Theories of relational lattices. AKA: Embeddability into relational lattices is undecidable¹

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Playing around with real world databases

Relational lattices

Quasiequational theories of relational lattices

The lattice of a frame

p-morphisms from lattice embeddings

More on equational theory

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Databases, tables, sqls ...

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Databases, tables, sqls . . .

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Operations on tables: the natural join

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Operations on tables: the inner union

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Lattices from databases

Proposition. [Spight & Tropashko, 2006] The set of tables, whose columns are indexed by a subset of A and values are from a set D, is a lattice, with natural join as meet and inner union as join.

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The relational lattices R(D, A)

A a set of attributes, D a set of values.

An element of R(D, A): • a pair (X, T) with $X \subseteq A$ and $T \subseteq D^X$.

We have

$$(X_1,\,T_1)\leq (X_2,\,T_2)$$
 iff $X_2\subseteq X_1$ and $T_1|_{X_2}\subseteq T_2$.

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Meet and join

$$(X_1, T_1) \land (X_2, T_2) := (X_1 \cup X_2, T)$$

where $T = \{ f \mid f_{\uparrow_{X_i}} \in T_i, i = 1, 2 \}$
 $= i_{X_1 \cup X_2}(T_1) \cap i_{X_1 \cup X_2}(T_2),$
 $(X_1, T_1) \lor (X_2, T_2) := (X_1 \cap X_2, T)$

$$(X_1, T_1) \lor (X_2, T_2) := (X_1 \cap X_2, T)$$

where $T = \{ f \mid \exists i \in \{1, 2\}, \exists g \in T_i \text{ s.t. } g_{\restriction X_1 \cap X_2} = f \}$
 $= T_1 \restriction_{X_1 \cap X_2} \cup T_2 \restriction_{X_1 \cap X_2} .$

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Representation via closure operators

The Hamming/Priess_Crampe-Ribenboim ultrametric distance on D^A :

$$\delta(f,g) := \{ x \in A \mid f(x) \neq g(x) \}.$$

NB: this distance takes values in the join-semilattice $(P(A), \emptyset, \cup)$.

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A subset X of $A + D^A$ is *closed* if $\delta(f, g) \cup \{g\} \subseteq X$ implies $f \in X$.

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Proposition. [Litak, Mikulás and Hidders 2015] R(D, A) is isomorphic to the lattice of closed subsets of $A + D^A$.

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Undecidable quasiequational theories

Theorem. [Litak, Mikulás and Hidders, 2015] The set of quasiequations in the signature (\land,\lor,H) that are valid on relational lattices is undecidable.

We refine here this to:

Theorem. The set of quasiequations in the signature (\land, \lor) that are valid on relational lattices is undecidable.

We actually prove a stronger result:

Theorem. It is undecidable whether a finite subdirectly irreducible lattice embeds into some R(D, A).

Related undecidable problems

Theorem. [Maddux 1980] The equational theory of 3-dimensional diagonal free cylindric algebras is undecidable.

Theorem. [Hirsch and Hodkinson 2001] It is not decidable whether a finite simple relation algebras embeds into a concrete one (a powerset of a binary product).

Theorem. [Hirsch, Hodkinson and Kurucz 2002] It is not decidable whether a finite mutimodal Kripke frame has a surjective pmorphism from a universal product frame.

Frames, universal product frames

- A (multimodal Kripke) A-frame is a pair $(X, \{R_a \mid a \in A\})$ with
 - X a set, and
 - $R_a \subseteq X \times X$, for each $a \in A$.

Frames, universal product frames

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A universal S5-product frame is an A-frame $(X, \{ R_a \mid a \in A \})$ with

- ► $X = \prod_{a \in A} Y_a$,
- $\vec{x}R_a\vec{y}$ iff $\vec{x}_b = \vec{y}_b$, for each $b \neq a$.

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A *p*-morphism from $(X, \{R_a \mid a \in A\})$ to $(X', \{R'_a \mid a \in A\})$ is a function $f : X \to X'$ such that

- ► xR_ay implies $f(x)R'_af(y)$, for each $a \in A$,
- f(x)R'_ay' implies xR_ay for some y ∈ X such that f(y) = y', for each a ∈ A.

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The lattice of a frame

Let $\mathcal{F} = (X, \{ R_a \mid a \in A \})$ be a finite A-frame.

If $\alpha \subseteq A$, then we say that $Y \subseteq X$ is α -closed if

$$x_0 R_{a_1} x_1 R_{a_2} x_2 \dots R_{a_n} x_n \in Y \text{ and } \{a_1, \dots, a_n\} \subseteq \alpha$$

implies $x_0 \in Y$.

We say that $Z \subseteq A + X$ is closed if $Z \cap X$ is $Z \cap A$ -closed.

Definition. The lattice $L(\mathcal{F})$ is the lattice of closed subsets of A + X.

The lattice of a frame

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Theorem. A full rooted **S4** mutimodal frame \mathcal{F} has a surjective *p*-morphism from a universal product frame iff $L(\mathcal{F})$ embeds into a relational lattice.

The easy part: embeddings from *p*-morphisms

L extends to a contravariant functor.

Moreover if $X = \prod_{a \in A} D$ (= D^A) and A is finite then $L(\mathcal{F}) = R(D, A)$.

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The easy part: embeddings from *p*-morphisms

L extends to a contravariant functor.

Moreover if
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 (= D^A) and A is finite then $L(\mathcal{F}) = R(D, A)$.

Corollary. If a finite multimodal frame \mathcal{F} has a surjective *p*-morphism from a universal product frame, then L(\mathcal{F}) embeds into some R(D, A).

