ACLT: Algebra, Categories, Logic in Topology

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- Grothendieck's generalized topological spaces (toposes)

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4. Toposes and geometric reasoning

How to "do generalized topology"

**TACL** Tutorial course

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4. Toposes and geometric reasoning

Classifying topos for T represents "space of models of T"

Outline of course

- 1. Sheaves: Continuous set-valued maps
- 2. Theories and models: Categorical approach to many-sorted first-order theories.
- 3. Classifying categories: Maths generated by a generic model
- Toposes and geometric reasoning: How to "do generalized topology"

Constructive! \_\_\_\_\_ No choice No excluded middle It is "geometric mathematics freely generated by generic model of T"

Map = geometric morphism = result constructed geometrically from generic argument

Bundle = space constructed geometrically from generic base point - fibrewise topology

Arithmetic universes for when you don't want to base everything on Set

# Point-free topology

Point-set topology says:

- 1 define collection of points as set
- 2 define topology, using open subsets

Point-free topology describes points and opens in one single structure

- a geometric theory
- points are models
- opens are propositions
- sheaves are "derived sorts"



open sets of completely prime filters are those of the form  $\{F \mid U \in F\}$  for some U

If Y also sober continuous maps X -> Y are determined by the frames: They are the frame homomorphisms  $\Omega Y \rightarrow \Omega X$ 

If f: X -> Y continuous, then inverse image f^{-1} is frame homomorphism

Other way round: note that completely prime filters of  $\Omega X$  are frame homomorphisms  $\Omega X \rightarrow \Omega =$  frame of truth values. Composing with a frame homomorphism  $\Omega Y \rightarrow \Omega X$ gives continuous map X -> Y

#### Locales

Any frame A can be treated as a point-free space.

From that point of view call it a locale

Continuous maps between locales are just frame homomorphisms backwards

As propositional geometric theory:

- A = signature. Write (a) for a as propositional symbol

1.

- axioms

$$(a) \vdash (b)$$
  
 $(a) \land (b) \vdash (a \land b)$   
 $T \vdash (T)$   
 $(V; a;) \vdash V; (a;)$ 

joins, finite meets in A become disjunctions, conjunctions in logic

models = completely prime filters of A Lindenbaum algebra = A

Continuous maps are geometric morphisms

For propositional case:

presheaves with pasting

Theorem Let A, B be frames, let Sh(A), Sh(B) be their toposes of sheaves. Then there is a bijection between

- frame homomorphisms B -> A
- isomorphism classes of geometric morphisms Sh(A) -> Sh(B)

#### Proof idea

Elements of A (opens) correspond to subsheaves of 1. If f:  $Sh(A) \rightarrow Sh(B)$  is a geometric morphism, then f\* maps opens of B to opens of A, and gives a frame homomorphism B -> A.

Every sheaf is a colimit of opens, so f\* is determined up to isomorphism by its action on opens.

Moreover, an arbitrary frame homomorphism gives rise to a geometric morphism.

Continuous maps are geometric morphisms

For propositional theories:

Geometric morphisms match continuous maps for locales - which match continuous maps for sober spaces

For general theories:

*Define* continuous map to be geometric morphism (between classifying toposes)

Remember: geometric morphisms are equivalent to -

- functors in the opposite direction
- preserving finite limits and arbitrary lolimits



Reasoning in point-free logic



Get map (geometric morphism) f: S[T\_1] -> S[T\_2]

Geometric morphism transforms points

#### Idea

Classifying topos S[T\_1] somehow "is" space of models of T. Its points are the models of T - but where?

Model in E is equivalent to geometric morphism M: E -> S[T\_1]

$$\mathcal{E} \longrightarrow \mathcal{B}[\mathbb{T}_1] \longrightarrow \mathcal{B}[\mathbb{T}_2]$$

Composing them gives model f(M) of T\_2 in E

f transforms models of T\_1 to models of T\_2, in any E.

