# Clones (3&4)

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## Galois connections

Let *A*, *B* be sets,  $R \subseteq A \times B$ . For any  $S \subseteq A$  and any  $T \subseteq B$  let

•  $S^u := \{ b \in B \mid \forall a \in S : aRb \}$ 

• 
$$T^{\ell} := \{ a \in A \mid \forall b \in T : aRb \}.$$

Then

- the maps  $T \mapsto T^{\ell}$  and  $S \mapsto S^{u}$  are  $\subseteq$ -antitone.
- the maps S → S̄ := S<sup>uℓ</sup> and T → T̄ := T<sup>ℓu</sup> are closure operators (S ⊆ S̄ = S̄)

$$\triangleright S^{u\ell u} = S^u, T^{\ell u\ell} = \ell.$$

Usually it is of interest to characterize the family of *closed* sets  $\{S \mid S = \overline{S}\}$  and the closure operator "from below".

## Galois connections, examples

Let *A*, *B* be sets,  $R \subseteq A \times B$ .  $S^{u} := \{b \in B \mid \forall a \in S : aRb\} = \bigcap_{a \in S} \{b \in B \mid aRb\}$   $T^{\ell} := \{a \in A \mid \forall b \in T : aRb\} = \bigcap_{b \in T} \{a \in A \mid aRb\}$ Examples

- ► *A* = vector space, *B* = dual space = set of linear forms.  $aRb \Leftrightarrow b(a) = 0$ .  $S \subseteq A \Rightarrow S^{u\ell}$  = linear hull of *S*.
- ► A = all formulas, B = all structures,  $aRb \Leftrightarrow b \models a$  (the formula *a* holds in the structure *b*).  $\overline{S} =$  all consequences of  $S = \{a : S \models a\}$
- ► A = operations on X, B = relations,  $f R \rho \Leftrightarrow f \triangleright \rho$ .  $S^{u\ell} = \langle S \rangle$  = clone generated by S.

## k-ary clones, clones

$$\mathbb{O}^{(k)} := \{ f \mid f : X^k \to X \}. \ \mathbb{O}_X := \bigcup_{k=1}^{\infty} \mathbb{O}_X^{(k)}.$$

#### Definition (*k*-ary clone)

A *k*-ary clone on X is a set  $T \subseteq O_X^{(k)}$  which is closed under "composition" and contains the *k* projections.

### Definition (Clone)

A *clone* on *X* is a set  $T \subseteq \mathcal{O}_X = \bigcup_{k=1}^{\infty} \mathcal{O}_X^{(k)}$  which is closed under "composition" and contains all projections.

## Definition (Composition)

Let  $f \in \mathcal{O}^{(k)}$ ,  $g_1, \ldots, g_k \in \mathcal{O}^{(m)}_X$ .  $f(g_1, \ldots, g_k)(\vec{x}) := f(g_1(\vec{x}), \ldots, g_k(\vec{x}))$  for all  $\vec{x} \in X^m$ . If *C* is a clone, then  $C^{(k)} := C \cap \mathcal{O}^{(k)}$  is a *k*-ary clone, the *k*-ary fragment of *C*.

# Vector-valued operations

*C* is a clone:  $f, g_1, \ldots, g_k \in C \Rightarrow f(g_1, \ldots, g_k) \in C$ . We can view  $(g_1, \ldots, g_k)$  as a single function  $\vec{g} : X^m \to X^k$ , and write  $f \circ \vec{g}$  instead of  $f(g_1, \ldots, g_k)$ .

#### Definition

For any set  $S \subseteq O_X$  let  $\tilde{S}$  be the set of all operations  $f: X^k \to X^n$  with the property that all "components" are in S:

$$\tilde{\boldsymbol{S}} := \bigcup_{k,n} \{ f : \boldsymbol{X}^k \to \boldsymbol{X}^n \mid \forall i \in \{1, \dots, n\} : \ \pi_i^n \circ f \in \boldsymbol{S}) \}$$

where  $\pi_i^n : X^n \to X$  is the *i*-th projection function.

