## Clones (1&2)

Martin Goldstern

Discrete Mathematics and Geometry, TU Wien

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## Base set X

Let *X* be a (nonempty) set.

Often finite:

```
X = {0,1}.
X = {0,*,1}.
X = { {}, {a}, {b}, {a,b} }.
X = {1,...,n}.
Etc.
```

Sometimes countably infinite:

► 
$$X = \mathbb{N} = \{0, 1, 2, \ldots\}.$$

Sometimes uncountably infinite:

$$X = \mathbb{R}$$
, etc.

# Operations on X

X = our base set.

- ▶ A unary operation is a (total) function  $f: X \to X$ .
- ▶ A binary operation is a function  $f: X^2 \to X$ .
- ternary, quaternary, . . .
- ▶ A  $\frac{k}{k}$ -ary operation is a function  $f: X^k \to X$  (for  $k \ge 1$ ).
- ► We write  $\mathcal{O}^{(k)}$  or  $\mathcal{O}_X^{(k)}$  for the set of all k-ary operations on X. (Sometimes also written  $X^{X^k}$ .)
- We let  $\mathcal{O}_X := \bigcup_{k=1}^{\infty} \mathcal{O}_X^{(k)}$ .

(For simplicity we will assume that the sets  $X^k$  are pairwise disjoint. We will ignore the 0-ary functions and replace them by constant 1-ary functions.)

## Transformation monoids

## Definition ((abstract) monoid)

A monoid or abstract monoid is a structure (M, \*, 1), where

- \* is a binary operation on M, associative
- ... together with a neutral element 1 (1 \* a = a \* 1 = a).

## Definition (transformation/concrete monoid, unary clone)

A *transformation monoid* is a subset  $T \subseteq \mathcal{O}_X^{(1)}$  (for some X) which is closed under composition and contains the identity function  $id: X \to X$ .  $((T, \circ, id))$  will be an abstract monoid.)

Conversely, a variant of Cayley's theorem shows that every abstract monoid is isomorphic to a transformation monoid.

# Binary clones

A transformation monoid or unary clone on X is a subset  $T \subseteq \mathcal{O}_X^{(1)}$  which is closed under composition and contains the identity function  $id: X \to X$ .

### Definition

A *binary clone* on X is a set  $T \subseteq \mathcal{O}_X^{(1)}$  which is closed under "composition" and contains the two projections

$$\pi_1, \pi_2: X^2 \to X.$$

## **Definition (Composition)**

Let  $f, g_1, g_2 \in \mathcal{O}_X^{(2)}$ . The composition  $f(g_1, g_2)$  is the function from  $X^2$  to X defined by

$$f(g_1g_2)(x,y) := f(g_1(x,y), g_2(x,y))$$

# *k*-ary clones

## Definition (k-ary clone)

A k-ary clone on X is a set  $T \subseteq \mathcal{O}_X^{(k)}$  which is closed under "composition" and contains the k projections

$$\pi_1,\ldots,\pi_k:X^k\to X.$$

### **Definition (Composition)**

Let  $f, g_1, \ldots, g_k \in \mathcal{O}_X^{(k)}$ . The composition  $f(g_1, \ldots, g_k)$  is the function from  $X^k$  to X defined by

$$\forall \vec{x} \in X^k : f(g_1, \ldots, g_k)(\vec{x}) := f(g_1(\vec{x}), \ldots, g_k(\vec{x}))$$

("Plugging  $g_1, \ldots, g_k$  into f")

## Clones

### Definition (Clone)

A *clone* on X is a set  $T \subseteq \mathcal{O}_X = \bigcup_{k=1}^{\infty} \mathcal{O}_X^{(k)}$  which is closed under "'composition" and contains all projections  $\pi_k^n : X^n \to X$ ,  $n = 1, 2, \ldots, 1 < k < n$ .

## **Definition (Composition)**

Let  $f \in \mathcal{O}^{(k)}$ ,  $g_1, \ldots, g_k \in \mathcal{O}_X^{(m)}$ . The composition  $f(g_1, \ldots, g_k)$  is the function from  $X^m$  to X defined by

$$\forall \vec{x} \in X^m : f(g_1, \dots, g_k)(\vec{x}) := f(g_1(\vec{x}), \dots, g_k(\vec{x}))$$

("Plugging  $g_1, \ldots, g_k$  into f") If C is a clone, then  $C^{(k)} := C \cap \mathcal{O}^{(k)}$  is a k-ary clone, the k-ary fragment of C.