## Reasoning in point-free logic



Role of geometricity

Construction of f(M\_G) out of M\_G was geometric

Non-geometric constructions (e.g. exponentials, powerobjects) are also available in S[T\_1]



They too could construct model of T\_2 and give geometric morphism f.

But they wouldn't be preserved by M\* - they wouldn't construct M\*(f(M\_G)) out of M

If construction on generic model M\_G is geometric, then it is uniform

- same construction also applies to all specific models M.

# To define a continuous map

- from space of T\_1 models to space of T\_2 models
- 1. Take as argument M a T\_1 model
- 2. Construct a T\_2 model f(M), geometrically
- 3. No continuity proof neededgeometricity guarantees continuity

"geometricity is continuity"

Aspects of continuity

For ordinary spaces:

Continuous maps preserve specialization order

For sober spaces:

Have all directed joins of points

Continuous maps preserve them

cf. Scott continuity



Aspects of continuity

For generalized spaces

For ordinary spaces: Continuous maps preserve specialization order For sober spaces: Have all directed joins of points Continuous maps preserve them

Specialization order becomes homomorphisms of models

Continuous maps (geometric morphisms) are functorial on points

 $\mathcal{E} \xrightarrow{\mathcal{M}_{2}} SEF_{1} \xrightarrow{\mathcal{F}} SEF_{2}$ Itomomorphism  $\mathcal{M}_{1} \rightarrow \mathcal{M}_{2}$  (models)  $\approx Natural transformation <math>\mathcal{M}_{1} \rightarrow \mathcal{M}_{2}$  (geometric  $\Rightarrow Natural transformation for <math>\mathcal{M}_{1} \rightarrow \mathcal{M}_{2}$  (morphisms)  $\Rightarrow Natural transformation for <math>\mathcal{M}_{1} \rightarrow \mathcal{F}\mathcal{M}_{2}$  $\approx ttomomorphism f(\mathcal{M}_{1}) \rightarrow \mathcal{F}\mathcal{M}_{2}$  Aspects of continuity

For generalized spaces

For ordinary spaces: Continuous maps preserve specialization order For sober spaces: Have all directed joins of points Continuous maps preserve them

Instead of directed joins, consider filtered colimits

Defn A category C is filtered if any finite diagram in it has a cocone

1. C has at least one object

Empty diagram has a cocone

2. For any two objects i and j, there are morphisms out to a third k

3. For any two parallel morphisms, there is a third that composes equally with them  $( \rightarrow )$ 

e.g. a poset is filtered iff it is directed

# **Filtered colimits**

A filtered colimit is a colimit of a filtered diagram, i.e. a functor from a filtered category

In Set: Suppose X: C -> Set a filtered diagram.



Filtered colimits

Facts

Filtered colimits commute with finite limits

For a geometric theory T, filtered colimits of models can be found by taking filtered colimits of the carriers

For any two Grothendieck toposes, we have filtered colimits of geometric morphisms between them

Those filtered colimits are preserved by composition on either side

As point transformers, geometric morphisms preserve filtered colimits

# **Object classifier**

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Let O be theory with one sort and nothing else
Model = set "set" here = object in a topos
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S[O] is the object classifier, the "space of sets"

Map F: E -> S[O] = (1) "continuous set-valued map on E" definition of classifying topos = (2) object of E

If E = Sh(X), this justifies "sheaf = continuous set-valued map"

If x is point of X, then F(x) = stalk at x (fibre of local homeomorphism)

Proof method Treat x as map x:  $1 \rightarrow X$ , calculate  $x^*(F)$ . It's an object of Sh(1) = Set

# Objects of S[O]

Intuition Object of S[O] is - continuous map from "space of sets" to itself

Continuity is at least functoriality and preservation of filtered colimits

- functor Set -> Set preserving filtered colimits

Every set is a filtered colimit of finite sets Set is ind-completion of Fin

- functor Fin -> Set Not a proof, but ...