The set S is a clone iff  $\tilde{S}$  contains all projection functions and is closed under composition:

$$\forall g: X^m 
ightarrow X^k \ \forall f: X^k 
ightarrow X^n: (f, g \in \tilde{S} \Rightarrow f \circ g \in \tilde{S})$$

## Examples of clones

- Every subset S ⊆ O<sub>X</sub> will generate a clone (S), the smallest clone containing S.
- For any relation ρ ⊆ X<sup>n</sup>: Pol(ρ) := {f ∈ O<sup>(|f|</sup><sub>X</sub> ⊳ ρ} is a clone.
- ► For any relation  $\rho \subseteq X^{K}$  (K infinite), Pol( $\rho$ ) is a clone.
- For any set R of relations, POL(R) := ∩<sub>ρ∈R</sub> Pol(ρ) is a clone.
- ⟨C⟩ = POL(INV(C)), where INV(C) := ∩<sub>f∈C</sub> Inv(f), Inv(f) := {ρ | f ⊳ ρ}. (For infinite X, need to allow infinitary relations; operations still have finite arity!)

# The lattice of all clones on X

For finite *X*,  $\mathcal{O}_X$  is countable. For infinite *X* of size  $\kappa$ ,  $\mathcal{O}_X$  has  $2^{\kappa}$  elements.

#### Definition

For any nonempty set X let CI(X) be the set of all clones on X. (CI(X) is a subset of the power set of  $\mathcal{O}_X$ .)

- CI(X) is a complete lattice. (meet = intersection, join = clone generated by union)
- CI(X) is Countable for |X| = 2. (Post's lattice. wikipedia!)
- CI(X) is of size  $|\mathbb{R}| = 2^{\aleph_0}$  for X finite with > 2 elements.
- For infinite X of size κ: |Cl(X)| ≤ 2<sup>2<sup>κ</sup></sup>. In fact: = 2<sup>2<sup>κ</sup></sup>. (Later)

# **Minimal clones**

### Definition

We call a clone *M* minimal if  $J \subsetneq M$  (*J* is the smallest clone, containing only the projections), but there is no clone *D* with  $J \subsetneq D \subsetneq M$ .

The minimal clones are the atoms of the clone lattice.

An operation *m* is minimal iff  $\langle m \rangle$  is a minimal clone.

Instead of minimal clones we consider minimal operations. If *m* is minimal, then  $\forall f \in \langle m \rangle \setminus J : m \in \langle f \rangle$ .

- If *m* is unary, then have *m* ∈ ⟨*m<sup>j</sup>*⟩ for all *j* except if *m<sup>j</sup>* = *id*. Hence *j*<sup>2</sup> = *id* ("retraction"), or *m* is a permutation of prime order.
- ► If *m* not essentially unary, then *m* must be idempotent. m(x,...,x) = m.

# Minimal operations, examples

- Every constant operation.
- Every permutation whose order is a prime number.
- The meet operation of any meet-semilattice.
- The median operation in any linear order.
- ... (many more. Some necessary conditions known, but no explicit criterion.)

#### Fact

If X is finite, then there are finitely many minimal operations. Every clone  $\neq$  J contains a minimal clone.

(This is not true for infinite sets. Let  $s : \mathbb{Z} \to \mathbb{Z}$  be defined by s(x) = x + 1, then every non-projection in  $\langle s \rangle$  is of the form  $s^j$   $(j \in \{1, 2, ...\}$ , and none of them is minimal, as  $\langle s^{2j} \rangle \subsetneq \langle s^j \rangle$ .)

# Complete sets

Theorem For every X:  $\langle \mathfrak{O}_X^{(2)} \rangle = \mathfrak{O}_X$ .

#### Proof for infinite X.

- Let  $p_2: X^2 \to X$  be a bijection.
- ▶ Find bijections  $p_j : X^j \to X$  for j = 3, 4, ..., with  $p_j \in \langle 0^{(2)} \rangle$ . For example,  $p_3(x, y, z) := p_2(x, p_2(y, z))$ .
- ► For every  $f: X^k \to X$ , let  $\hat{f} := f \circ p_k^{-1}$ . So  $f(\vec{x}) = \hat{f}(p_k(\vec{x}))$  for all  $\vec{x} \in X^k$ . As  $\hat{f}$  is unary,  $\hat{f} \in \langle 0^{(2)} \rangle$ .
- From  $\hat{f} \in \langle \mathbb{O}^{(2)} \rangle$  and  $p_k \in \langle \mathbb{O}^{(2)} \rangle$  conclude  $f \in \langle \mathbb{O}^{(2)} \rangle$ .

# Complete sets

For every X:  $\langle \mathfrak{O}_X^{(2)} \rangle = \mathfrak{O}_X$ .

