An opinionated look at

Description Logics

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Outline

- >> 1. Motivation
 - 2. Fundamental notions of DLs + syntax
 - 3. Formal Properties
 - 4. An Application of DLs
 - 5. Importing Knowledge from DL KBs

Motivation

Conceptual models are needed in

- *artificial intelligence* (meaning of natural language sentences, representing knowledge in general)
- database design (Entity Relationship diagrams)
- software engineering (requirements, UML)
- in the age of the Internet:
 - information integration
 - finding and composing web services

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Motivation (cont'd)

"How are Sean Lennon and Mick Jagger connectd?" http://www.pumpthemusic.com/oracle/index_post.php

Links from Sean Lennon to Mick Jagger				
Sean Lennon is t	the child of	John Lennon		
John Lennon con	mposed	Imagine		
Imagine wa	as composed by	Paul McCartney		
Paul McCartney col	llaborated minorly on	Band Aid		
	as a minor collaboration tween	David Bowie		
David Bowie col	llaborated on	David Bowie & Mick Jagger		
David Bowie & Mick wa Jagger	as a collaboration between	Mick Jagger		

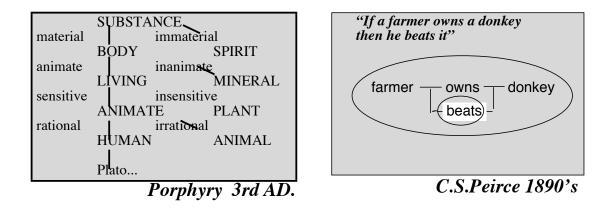
Graphical representation

Sean Lennon	is the child of	John Lenno	n
John Lennon	composed	Imagine	
Imagine	was composed by	Paul McCa	rtney
Paul McCartney	collaborated minorly on	Band Aid	
Band Aid	was a minor collaboration	on between	David Bowie
David Bowie	collaborated on		David Bowie & Mick Jagger
David Bowie & Mick Ja	ngger was a colla	boration betw	veen Mick Jagger

Sean Lennon is_child_of, John Lennon	composed Imagine
	composed_by
BandAid < collab_minorly_on Paul McCartn	ev
was_minor_collab_of	
	b_on Bowie & Jagger
Mick Jagger *	llab_between

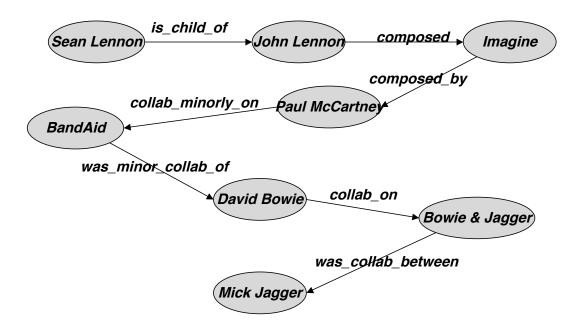
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Not entirely new idea



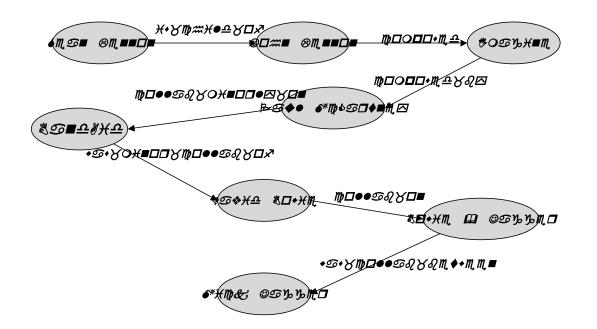
SEMANTIC NETWORKS in Artificial Intelligence/Cognitive Science *Quillian 1966* ...