Theorem S[O] is equivalent to [Fin, Set]

More generally For any cartesian theory T, S[T] is equivalent to category of set-valued functors on category of finitely presented T-models Reasoning in point-free topology: examples

+: RXR -> R Dedekind sections, e.g. (L\_x, R\_x) et x, y E R Then scrye R where Lxry = Eqtr ge Loc, re Ly} Rxry = Eqtr ge Rx, re Ryj

# Why is real line R geometric?

1. Propositional theory

Propositional symbols for subbasic open intervals  $(q, \infty)$ ,  $(-\infty, q)$  (q rational)

Axioms to express relations between these, e.g.

$$(q, \infty) \vdash \bigvee (q', \infty)$$
 infinite disjunction!  
 $q \leq q'$ 

2. First order theory

- sorts N, Q

need infinite disjunctions to do this

- structure and axioms to force them to be modelled as natural numbers and rationals
- predicates L(q:Q) and R(q:Q) for left and right parts of a Dedekind sections
- appropriate axioms

Can show (1) and (2) are equivalent - mutually inverse maps between them.

Reasoning in point-free topology: examples



Fibrewise topology

Let X be a space

Imagine topology "continuously parametrized by points x of X" Hope to do topology as usual, but with parameters x everywhere

e.g. define spaces Y\_x
bundle them together to make space Y with map p: Y -> X
Y\_x = fibre of p over x

- each fibre Y\_x has given topology
- but what about topology of Y across the fibres?
- somehow comes from "continuity" of x |-> Y\_x ???

Makes sense if -

- spaces are point-free
- construction of Y\_x is geometric

James: "Fibrewise Topology" - classical, point-set



- Externally: get theory T2, models = pairs (M, N) where
- M a model of T1
- N a model of F(M)

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Map p: S[T2] -> S[T1]
- (M,N) |-> M
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Fibres

Can generalize to models in other toposes

[F(M)

M

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Suppose M a model of T1 in Set = Sh(1) Get map M: Set -> S[T1]

S[F(M)] = "fibre of p over M"

Get square ·

Fact It's a (pseudo)pullback in category Top of Grothendieck toposes, geometric morphisms

cf. pullback square for ordinary fibres

Geometric morphisms: two views

f: S[T2] -> S[T1]

1. Map

argument y |-> result f(y)

2. Bundle

base point x |-> fibre f^{-1}{x}

Technicality when generalizing to elementary toposes:

for bundle view, f must be bounded

Localic bundle theorem (Joyal and Tierney)

Let E be any topos

Then there is a duality between:

- internal frames in E

- bundles p: F -> E that are localic

Technically: every object of F is quotient of a subobject of some p\*(X)

Hence:

internal locale maps correspond to external bundle morphisms

#### Note -

frames are not part of the geometric mathematics

Need powersets to construct them

Frame presentations (propositional geometric theories) are geometric

# Geometric properties of bundles

Some topological properties C of spaces - e.g. discreteness, compactness, local connectedness, ... are preserved under (pseudo)pullback of bundles

Then say C is a geometric property

Say bundle p: Y -> X is fibrewise C iff it is internally C in Sh(X)

Then all its fibres are also C

## e.g. discreteness

Discrete space = set (or object in topos)

Object X (in topos E)

- powerobject P(X) = frame for discrete space
- geometric theory T signature: propositional symbols s\_x (x in X) axioms:

point = model of theory = singleton subset of X = element of X open = formula = arbitrary subset of X (discrete topology)

Corresponding bundle over E is fibrewise discrete Note: in this case bundle topos is equivalent to slice E/X.

Local homeomorphisms (Joyal and Tierney)

Let A be an internal frame in topos E, let p: F -> E be the bundle

Theorem The following are equivalent.

- p is fibrewise discrete (A is isomorphic to some P(X))
- p is open and so is  $\Delta$ : F -> F x\_E F



cf. Lecture 1! Use this as point-free definition of local homeomorphism

(But first have to define open maps.)

Fibrewise discrete = local homeomorphism

Point-set topology is inadequate!