Proof for finite X ("Lagrange interpolation").

Let  $(X, +, \cdot, 0, 1)$  be a finite lattice with smallest element 0 and greatest element 1. So  $x + 0 = 0 + x = x = 1 \cdot x$  for all x.

- For each a ∈ X let χ<sub>a</sub> : X → X be the characteristic function of the set {a}. So χ<sub>a</sub> ∈ O<sup>(1)</sup> ⊆ (O<sup>(2)</sup>).
- ► For each  $\vec{a} \in X^k$  let  $\chi_{\vec{a}} : X^k \to X$  be the characteristic function of  $\{\vec{a}\}$ :  $\chi_{\vec{a}} = \prod_i \chi_{a_i}(x_i)$ . So  $\chi_{\vec{a}} \in \langle 0^{(2)} \rangle$ .
- ▶ For any  $b \in X$  let  $c_b \in O^{(1)}$  be constant with value b.
- ► Every operation  $f \in \mathcal{O}^{(k)}$  can now be written as  $f = \sum_{\vec{a} \in X^k} (\chi_{\vec{a}} \cdot c_{f(\vec{a})})$ . So  $f \in \langle \mathcal{O}^{(2)} \rangle$ .

(Remark: This proof also works for strongly amorphous sets.)

# Precomplete clones

#### Definition

A clone  $C \subseteq \mathcal{O}_X$  is "precomplete" (or "maximal") if  $C \neq \mathcal{O}_X$ , but there is no clone *D* satisfying  $C \subsetneq D \subsetneq \mathcal{O}_X$ .

#### Theorem

For any clone  $C \subsetneq \mathfrak{O}_X$  there is a precomplete clone C' with  $C \subseteq C'$ .

(Remark: Not true for infinite sets! At least if the continuum hypothesis holds.)

# Post's lattice

The lattice of all clones on a 2-element set is countably infinite.

It has 5 coatoms ("precomplete" clones) and 7 atoms.



### Precomplete clones, example 1

Let  $\rho$  be a nontrivial unary relation, i.e.  $\emptyset \subsetneq \rho \subsetneq X$ . Then Pol( $\rho$ ) is the set of all operations *f* such that  $\rho$  is a subalgebra of (*X*, *f*). This clone is precomplete.

Proof.

Let  $g: X^k \to X, g \notin \operatorname{Pol}(\rho)$ . Let  $C := \langle \operatorname{Pol}(\rho) \cup \{g\} \rangle$ . We show  $C = \mathcal{O}_X$ . Sufficient:  $C \supseteq \mathcal{O}_X^{(2)}$ . For  $v \in X$ , let  $c_v$  be the constant function with value v. There are  $\vec{a} = (a_1, \dots, a_k) \in \rho^k$ ,  $b \notin \rho$  with  $g(\vec{a}) = b$ , So  $c_b = g(c_{a_1}, \dots, c_{a_k})$  is in C. For  $f \in \mathcal{O}_X^{(2)}$  define  $\hat{f}(x_1, x_2, y) := \begin{cases} x_1 & \text{if } y \in \rho \\ f(x_1, x_2) & \text{if } y \notin \rho \end{cases}$ . So  $\hat{f} \in C$ . Now  $f = \hat{f}(\pi_1^2, \pi_2^2, c_b)$ , i.e.,  $f(x_1, x_2) = \hat{f}(x_1, x_2, b)$ . So  $f \in C$ .

## Precomplete clones, example 2

 $\sim$  a nontrivial equivalence relation  $\Rightarrow$  Pol( $\sim$ ) is precomplete.

Proof.

For  $\vec{a}, \vec{b} \in X^k$  write  $\vec{a} \sim \vec{b}$  iff  $\forall i \ a_i \sim b_i$ . This is an equivalence relation on  $X^k$ .

Let  $g: X^k \to X$ ,  $g \notin \text{Pol}(\sim)$ . Let  $C := \langle \text{Pol}(\sim) \cup \{g\} \rangle$ . We have to show  $C = \mathcal{O}_X$ . Sufficient:  $C \supseteq \mathcal{O}_X^{(2)}$ .

