 \rightarrow [0, 1]$.

Theorem (Lovász–Szegegy) sampling conv. \Leftrightarrow graphical conv.

Theorem (L–Sz.) Every graphon W can be achieved in the limit.

Proof:



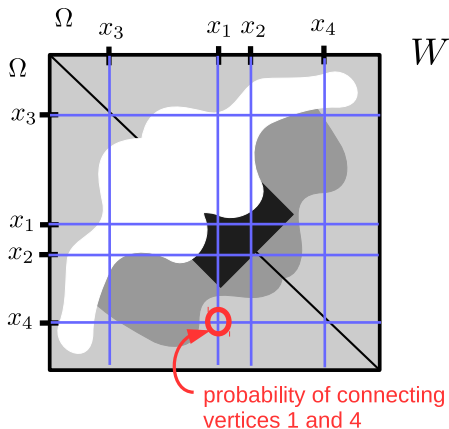
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Proof: Random graphs G_1, G_2, \dots ; $V(G_n) = \{1, \dots, n\}$; sample $x_1, \dots, x_n \in \Omega$ and connect i with j with probability $W(x_i, x_j)$.



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It can be shown that almost surely, $G_1, G_2, \dots \rightarrow W$. □

$\mathbb{G}(n, W)$ as a generalization of the Erdős–Rényi model $\mathbb{G}(n, p)$.

Interesting model *per se!*

Bollobás–Janson–Riordan'07, Doležal–H.–Máthé'15