# Examples of clones

- ▶ The smallest clone  $J_X$  contains only the projections.
- ▶ The largest clone  $\mathcal{O}_X$  contains all operations.
- ▶ Every subset  $S \subseteq \mathcal{O}_X$  will *generate* a clone  $\langle S \rangle$ , the smallest clone containing S. The clone  $\langle S \rangle$  can be obtained from below by closing S under composition, or from above as  $\langle S \rangle = \bigcap \{ M \mid S \subseteq M \subseteq \mathcal{O}_X, M \text{ is a clone } \}$ .
- ▶ If V is a vector space over the field K, then the set of all linear functions  $f_{\vec{a}}: V^k \to V$

$$f_{\vec{a}}(v_1,\ldots,v_k):=a_1v_1+\cdots+a_kv_k$$
 (with  $\vec{a}=(a_1,\ldots,a_k)\in K^k$ ) is a clone.

## Examples of clones, continued

For every algebra  $\mathcal{X} = (X, f, g, ...)$  (=universe X with operations f, g, ... — for example  $\mathcal{X}$  might be a group, a ring, etc) we consider

- ▶ the clone of *term operations* on X, the smallest clone containing all the basic operations f, g, ... of  $\mathfrak{X}$ ;
- the clone of polynomial operations on X, the smallest clone containing all terms as well as all constant unary functions on X.

Many properties of the algebra  $\mathcal{X}$  depend only on the clone of term functions, and not on the specific set of basic operations which generates this clone. (E.g. subalgebras, congruence relations, automorphisms, etc)

For example, a Boolean algebra will have the same clone as the corresponding Boolean ring.

## The family of all clones

For any nonempty set X let CI(X) be the set of all clones on X.

- ► The intersection of any subfamily of CI(X) is again in CI(X).
- (CI(X), ⊆) is a complete lattice. Meet = intersection, join = generated by union.
- ▶  $J_X$  is the smallest clone,  $\mathcal{O}_X$  the largest.
- ▶ If  $X = \{0\}$ , then there is a unique clone:  $J_X = \mathcal{O}_X$ .
- ▶ If  $X = \{0, 1\}$ , then CI(X) is countably infinite.
- ▶ If X is finite and has at least three elements, then CI(X) is uncountable. (In fact:  $|CI(X)| = |\mathbb{R}|$ .)
- ▶ If X is infinite, then ... (later)

# Uncountably many clones

If  $X = \{0, 1, 2\}$ , then CI(X) is uncountable.

### Proof sketch.

- ▶ We call a k-tuple  $(a_1, ..., a_k) \in \{0, 1, 2\}^k$  proper, if exactly one of the  $a_i$  is equal to 1, and all the others are 2.
- ► For every  $k \ge 3$  let  $f_k : X^k \to X$  be the function that assigns 1 to every proper k-tuple, and 0 to everything else.
- ▶ For every  $A \subseteq \{3,4,\ldots\}$  let  $C_A := \langle \{f_i \mid i \in A\} \rangle$ .
- Check that for k ∉ A we have f<sub>k</sub> ∉ C<sub>A</sub>. (Every composition of functions f<sub>i</sub>, i ≠ k will assign 0 to some proper k-tuple.)
- ▶ Hence the map  $A \mapsto C_A$  is 1-1.

## Completeness

Fix a base set X.

### Definition

A set  $S \subseteq \mathcal{O}_X$  is *complete* if  $\langle S \rangle = \mathcal{O}_X$ , i.e., if every operation on X is term function of the algebra with operations S.

### Example

Let  $X = \{0, 1\}, \mathcal{X} = (X, \vee, \wedge, \neg, 0, 1).$ 

- ▶ The set  $\{\lor, \land, \neg\}$  is complete.
- ▶ The set  $\{\land, \neg\}$  is complete.
- ► The set {|} is complete, where  $x|y := \neg(x \land y)$ . (Sheffer stroke)

# Completeness, more examples

#### **Theorem**

For every  $X: \langle \mathcal{O}_X^{(2)} \rangle = \mathcal{O}_X$ .

### Proof.

- finite: Lagrange interpolation
- ▶ infinite: use  $X \times X \approx X$ .

Caution: Most clones C are NOT generated by their binary fragment  $C \cap \mathcal{O}^{(2)}$ . (Not even finitely generated.)

#### **Theorem**

If 
$$X = \{1, ..., k\}$$
, then there is a single function  $f \in \mathcal{O}_X^{(2)}$  with  $\langle f \rangle = \mathcal{O}_X^{(2)}$ : Let  $f(x, x) = x + 1$  (modulo k),  $f(x, y) = 0$  otherwise.

# (Completeness on infinite sets)

If X is infinite, then  $\mathcal{O}_X$  is uncountable. Hence a finite/countable set of operations cannot generate all of  $\mathcal{O}_X$ .

### However:

### **Theorem**

Let  $X \neq \emptyset$ . For any finite or countable set  $T \subseteq \mathcal{O}_X$  there is a single function  $f_T$  (not necessarily in T) such that  $T \subseteq \langle f \rangle$ .