There is *k* and  $\vec{a} \sim \vec{b} \in X^k$  with  $1 := g(\vec{a}) \not\sim g(\vec{b}) =: 0$ . We claim that for each  $p \in X^2$  there is a function  $\chi_p : X^2 \to X$ which maps *p* to 1, everything else to  $0 \not\sim 1$ . For each  $p \in X^2$  let  $h_p : X^2 \to X^k$  be defined by  $h_p(p) = \vec{a}$ ,  $h_p(x) = \vec{b}$  otherwise. Clearly  $h_p \in Pol(\sim)$ . So  $\chi_p := g \circ h_p \in C$ . (continued on next page)

#### Proof that $Pol(\sim)$ is precomplete, continued.

We started with a clone  $C \supseteq \operatorname{Pol}(\sim)$ . For each  $p \in X^2$  we have found  $\chi_p \in C$ ,  $\chi_p : X^2 \to X$  with  $\chi_p(p) = 1$ ,  $\chi_p(x) = 0$  for  $x \neq p$ . (And  $0 \not\sim 1$ ) Define  $\chi : X^2 \to X^{|X|^2}$  by  $\chi(\vec{x}) = (\chi_p(x) : p \in X^2)$ . So  $\chi \in \tilde{C}$ . Let  $f \in \mathcal{O}_X^{(2)}$  be arbitrary. We will show  $f \in C$ . Define  $\hat{f} : X^{2+|X|^2} \to X$  as follows:

# ▶ $\hat{f}$ is constant on each ~-class. (So $\hat{f} \in Pol(\sim) \subseteq C$ )

• 
$$\hat{f}(\vec{x},\chi(\vec{x})) = f(\vec{x}).$$

This two requirements are compatible, as  $\vec{x} \neq \vec{x}'$  implies that  $\chi(\vec{x}) \not\sim \chi(\vec{x}')$ . Clearly  $f(\vec{x}) = \hat{f}(\vec{x}, \chi(\vec{x}))$ . So  $f \in C$ .

## Precomplete clones, example 3

#### Definition

Let  $r : X \to X$ ,  $f : X^k \to X$ . We say that f commutes with r if:

$$\forall x_1,\ldots,x_k\in X:f(r(x_1),\ldots,r(x_k))=r(f(x_1,\ldots,x_k))$$

Writing  $r^{\bullet}$  for the relation  $\{(x, r(x)) | x \in X\}$ , f commutes with r iff  $f \triangleright r^{\bullet}$ . (We may write  $f \triangleright r$  instead of  $f \triangleright r^{\bullet}$ )

Clearly  $f \triangleright r \Rightarrow f \triangleright r^j$  for all *j*. Hence e.g.  $Pol(r) \subseteq Pol(r^2)$ . But if *r* is a permutation of order *p*, then  $Pol(r) = Pol(r^j)$  whenever *p* does not divide *j*.

#### Theorem

Assume that  $r : X \to X$  is a permutation and all cycles have the same prime length. Then Pol(r) is precomplete.

### Precomplete clones, examples 4,5

- "monotone": Let ρ ⊆ X × X be a partial order with smallest and greatest element.
   Pol(ρ) is the set of all pointwise monotone operations.
- "affine" Assume |X| = p<sup>m</sup>, so wlog X is a finite field X = GF(p<sup>m</sup>).
   Let ρ = {(a, b, c, d) ∈ X<sup>4</sup> | a + b = c + d}. Then Pol(ρ) is the set of all operations f of the form

$$f(x_1,...,x_k) = a_0 + \sum_{i=1}^k \sum_{j=0}^{m-1} x_i^{p^j}$$

All these clones are precomplete.

Clones (3&4)

# Post's lattice, again

The 5 precomplete clones in  $Cl(\{0,1\})$ :

- operations preserving {0}.
- operations preserving {1}.
- monotone operations
- "commuting":  $f(\neg x) = \neg f(x)$ .
- affine operations



# Rosenberg's list

#### Theorem

Let  $X = \{1, ..., k\}$ . Then there is an explicit finite list of relations  $\rho_1, ..., \rho_m$  such that every precomplete clone on X is one of  $Pol(\rho_1), ..., Pol(\rho_m)$ . The list includes

- all "central relations" (generalisations of  $\rho \subsetneq X$ )
- all nontrivial equivalence relations
   (∄ if |X| = 2)
- all prime permutations
- All bounded partial orders
- affine relations (only if  $|X| = p^n$ )
- (others. more complicated but still explicit)

# Rosenberg's list

#### Theorem Let $X = \{1, ..., k\}$ . Then there is an explicit finite list of relations $\rho_1, ..., \rho_m$ such that every precomplete clone on X is one of Pol( $\rho_1$ ), ..., Pol( $\rho_m$ ).