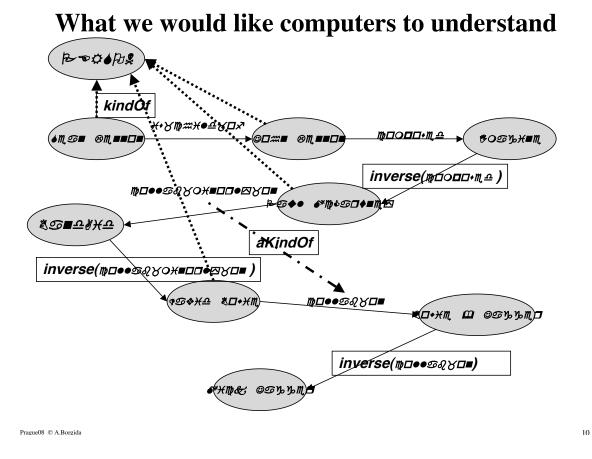
What we say to computers



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What computers "hear"





Outline

- 1. Motivation
- >> 2. Fundamental notions of DLs + syntax
 - 3. Specification of reasoning + some formal properties
 - 4. Applications of DLs

Description Logics

• A precise notation for representing "noun phrases" [Brachman 70's: KL-ONE]

Fundamental ontology: conceptual model is populated by

- <u>individuals</u>
- related by binary relationships (called *roles & features*)
- grouped into classes (concepts)
- So we need the ability to describe concepts, relationships, individuals.

First Order Logic would be fine, but it is impossible to reason with it decidably.

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Description Logics (cont'd)

- **Fundamental observation 1:** In addition to <u>primitive</u> concepts, such as **PERSON**, **CHAIR**, ... there are <u>defined</u> concepts
 - some have names:
 - "person with age between 13 and 17" = TEENAGER
 - "person who eats only non-meat foods" = VEGETARIAN
 - others are describable only by relative clauses or compound nouns:
 - "person who has at least 3 children"
 - "towns located in MA or NH or VT,..." (NEW_ENGLAND_TOWNS)

Description Logics (cont'd)

Fundamental observation 2: Both primitive and defined concepts can have additional assertions made about them, representing *necessary* conditions.

A standard way to make such assertions is to use

is-a / is-subconcept-of / is-subsumed-by / is-a-kind-of

PERSON is-a ANIMATE PERSON is-a ("age having an integer value") TEENAGER is-a LIKES_MTV

Description Logics (cont'd)

We need a language for defining concepts. (Based on empirical experience on what has been useful in many applications):

- atomic/primitive concepts: PERSON, COURSE, BOOK
- boolean combinations of these:
 - AFRICAN and HERBIVORE
 - PERSON or CORPORATION
- concepts defined by enumeration of individuals: {Masc,Fem}
- concepts from "concrete domains" (numbers, strings, ...)
- primitive binary relationships graduateOf, locatedIn, likes, hasPart
- sets of objects satisfying restrictions on their role fillers
 - objects all of whose locatedIn values are in NEW_ENGLAND_TOWNS
 - objects some of whose graduateOf values are in UNIVERSITY
 - objects with <u>at least</u> 3 hasPart fillers
 - objects whose firstName <u>same as</u> father's firstName
 - objects whose name values include "Jr."

Description Logics - syntax (1)

Just like {and,or,not} are logical *formula* constructors, DLs offer *concept constructors*. Will use term/prefix notation here:

•and(AFRICAN, HERBIVORE)
<pre>.not(ANIMATE)</pre>
•or(PERSON, CORPORATION)
<pre>•and(PERSON, not(TEEN))</pre>
<pre>•enum(Masc , Fem)</pre>
·INTEGER
V all (locIn,NEW_ENGLAND_TWN)
<pre>.some(graduateOf,UNIVERSITY)</pre>
<pre>•at-least(3,hasPart)</pre>
<pre>•same-as([name],[father name])</pre>

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Description Logics -syntax

Can describe concepts of arbitrary complexity by nesting. (Unlike OO, etc. no need to name concepts)

"Courses taken by 60 to 90 students, who are all juniors or seniors, and taught by a CS professor"

```
• and(
```