#### **Theorem**

- If X is countable, then there is a countable dense subset of ①<sub>X</sub> (in the natural topology), hence there is a single function f such that the topological closure of ⟨f⟩ is all of ①<sub>X</sub>.
- If X is uncountable, then  $O_X$  will not be separable any more.

# Completeness, continued

Let  $X=\{0,1\}$  be the 2-element Boolean algebra, with Boolean operations  $\land,\lor,\lnot,\to,|,\ldots$ 

## Example

The set  $\{\lor, \land, \to\}$  is not complete.

### Proof.

Each of the three operations preserves the set  $\{1\}$ , i.e., this set is a subalgebra of the algebra  $(\{0,1\}, \land, \lor, \rightarrow)$ .

Hence every function in  $\langle \{ \land, \lor, \rightarrow \}$  will also preserve this set, but  $\neg$  does not. So  $\neg \notin \langle \{ \land, \lor, \rightarrow \} \rangle$ .

# Polymorphisms, example

### Example

The set  $\{\lor, \land, 0, 1\}$  is not complete.

### Proof.

All four functions are monotone in both arguments.

### **Definition**

Let  $\rho \subseteq X \times X$  be a relation (Example:  $\leq$  on  $\{0, 1\}$ .)

A function  $f: X^k \to X$  preserves  $\rho$  iff:

for all 
$$\begin{pmatrix} x_1 \\ y_1 \end{pmatrix}, \dots, \begin{pmatrix} x_k \\ y_k \end{pmatrix} \in \rho$$
, we have  $\begin{pmatrix} f(x_1, \dots, x_k) \\ f(y_1, \dots, y_k) \end{pmatrix} \in \rho$ .

### Lemma

If all  $f \in S \subseteq \mathcal{O}_X$  preserve  $\rho$ , then all  $f \in \langle S \rangle$  preserve  $\rho$ .

# Polymorphisms, definition

#### Definition

Let  $\rho \subseteq X^m$  be an m-ary relation, and let  $f: X^k \to X$  be a k-ary function. We say that "f preserves  $\rho$ " ( $f \triangleright \rho$ ,  $f \in \text{Pol}(\rho)$ ) if:

- for all  $(a_{i,j}: i \leq m, j \leq k) \in X^{m \times k}$ :
  - whenever  $a_{*,1} \in \rho, \ldots, a_{*,k} \in \rho$

• then also 
$$\begin{pmatrix} f(a_{1,*}) \\ \vdots \\ f(a_{m,*}) \end{pmatrix} \in \rho$$
.

(We let 
$$a_{*,j}:=\left(egin{array}{c} a_{1,j} \\ \vdots \\ a_{m,i} \end{array}\right)$$
, similarly  $a_{i,*}=(a_{i,1},\ldots,a_{i,k})$ .)

# Polymorphisms, examples

- Let  $\rho$  be a nontrivial unary relation, i.e.  $\emptyset \subsetneq \rho \subsetneq X$ . Then  $Pol(\rho)$  is the set of all operations f such that  $\rho$  is a subalgebra of (X, f).
- Let  $\rho \subseteq X \times X$  be an equivalence relation. Then  $Pol(\rho)$  is the set of all operations f such that  $\rho$  is a congruence relation of the algebra (X, f).
- ▶ Let  $\rho \subseteq X \times X$  be a (reflexive) partial order. Then Pol( $\rho$ ) is the set of all pointwise monotone operations.
- Let ρ ⊆ X × X be the graph of a function r:
   ρ = {(x, r(x)) : x ∈ X}.
   Then Pol(ρ) is the set of all functions f such that r is an endomorphism of (X, f), i.e., f commutes with r.

Fix a finite base set X.

#### Definition

For any relation  $\rho \subseteq X^m$  let  $\operatorname{Pol}(\rho)$  be the set of all operations preserving  $\rho \colon \operatorname{Pol}(\rho) := \{ f \in \mathfrak{O}_X \mid f \rhd \rho \}$ For a set R of relations, let  $\operatorname{POL}(R) := \bigcap_{\rho \in R} \operatorname{Pol}(\rho)$ .

#### Lemma

If  $S \subseteq Pol(\rho)$ , then also  $\langle S \rangle \subseteq Pol(\rho)$ . In particular,  $Pol(\rho)$  and also POL(R) are always clones.

### **Theorem**

For every clone  $C \subseteq \mathcal{O}_X$  there exists:

- ▶ A set  $S \subseteq \mathcal{O}_X$  such that  $C = \langle S \rangle$ . (Trivial)
- ▶ A set R of relations such that C = POL(R).

(Helpful to show incompleteness.)