Completeness criterion  $\langle S \rangle \neq \mathfrak{O}_X$  iff there is some  $\rho_i$  from the list with  $\forall f \in S : f \rhd \rho_i$ .

# A complicated interval in the clone lattice

## Definition

Let  $C_{\text{idem}}$  be the clone of all idempotent operations: f(x, ..., x) = x. (Assume  $|X| \ge 3$ .) Find all clones between  $C_{\text{idem}}$  and  $\mathcal{O}_X$ !

#### Example

Let  $Y \subseteq X$ . Then  $C_{\text{idem} \upharpoonright Y} := \{f \mid \forall x \in Y : f(x, \dots, x) = x\}$  is a clone  $\supseteq C_{\text{idem}}$ .

#### Theorem

Every clone between  $C_{idem}$  and  $\mathfrak{O}_X$  is of the form  $C_{idem \upharpoonright Y}$ . Hence: the interval  $[C_{idem}, \mathfrak{O}_X]$  is (anti-)isomophic to the power set of *X*.

(Precomplete clones correspond to singletons,  $O_X$  to  $\emptyset$ .)

# $[C_{idem}, \mathfrak{O}_X]$ , proof sketch

Let *C* be a clone containing all idempotent operations f(x, ..., x) = x. We want to find *Y* such that  $C = C_{idem|Y} = \{f \mid \forall y \in Y : f(y, ..., y) = y\}.$ 

- $fix(f) := \{a \in X \mid f(a, \ldots, a) = a\}, nix(f) := X \setminus fix(f).$
- Let  $R := { nix(f) | f \in C }.$
- R is downward closed.
- R is upward directed, hence an ideal.
- Let Z be the largest element of R,  $Y := X \setminus Z$ .
- So  $C \subseteq C_{\text{idem} \upharpoonright Y}$ .
- If  $nix(f) \subseteq nix(g)$  and  $g \in C$ , then  $f \in C$ .
- Hence  $C = C_{\text{idem} \upharpoonright Y}$ .

# A complicated interval, continued

 $C_{idem}$  = the clone of all idempotent operations: f(x, ..., x) = x. Theorem (X finite) Every clone between  $C_{idem}$  and  $\bigcirc_X$  is of the form  $C_{idem|Y}$ . For infinite X:

Definition For every filter  $\mathcal{F}$  on X, let

$$\mathcal{C}_{\mathcal{F}} := \bigcup_{Y \in \mathcal{F}} \mathcal{C}_{\operatorname{idem} \upharpoonright Y} = \{ f \mid \exists Y \in \mathcal{F} \; \forall y \in Y \; f(y, \ldots, y) = y \}$$

Each  $C_{\mathcal{F}}$  is a clone above  $C_{\text{idem}}$ .

## A complicated interval, conclusion

$$\mathcal{C}_{\mathcal{F}} := \{ f \mid \exists Y \in \mathcal{F} \; \forall y \in Y \; f(y, \ldots, y) = y \}$$

#### Theorem

Let X be any set. Then the map  $\mathcal{F} \mapsto C_{\mathcal{F}}$  is an order-preserving bijection between the filters on X and the clones above  $C_{\mathcal{F}}$ . Ultrafilters correspond to precomplete clones in this interval, and the improper filter corresponds to  $\mathfrak{O}_X$ . (For finite sets, all filters are principal.)

Translation to topology: the interval  $[C_{idem}, \mathcal{O}_X]$  is anti-isomorphic to the family of closed sets of  $\beta X$ , the Čech-Stone compactification of the discrete space *X*. (Precomplete clones correspond to points,  $\mathcal{O}_X$  to  $\emptyset$ .)