Description Logics -syntax variants

"Persons who eat only non-meat stuff"

```
• (and PERSON (all eats (not MEAT)))
· PERSON || Veats.-MEAT

• <concept> <and>

           <primitive name="PERSON"/>
           <all>
                 <primrole name="eats"/>
                  <not> <primitive name="MEAT"/> </not>
           </all> </and> </concept>
· <owl:intersectionOf rdf:parseType="Collection">
    <owl:Class rdf:about="#PERSON" />
    <owl:Restriction>
       <owl:onProperty rdf:resource="#eats" />
       <owl:allValuesFrom>
           <owl:complementOf rdf:resource="#MEAT" />
       </owl:allValuesFrom/>
     </owl:Restriction>
  </owl:intersectionOf>
```

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Description logics: roles/properties

DL Fundamental observation 3: Relationships are like concepts. Hence they can also be structured and defined, using *role constructors*.

- *loves* is-a-kind-of *likes* loves *is-a* likes
- childOf is the inverse of parentOf
 inverse(parentOf)
- descendantOf is the transitive closure of childOf
 trans(childOf)
- nephewOf is the composition of sonOf and siblingOf
 compose(sonOf,siblingOf)

Concept/Description *Languages:* summary

- Descriptions are composite, variable-free *terms*, which can be recursively built up from primitive symbols, using *constructors*
- There are constructors for both concepts and roles (binary relationships)
- There is a collection of constructors that have been empirically found useful over the years

So what can one do with descriptions?

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Standard "judgements" about Descriptions