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Lattice embeddings into the R(D, B)s

We study lattice embeddings of the form

$$i: L(\mathcal{F}) \longrightarrow R(D, B)$$

where \mathcal{F} is an A-frame.

Lattice embeddings into the R(D, B)s

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$$i: L(\mathcal{F}) \longrightarrow R(D, B)$$

where \mathcal{F} is an A-frame.

We can suppose that:

- 1. A = B is both the set of join-prime elements of L(\mathcal{F}) and the set of join-prime elements of R(D, B) (= R(D, A));
- 2. *i* preserves \bot , \top , so $\mu \dashv i$ (use L(\mathcal{F}) subdirectly irreducible).



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For $f \in D^A$, we have $\mu(f) \in X_F$ if and only if $f \in X_{\nu}$, where

 $X_{\nu} := \{ f \in D^{\mathcal{A}} \mid \nu(f) = \emptyset \}, \quad \nu(f) := \{ j \in \mathcal{A} \mid j \leq \mu(f) \}.$



For $f \in D^A$, we have $\mu(f) \in X_F$ if and only if $f \in X_{\nu}$, where

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ight\}.$$

Moreover ν is a module on the space (D^A, δ) , where $\lambda \in \mathbb{R}^{n}$ is a module of the space (D^A, δ) .

Some theory of generalized ultrametric spaces

An ultrametric space over P(A) is a pair (X, δ) such that

•
$$\delta(x, y) = \emptyset$$
 iff $x = y$,

- $\delta(x,z) \subseteq \delta(x,y) \cup \delta(y,z)$,
- $\delta(x, y) = \delta(y, x)$.

A space (X, δ) is *pairwise-complete* if

δ(x, z) ⊆ α ∪ β implies δ(x, y) ⊆ α and δ(y, z) ⊆ β, for some y ∈ X,

A space (X, δ) is *spherically-complete* if every chain of balls has non empty intersection.

Universal product frames as GUMSs

Theorem. [Ackerman 2013] For a GUMS (X, δ) over P(A), TFAE:

- (X, δ) is an injective object in the category of GUMS over P(A),
- (X, δ) is pairwise-complete and spherically-complete,

Universal product frames as GUMSs

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Theorem. [LS] For a GUMS (X, δ) over P(A), TFAE:

- (X, δ) is pairwise-complete and spherically-complete,
- (X, δ) are spaces of sections (universal product frames, Hamming graphs, dependent product types, ...)

Modules

An ultrametric space (X, δ) is a category enriched over $(P(A), \emptyset, \cup)$.

A module on (X, δ) is an enriched functor $v : (X, \delta) \to (P(A), \triangle)$. That is, a function $v : X \to P(A)$ such that:

 $v(x) \subseteq \delta(x,y) \cup v(y)$.

Lemma

If (X, δ) is spherically-complete and pairwise-complete and $v : (X, \delta) \rightarrow P(A)$ is a module, then its kernel

$$X_{\nu} = \{ x \in X \mid \nu(x) = \emptyset \}$$

induces a spherically-complete and pairwise-complete subspace of (X, δ) .

Completing the proof of the converse

The subspace induced by

$$X_{\nu} = \{ f \in D^{\mathcal{A}} \mid \mu(f) \in X_{\mathcal{F}} \} = \{ f \in D^{\mathcal{A}} \mid \nu(f) = \emptyset \},\$$

is the kernel of a module, therefore it is pairwise-complete (and spherically-complete), that is, a universal product frame.

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Then

$$u_{\restriction_{X_{\nu}}}:X_{\nu}\longrightarrow X_{\mathcal{F}}$$

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yields the desired surjective map.

Completing the proof of the converse

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is the kernel of a module, therefore it is pairwise-complete (and spherically-complete), that is, a universal product frame.

Then

$$\mu_{\restriction_{X_{\nu}}}:X_{\nu}\longrightarrow X_{\mathcal{F}}$$

yields the desired surjective map.

This map is a *p*-morphism since (roughly) this property corresponds to μ preserving joins.

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Theorem. If A or D is finite, then R(D, A) belongs to the variety generated by the finite R(D', A').

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▶ If A is finite, then R(D, A) is an algebraic lattice.

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Theorem. If A or D is finite, then R(D, A) belongs to the variety generated by the finite R(D', A').

- If A is finite, then R(D, A) is an algebraic lattice.
- If A is infinite, then R(D, A) is not an algebraic lattice.

Functorial properties

Let $f : A \rightarrow B$ be a (set theoretic function). Then

$$\mathsf{R}(D, f) : \mathsf{R}(D, A) \to \mathsf{R}(D, B)$$

defined by

$$\mathsf{R}(D,f)(\alpha,X) := (\forall_f(\alpha), f^{*-1}(X))$$

makes R(D, -) into a functor from **Set** to ASL.

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Proposition. If D is finite, then the canonical map

 $\mathsf{R}(D,A) o \lim_{Q ext{ a finite partition of } \mathsf{Q}} \mathsf{R}(D,A/Q)$

is injective and preserves finite joins.

From meet-semilattices to lattices

- The projective limit $\lim_{Q} R(D, A/Q)$ is an algebraic lattice.
- Compact elements are of the form j_Q(β, Y), for some (β, Y) ∈ R(D, A/Q), for some finite partition Q of A. Here j_Q is left adjoint to R(D, π_Q) : R(D, A) → R(D, A/Q) with π_Q : A → A/Q.

Proposition. If $\pi : A \to B$ is surjective, then the left adjoint to $R(D, \pi)$ is a right adjoint (that is, it preserves meets). Theorem. The projective limit $\lim_Q R(D, A/Q)$ is (up to isomorphism) the ideal completion of inductive $\operatorname{colim}_Q R(D, A/Q)$, where the latter lives in the category of lattices.

Thanks ! Questions ?

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