## Another complicated interval

Let X be infinite. We will find "very many" clones with trivial unary fragment, i.e., below  $C_{idem}$ , the clone of all idempotent operations. (Unfortunately: no complete classification.) In fact all our operations will be "conservative":

 $f(x_1,\ldots,x_k)\in\{x_1,\ldots,x_k\}.$ 

- ► Let  $(A_i : i \in I)$  be a family of sufficiently independent sets. (In particular: we demand that for any finite  $I_0 \subseteq I$  and any  $j \in I \setminus I_0$  the set  $(\bigcup_{i \in I_0} A_i) \cap (X \setminus A_j)$  contains at least 2 elements. It is possible to find such a family with  $2^{|X|}$  elements, in particular: an uncountable such family.)
- Fix a linear order  $\leq_i$  on  $A_i$ , with minimum operation  $\wedge_i$ .
- Extend  $\wedge_i$  to X by requiring  $x \wedge_i y = x$  outside  $A_i$ .
- For any I' ⊆ I let C<sub>I'</sub> := ({∧<sub>i</sub> | i ∈ I'}). Then all C<sub>I'</sub> are distinct. (Note: the numbers of such clones = 2<sup>2|X|</sup>!)

# Local clones

Let *X* be infinite. A clone *C* is local if each fragment  $C \cap O_X k$  is closed in the product topology (pointwise convergence) on  $X^{X^k}$  (with discrete *X*). Equivalently: If there is a set *R* of relations of finite arity such that C = POL(R).

The lattice of local clones has only  $2^{|X|}$  elements; the lattice of all clones:  $2^{2^{|X|}}$ .

Example:

On a finite set with *k* elements, the intervall  $[O_X^{(1)}, O_X]$  has k + 1 elements.

On any infinite set X, the intervall  $[\mathcal{O}_X^{(1)}, \mathcal{O}_X]$  in the lattice of all clones has at least  $2^{2^{|\mathbb{N}|}}$  elements.

On any infinite set X, the intervall  $[\mathcal{O}_X^{(1)}, \mathcal{O}_X]$  in the lattice of *local clones* has at only countably many elements.

## Bonus round: non-AC

We used  $X \times X \approx X$  to show that  $\langle \mathbb{O}_X^{(2)} \rangle = \mathbb{O}_X$  (for infinite sets *X*). But  $X \times X \approx X$  uses the axiom of choice (and in fact  $\forall X$  infinite :  $X \times X \approx X$  is equivalent to AC). Was that necessary? Yes, probably.

Proof sketch. Really: a hint. An idea of a hint. No satisfaction guaranteed.

Let  $(M, R_3)$  be the "random 3-uniform hypergraph". That is,  $R_3$  is a totally symmetric totally irreflexive relation which is "as random as possible". For example: For all (reasonable) finite sets  $\{a_1, b_1, \ldots, a_k, b_k, c_1, d_1, \ldots, c_n, d_n\} \subseteq M$  there is some  $e \in M$  with  $R(a_i, b_i, e)$  for all *i*, and  $\neg R(c_j, d_j, e)$  for all *j*. (Technically: the Fraïssé limit of all finite 3-uniform hypergraphs.)

Clones (3&4)

# non-AC, continued

## Continuation of the proof.

Let  $(M, R_3)$  be the "random 3-uniform hypergraph". (M countable,  $R_3 \subseteq M^3$  is "random" or "generic".) Let  $f_1, \ldots, f_m$  be first order definable binary operations, say definable from  $m_1, \ldots, m_k$  in the structure (M, R). Then the set  $X \times X$  can be partitioned into finitely many sets according to the "type" a pair (x, y) can have over  $m_1, \ldots, m_k$ . On each type each operation  $f_i$  must be either constant or a projection, so the same is true for any element of  $\langle f_1, \ldots, f_k \rangle$ . But the function  $\chi_B$ is neither a projection or a constant on any type. So we have found a definable ternary function not in the clone generated by the definable binary functions.

# non-AC, conclusion

We have found a definable ternary function on (M, R), definable from R, but not in the clone generated by the definable binary functions.

Now construct a model of  $ZF_{+\neg}AC$  in which all operations on M are definable from R and finitely many parameters. In this model, all binary operations are trivial on a large set, but not all ternary operations.

# Summary

- The clone lattice on {0, 1} is well understood. (But nontrivial.)
- Cl(X) for larger finite sets X: many fragments are explicitly known (certain intervals, coatoms, ...), others only partially (atoms), or only for very small sets (say, |X| ≤ 4,5).
- ► To analyse k-ary operations, it is often helpful to consider k + 1-ary operations. (Or 2k-ary. or (k + |X|<sup>2</sup>)-ary, etc.)
- Many open questions.
- For infinite X: set theory kicks in. Local clones more interesting than all clones?

Thank you for your attention! and for your questions! ... and for your corrections!!