1. Does C subsume D? D :< C $D \subseteq C$

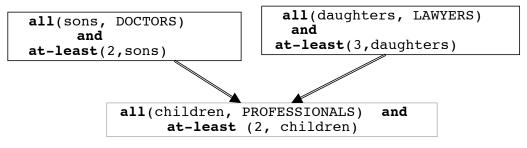
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- and(PERSON, MALE) :< PERSON
- at-least(3, hasChildren):<at-least(1, hasChildren)
- and(all(p,C) , all(p,D)) :< all(p , and(C,D))
- fills(loves, Eve) :< at-least(1, likes)</pre>
```

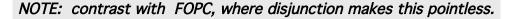
2. Is concept C incoherent?

and(PERSON
 at-least(3, hasDegree)
 all(hasDegree , enum("BA", "BS"))

Non-standard judgements

3. What is the <u>least common subsumer</u> of concepts C and D: lcs(C,D) in the (infinite) lattice of all description terms!!! [B] (Useful in machine learning.)





4. <u>Matching/Unification</u> [B] (Useful in printing relevant aspects)

e.g., matching all(hasParts, ?Y) against ARCH yields ?Y ↔ BLOCK; But macthing is against "semantic complection" of ARCH !

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How does one use DLs?

- (A specific DL consists of a particular set of concept & role constructors)
- Then create a theory \mathcal{T} of subsumption and definition assertions (or other kinds of assertions

e.g., A disjoint_from B \equiv and(A,B) :< \perp \mathcal{T} is usually called a T-box ("ontology", "knowledge base")

- As part of creating \mathcal{T} , concepts in it are
 - automatically pre-classified into a subsumption hierarchy
 - tested for "reasonableness" (satisfiable)
- \mathcal{T} can then be queried to see if it entails other judgements

DLs and individuals/nominals

Two new judgements ٠

> Mimi : HAPPY sisterOf(Anna,Mimi)

ind membership roles relating inds

- Create a theory A of assertions about individuals, usually called an A-box ("database")
- As part of creating *A*, individuals in it are (often)
 - automatically pre-classified under the most specific named *concept* in T-box taxonomy
 - tested for "reasonableness" (satisfiable)
 - some propagations cached
- A can then be queried to see if it entails other judgements

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Sample Individual Reasoning

• <u>Assertions</u> : individuals can be asserted to satisfy descrip	otions
Calvin : PERSON	
Calvin : all(friendOf, the(age and(min(5),max	(7)))
• <u>Consistency checking</u> : given additional assertion	Open World Assumption
friendOf(Calvin, Susie)	
verify that Susie's age is not known to be under 5 or over 7	
• <u>Propagation</u> if Susie's age is not known, then infer partial infe	ormation
<pre>Susie : the(age , and(min(5),max(7)))</pre>	
• Individual Classification in either case, if we have a definition	n like
CHILD = _{def} the(age , and(min(0),max(12)))
then Susie is inferred to be a child	
Susie : CHILD	

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- 1. Motivation
- 2. Fundamental notions of DLs + syntax
- >>3. Formal properties
 - 4. Applications of DLs
 - **5. DDL**

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Expressive power

 Even the most expressive DL ever proposed = FOL + counting quantifiers + fix point <u>but only 3 variable</u> <u>symbols</u>

- so cannot represent 4-clique

happy(X) :- likes(X,Y1),likes(X,Y2),likes(Y1,Y2),...

• But open-world assumption, ALL-restrictions, definitions, put it beyond Datalog

• Subsets of DL are variants of

- modal logic K
- Propositional Dynamic Logic
- Guarded Fragment of FOL

Some well known DLs

• Classic (early 1990's, AT&T Bell Labs [B]

- low-order polynomial time reasoning
- used in industrial application at AT&T to configure switching equipment

• FaCT *SHIQ* (late 90's, Manchester)

- optimized tableaux implementation
- used for large (5000 concept) medical ontology, which is not just a tree
- although logic is EXPTIME-complete, in practice not a problem!?

• OWL-DL

- the ontology language of the semantic web
- *SHOIQ*(C)

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Some complexity results

Constructors		T-Box	S	ubsumes?	Member?
	(prim :< D)	(D :< C)	cyclic		
$\mathcal{AL}(and,all)$				O(n ²)	
AL	*			co-NP-complet	e
CLASSIC with ho	het			O(n ³)	
individuals	551			0(11*)	
ALE(and,all,sor	ne)			NP-compl.	PSPACE
ALC (and,all,not ALC(and,all,not		*		PSPACE-comp EXPTIME-com	
ALCNR(r-and,r ALCNR,SHIQ NEXPTIME	nrs)	*	*	PSPACE * NEXPT	
ALCQ, ALCN	+complex rc	oles but no	<i>t</i> r-and	EXPTIME-com	plete
AL & role same				undecidable	
.AL & tunc'n role	same-as		*	poly-time	

Outline

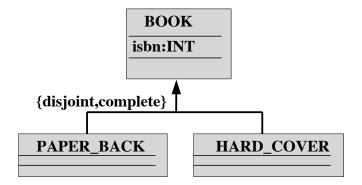
- 1. Motivation
- 2. Fundamental notions of DLs