With respect to base space X:

- space = bundle over X
- set = discrete space = local homeomorphism over X

"set of points" for a bundle p = approximation by a local homeomorphism

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Sometimes no approximation is good enough.

For local homeomorphism, specialization in base gives map between fibres

pt(p) must have empty fibre over top - so must also have empty fibre over bottom



Reasoning in point-free topology: examples

Spec: [BA] -> Spaces



Reasoning in point-free topology: examples

B a pt of space of Boolean algebras
internal point-free space
external bundle

Spec(B) is fibre over B

Geometricity => construction is uniform:

- single construction on generic B
- also applies to specific B's
- get those by pullback

Spec(B) ~ [BA+ prime filler] & J EBA] B

pullback
= generalized fibre of generalized point

Localic hyperspaces (powerlocales)

Very useful geometrically

Point-free treatment of Vietoris topology - split into two halves

Lower powerlocale P\_L(X)

- point = "overt, weakly closed sublocale of X"

- specialization order = sublocale order

Upper powerlocale P\_U(X)

Think: compact, and up-closed under specialization order

Think: closed subspace

- point = "compact, fitted sublocale of X"

- specialization order = opposite of sublocale order

Both work internally in geometric way

- giving fibrewise hyperspaces of bundles

Hyperspace applications - examples Lower powerlocale P\_L(X) - point = "overt, weakly closed sublocale of X"

- X is compact iff
- there is a point K of P\_U(X)
- such that

K is biggest compact fitted sublocale, and can show it must be the whole of X.

p: Y -> X is open iff - there is a section L of  $(P_L)^X(Y)$ - such that  $\left( \frac{1}{2} \times (x) \right) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \eta$  Upper powerlocale P\_U(X)

- point = "compact, fitted sublocale of X"

- specialization order = sublocale order

- specialization order = opposite of sublocale order (



fibrewise lower powerlocale of p p open iff ...

intuition:

 $Y = P_{L}^{x}(Y)$   $P = \frac{q/1}{L}$ For any Y: subbase of opens for PLY DV: {Al AnV = \$} (V open in Y) For  $p: Y \rightarrow X$  - write  $Y_{\infty}$  for fibre over XIf V open in Y,  $VnY_{\infty}$  open in  $Y_{\infty}$   $S(VnY_{\infty})$  open in  $P_{\perp}(Y_{\infty})$   $(S(VnY_{\infty}))_{X \in X}$  open in  $P_{\perp}(Y)$ Inequality  $[a_{1}x_{1}x_{2}] = L^{\circ}q$  says  $L(x) = Y_{x}$   $L^{-1} ( \langle V (V \wedge Y_{x}) \rangle_{x \in x} = \{ sc | Y_{x} \in \langle V (V \wedge Y_{x}) \} : \{ x | V \wedge Y_{x} \neq \phi \}$  = image p(V)  $L continuous \Rightarrow p open$ 

Topos theory to do fibrewise topology of bundles

Programme:

Carry out topology in a way that is

- point-free (but can still use points!)
- geometric (reasoning must be constructive)
- In scope of declaration "Let x be a point of X"
- x = generic point of X (in Sh(X))
- space = generic fibre of bundle over X
- geometric properties of space hold fibrewise

Hyperspaces are very useful

- internal hyperspace works fibrewise

## Conclusions

Grothendieck's generalized spaces:

- can understand topology and continuity much more broadly than before

- sheaves are more important than opens
- sheaves provide a rich geometric mathematics for performing generic constructions on generic points

Even for ungeneralized spaces:

Topos theory constructive (geometric) point-free reasoning using sheaves over base

gives natural fibrewise topology of bundles - topology parametrized by base point Further reading

Topos theory -Mac Lane and Moerdijk "Sheaves in geometry and logic" Johnstone "Sketches of an elephant" + a readers' guide to those two -Vickers "Locales and toposes as spaces"

Constructive reasoning for locales -Joyal and Tierney "An extension to the Galois theory of Grothendieck"

Powerlocales -Vickers - various papers; in particular "The double powerlocale and exponentiation"