## Galois connection

#### **Theorem**

For every clone  $C \subseteq \mathcal{O}_X$  there exists a set R of relations such that  $C = \mathsf{POL}(R) = \{f \mid \forall \rho \in R : f \rhd \rho\}.$ 

#### Proof sketch.

The largest set R satisfying  $\forall \rho \in R : C \subseteq Pol(\rho)$  is the set

$$\mathsf{INV}(C) := \{ \rho \mid \forall f \in C : f \rhd \rho \}$$

For finite sets X, we can check that C = POL(INV(C)).

even:  $\langle S \rangle = \text{POL}(\text{INV}(S))$  for all  $S \subseteq \mathcal{O}_X$ . We will see a construction of a "better" set R with C = POL(R) later.

# Pol: completeness criterion

Fix a finite base set *X*.

#### **Theorem**

For every clone  $C \subseteq \mathcal{O}_X$  there exists a set R of relations such that C = POL(R).

## Corollary

If  $S \subseteq \mathcal{O}_X$  is not complete (i.e.,  $\langle S \rangle \neq \mathcal{O}_X$ ), then there is a nontrivial relation  $\rho$  such that  $S \subseteq \text{Pol}(\rho)$ , hence  $\langle S \rangle \subseteq \text{Pol}(\rho)$ .

(But there are so many candidates for  $\rho$ ! Want to search a small set.  $\rightarrow$  precomplete clones)

# Precomplete clones

#### **Definition**

A clone  $C \subseteq \mathcal{O}_X$  is "precomplete" (or "maximal") if  $C \neq \mathcal{O}_X$ , but there is no clone D satisfying  $C \subsetneq D \subsetneq \mathcal{O}_X$ .

#### **Theorem**

For any clone  $C \subsetneq \mathfrak{O}_X$  there is a precomplete clone C' with  $C \subset C'$ .

(Remark: Not true for infinite sets!)

#### Proof.

(Use Zorn's lemma??) Let  $\mathcal{O}_X = \langle f \rangle$ . Among all clones D with  $C \subseteq D$ ,  $f \notin D$ , find a maximal element. (Better proof: later)

# Examples of precomplete clones

## Example

Let  $\emptyset \subsetneq \rho \subsetneq X$ . Then Pol( $\rho$ ) is precomplete.

### Proof.

Assuming  $g \notin Pol(\rho)$ , we let  $C := \langle Pol(\rho) \cup \{g\} \rangle$ ; we show  $C = \mathcal{O}_X$ .

First show that there is  $b \notin \rho$  such that the constant operation  $c_b$  with value b is in C.

For any function  $f: X^k \to X$  let  $\hat{f}: X^{k+1} \to X$  be defined by  $\hat{f}(\vec{x}, b) = f(\vec{x})$ , and  $\hat{f}(\vec{x}, y) \in \rho$  arbitrary for  $y \neq b$ . Then  $\hat{f} \in C$ , and  $f(\vec{x}) = \hat{f}(\vec{x}, c_b(x_1))$ , so  $f \in C$ .

### Example

Let  $\rho$  be a bounded partial order. Then  $Pol(\rho)$  is precomplete.

# Rosenberg's list

### **Theorem**

Let  $X = \{1, \dots, k\}$ . Then there is an explicit finite list of relations  $\rho_1, \dots, \rho_m$  (including, for example, all nontrivial unary relations, all bounded partial orders) such that every precomplete clone on X is one of  $Pol(\rho_1), \dots, Pol(\rho_m)$ .

Completeness criterion If  $\langle S \rangle \neq \mathfrak{O}_X$  iff there is some i with  $\forall f \in S : f \rhd \rho_i$ .

# k-ary fragments

Let *D* be a *k*-ary clone. The smallest clone *C* with  $C \cap \mathcal{O}_X^{(k)} = D$  is  $\langle D \rangle$ .

 $D \subseteq X^{X^k}$  can be viewed as a relation on X.

The largest clone C with  $C \cap \mathcal{O}_X^{(k)} = D$  is

$$\mathsf{Pol}(D) = \bigcup_{n} \{ f \in \mathcal{O}_X^{(n)} \mid \forall d_1, \dots, d_n \in D : f(d_1, \dots, d_n) \in D \}$$

For any clone E, the clones  $Pol(E \cap \mathcal{O}_X^{(k)})$  approximate E from above, agreeing with E on larger and larger sets:

$$\mathsf{Pol}(E \cap \mathcal{O}_X^{(k)}) \cap \mathcal{O}_X^{(k)} = E \cap \mathcal{O}_X^{(k)}.$$

### **Theorem**

For all clones  $E: E = \bigcap_k Pol(E \cap \mathcal{O}_X^{(k)})$ .

# CI(X) is dually atomic

#### **Theorem**

Let X be finite,  $C \neq 0_X$  a clone. Then there is a precomplete clone  $D \supseteq C$ .

### Proof.

Let  $C' \supseteq C$  be such that  $C' \cap \mathcal{O}_X^{(2)}$  is maximal. (finite!) Let  $D := \operatorname{Pol}(C')$ .