- 3. Syntax, semantics, some formal properties

>> 4. Application of DLs

- (representing UML class diagrams-- hence reasoning about consistency)
- describing e-services/programs

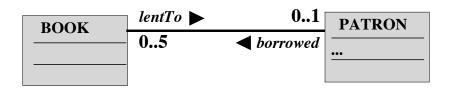
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Representing UML in SHIQ / DL-lite



BOOK :< the(isbn, INT) PAPER_BACK :< BOOK HARD_COVER :< BOOK BOOK :< or(PAPER_BACK, HARD_COVER) ;;complete and(PAPER_BACK, HARD_COVER) :< NOTHING ;;disjoint

Representing UML in SHIQ / DL-lite



BOOK :< all(lentTo,PATRON) and at-most(1,lentTo)
borrowed =def= inverse(lentTo)
PATRON :< all(borrowed,BOOK) and at-most(5,borrowed)</pre>

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An application: e-service description [B]

CORBA interface: $\frac{attrib}{attrib} CAR-MODEL model;$ $\frac{attrib}{attrib} OWNER ownedBy;$ $\frac{attrib}{attrib} MANUFACT madeBy;$ \cdots deliver(in MANUFACT src, in DEALER dest, in DATE time)signals (BadDealer); sell(...); destroy(...);

1. Create class for CAR with attributes and methods as properties:

(model **some** CAR_MODELS) (ownedBY **some** OWNER) (madeBy **some** MANUFACT) (deliver **some** DELIVER)

SE application: e-service description

CORBA interface: $\frac{\text{interface CAR}{\text{attrib CAR-MODEL model};}{\text{attrib OWNER ownedBy};}{\text{attrib MANUFACT madeBy};}{\dots}$ $\frac{\text{deliver}(\text{ in MANUFACT src,}{\text{ in DEALER dest,}}{\text{ in DATE time}}){\text{signals (BadDealer)};}{\text{sell}(\dots);}{\text{destroy}(\dots);}$

2. Reify methods, to describe parameters as attributes

DELIVER :< ACTION and (this some CAR) (src some MANUFACT) (dest some DEALER) (time some DATE)

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SE application: e-service description

DELIVER :< ACTION **and** (this **some** CAR) (src **some** MANUFACT) (dest **some** DEALER) (time **some** DATE)

3. Describe service semantics by giving pre- and post-conditions, conditions for exceptions,...

CAR :<
 (model some CAR_MODELS)
 (ownedBY some OWNER)
 (madeBy some MANUFACT)
 (deliver some DELIVER)
 //preconds include
 (madeBy same-as deliver.src)
 //postconds include
 (ownedBy same-as deliver.dest)
 //exception BadDealer signalled when
 (not (src overlaps dest.represents))</pre>

Pros and Cons of DLs

Pros

- Has been found empirically useful to describe "natural" domains we talk about (can represent and reason with ER and UML diagrams)
- "Open World Assumption" helps with reasoning in the presence of incomplete knowledge
- Syntax avoids variables, quantifiers, and supports nested complex concepts without having to name them
- Distinguishes *definitions* from *primitive concepts*, and applies uniformly to relationships and concepts
- Intermediate in expressive power between propositional and full First Order Predicate Calculus
- Well-explored complexity picture for many combinations of constructors

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Pros and Cons of DLs

Cons

- Expressive limitation: 3FOL + fixed point logic
- **Poor at describing mathematical concepts** (algebraic equations and reasoning with them)
- Cannot express even conjunctive queries (non-recursive Datalog)
- Vast majority of 'ontologies' being built are simple (simple hierarchies of terms (e.g., DMOZ, Yahoo), or at most UML). For these, OWL is overkill

References

- Description Logic Handbook, F. Baader et al, Cambridge Press, 2003
- Annual Description Logic workshops (20 so far). Electronic proceedings on web -- search for

dblp DL 2006

some [Borgida...] papers:

- "CLASSIC: A Structural Data Model for Objects", *SIGMOD Database Conf.* 1989 (*with* R. Brachman, D. McGuinness, L. Alperin Resnick)
- "Description Logics in Information Management", *IEEE Trans. Knowl. & Data Engineering* (1995)
- "On the Relative Expressiveness of Description Logics and Predicate Logics", *Artif. Intelligence Journal* (1996)
- "Adding more 'DL' to IDL: Towards More Knowledgeable Component Inter-Operability", Int. Conf. on Software Engineering (ICSE) 1999 (with Prem Devanbu)
- **"Explaining** *ALC* **subsumption"**, *ECAI*'2000, (*with* E. Franconi, I. Horrocks, D. McGuinness)
- "Distributed Description Logics", Journal of Data Semantics 1(1), 2004, (with L.Serafini)
- "On Concept Similarity" (DL'2006) (with T.Walsh, H.Hirsh)
- "Importing Knowledge from T-Boxes" (DL'2007)

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Outline

- 1. Motivation
- 2. Fundamental properties of DLs
- 3. Syntax and reasoning with DLs
- 4. Using DLs in Information Management
- 5. Importing knowledge from DL KBs

On Importing Knowledge

- It is important to **reuse** knowledge from previous KBs when building new ones.
- Study the notion *"KB1 imports identifiers S={N,...} from KBexp"*

Basic Desiderata:

- behave as if all of KBexp was included in KB1
- *but* minimize import to make understanding easier and reasoning faster
- accept possibly additional names & axioms imported, not just *S*
 - $S = \{ \text{Dog}, \text{Cat} \},$
 - Dog :< Carnivore :< Animal, Cat :< Carnivore

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On Importing Knowledge

"KB1 imports identifiers *S*={N,...} from KBexp"

Approach 1: based on the notion of "module"

- **KBexp** partitioned into modules *M1*,... which are exported *a priori*.
- Each needed module is then imported as a unit (so imported concept name N comes with everything in its module)
- I. Modules are created by hand, by the developer
- II. Automatic modularization
 - based on more or less syntactic (graph theoretic) grounds
 - based on logical properties

On Importing Knowledge

"KB1 imports identifiers *S*={N,...} from KB_{exp}"

Approach 2: Use list S of names to customize material imported

- III. Define and compute *import(KB1,S,KBexp*)
- IV. Use names in *S* to write special axioms ("bridge rules") connecting KB1 and KBexp, and treat KB1 and KBexp as independent, communicating sources

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Defining *import*(*S*,*KB2*)

Borgida [DL'07,WOMO'07] Grau et al [WWW'07] *Issues*

• Axioms imported form a subset of

- 1. theorems(KBexp)
- 2. KBexp
- 3. expanded(KBexp)
 - to deal with dependence on syntax, avoid irrelevant material
- How to define "minimal amount of knowledge to be imported"

 $\varphi \mid \operatorname{vocab}(\operatorname{KB} \cup \{\varphi\}) \cap \operatorname{vocab}(\operatorname{KB2}) \subseteq S \text{ and } \operatorname{KB} \cup \operatorname{KB2} \models \varphi$

- just for this importing KB? or for all possible ones?
- Influence of importing KB
 - limit the places where symbols from S can appear (this may limit the set of axioms that need to be brought)

Computing *import*

- Even in very simple cases (hierarchies with disjointness), cost of *minimizing* makes problem co-NP hard
- [Grau et al] have syntactic condition on KBexp ("locality") which allows import to be found effectively
- In general, problem related to "conservative extensions", and is hard

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IV. Multi-logics with "connections"

Local Semantics

- **1.** DDL (Distributed DL)
- 2. E-connections
- **3. P-DL**
- 4. [Stuckenschmidt&Klein ISCW 04]

Characteristic:

• denotational semantics does not assume the same domain of interpretation for all ontologies

Distributed Desription Logics

Borgida & Serafini [J of Data Semantics 2004] Serafini, Borgida & Tamlin [IJCAI 2005]

GIVEN: 〈T1, T2, { T2 imports T1\$A} 〉 + very restricted use of these imported names in T2! Only in axioms of the form H⊆1:A /*A onto H*/ 1:B ⊆G /*B into G*/ (and actually, ⊑ is not real subsumption: it is mediated by *domain relation* r₁₂ connecting Domain1 to Domain2
1:teamA |---> {2:Pele, 2:Julinho,...}
RESULTS:
specification of DDL entailment ⟨T1,T2,imports ⟩ |=_{ddl} 2: E F ⊑
implementation as distributed tableaux theorem prover
fixed point characterization using H :< G1 ∨ ... ∨ Gn derived from bridge axioms and T1

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E-connections

Grau, Parsia, & Sirin [ISWC 2004]

Analogy to DDL:

GIVEN: (T1, T2, {T2 imports concept T1\$A}) + *R* Somewhat restricted use of these imported names in T2!

Imported concepts can only be used in T2 to create new restrictions on the special roles in R, using a specific set of constructors. (But once defined, such concepts can be used anywhere in T2.)

RESULTS: spec and implementation for OWL-DL importers

>Can *simulate DDL* by using $R = \{ \mathbf{r}_{12}^{-} \}$ *"into":* T1\$A :< ($\forall \mathbf{r}_{12}$. G) = ($\exists \mathbf{r}_{12}^{-}$. T1\$A) :< G *"onto":* H :< ($\exists \mathbf{r}_{12}^{-}$. T1\$B)

Summary

- Exciting times in Description Logics (too exciting for my taste ;-)
- Lots of work on modularization
- Return to interest on low-expressivity DLs
 - EL – DL Lite

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