

The Faber-Manteuffel Theorem and its Consequences

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joint work with

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Optimal Krylov subspace methods

and low memory requirements?

- Consider a system of linear algebraic equations

$$\mathbf{A}x = b$$

$\mathbf{A} \in \mathbb{R}^{n \times n}$ is nonsingular, $b \in \mathbb{R}^n$.

- Given x_0 , find an *optimal*

$$x_j \in x_0 + \mathcal{K}_j(\mathbf{A}, r_0)$$

so that the error is minimized in a given vector norm.

- What are **necessary and sufficient conditions** on \mathbf{A} so that the optimal x_j can be computed using **short recurrences**? (only a constant number of vectors is needed)

Examples of optimal Krylov subspace methods

with short recurrences

CG [Hestenes, Stiefel 1952], MINRES, SYMMLQ [Paige, Saunders 1975]

- *Optimal* in the sense that they minimize some error norm:

$$\|x - x_j\|_{\mathbf{A}} \text{ in CG,}$$

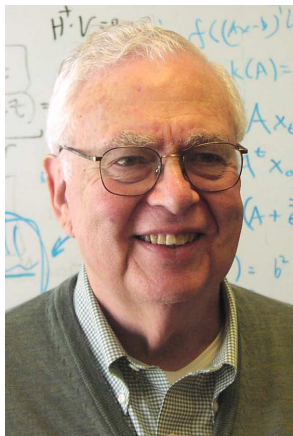
$$\|x - x_j\|_{\mathbf{A}^T \mathbf{A}} = \|r_j\| \text{ in MINRES,}$$

$$\|x - x_j\| \text{ in SYMMLQ - here } x_j \in x_0 + \mathbf{A}\mathcal{K}_j(\mathbf{A}, r_0).$$

- Generate orthogonal (or \mathbf{A} -orthogonal) Krylov subspace basis using a three-term recurrence,

$$r_{j+1} = \gamma_j \mathbf{A} r_j - \alpha_j r_j - \beta_j r_{j-1}.$$

- An important assumption: \mathbf{A} is **symmetric** (MINRES, SYMMLQ) and **positive definite** (CG).



G. H. Golub, 1932–2007

- By the end of the 1970s it was unknown if such methods existed also for general unsymmetric A .
- Gatlinburg VIII (now Householder Symposium) held in Oxford in 1981.
- “A prize of \$500 has been offered by Gene Golub for the construction of a 3-term conjugate gradient like descent method for non-symmetric real matrices or a proof that there can be no such method”.

What kind of method Golub had in mind

- We want to solve $\mathbf{A}x = b$ using **CG-like descent method**: error is minimized in some given inner product norm,

$$\|\cdot\|_{\mathbf{B}} = \langle \cdot, \cdot \rangle_{\mathbf{B}}^{1/2}.$$

- Starting from x_0 , compute

$$x_{j+1} = x_j + \alpha_j p_j, \quad j = 0, 1, \dots,$$

p_j is a direction vector, α_j is a scalar (to be determined),

$$\text{span}\{p_0, \dots, p_j\} = \mathcal{K}_{j+1}(\mathbf{A}, r_0), \quad r_0 = b - \mathbf{A}x_0.$$

- $\|x - x_{j+1}\|_{\mathbf{B}}$ is minimal iff

$$\alpha_j = \frac{\langle x - x_j, p_j \rangle_{\mathbf{B}}}{\langle p_j, p_j \rangle_{\mathbf{B}}} \quad \text{and} \quad \langle p_j, p_i \rangle_{\mathbf{B}} = 0.$$

- p_0, \dots, p_j has to be a **B-orthogonal basis** of $\mathcal{K}_{j+1}(\mathbf{A}, r_0)$.

Optimal Krylov subspace method with short recurrences

The question about

the existence of an optimal Krylov subspace method with short recurrences

can be reduced to the question:

For which \mathbf{A} is it possible to generate a **B-orthogonal basis** of the Krylov subspace using short recurrences?

(for each initial starting vector)

NECESSARY AND SUFFICIENT CONDITIONS FOR THE EXISTENCE OF A CONJUGATE GRADIENT METHOD*

VANCE FABER[†] AND THOMAS MANTEUFFEL[†]

Abstract. We characterize the class $CG(s)$ of matrices A for which the linear system $A\mathbf{x}=\mathbf{b}$ can be solved by an s -term conjugate gradient method. We show that, except for a few anomalies, the class $CG(s)$ consists of matrices A for which conjugate gradient methods are already known. These matrices are the Hermitian matrices, $A^*=A$, and the matrices of the form $A=e^{i\theta}(dI+B)$, with $B^*=-B$.

- Faber and Manteuffel gave the answer in 1984:
For a general matrix A there exists *no* short recurrence
for generating orthogonal Krylov subspace bases.
- What are the details of this statement?

Outline

- 1 The Faber-Manteuffel theorem
- 2 Ideas of a new proof
- 3 Consequences
- 4 Other types of recurrences

Formulation of the problem

B-inner product, Input and Notation

Without loss of generality, $\mathbf{B} = \mathbf{I}$. Otherwise change the basis:

$$\langle x, y \rangle_{\mathbf{B}} = \langle \mathbf{B}^{1/2}x, \mathbf{B}^{1/2}y \rangle, \quad \hat{\mathbf{A}} \equiv \mathbf{B}^{1/2}\mathbf{A}\mathbf{B}^{-1/2}, \quad \hat{v} \equiv \mathbf{B}^{1/2}v.$$

Input data:

- $\mathbf{A} \in \mathbb{C}^{n \times n}$, a nonsingular matrix.
- $v \in \mathbb{C}^n$, an initial vector.

Notation:

- $d_{\min}(\mathbf{A}) \dots$ the degree of the minimal polynomial of \mathbf{A} .
- $d = d(\mathbf{A}, v) \dots$ the grade of v with respect to \mathbf{A} , the smallest d s.t. $\mathcal{K}_d(\mathbf{A}, v)$ is invariant under multiplication with \mathbf{A} .

Formulation of the problem

Our Goal

- Generate a basis v_1, \dots, v_d of $\mathcal{K}_d(\mathbf{A}, v)$ s.t.
 1. $\text{span}\{v_1, \dots, v_j\} = \mathcal{K}_j(\mathbf{A}, v)$, for $j = 1, \dots, d$,
 2. $\langle v_i, v_j \rangle = 0$, for $i \neq j$, $i, j = 1, \dots, d$.

The Arnoldi algorithm:

Standard way for generating the orthogonal basis
(no normalization for convenience): $v_1 \equiv v$,

$$v_{j+1} = \mathbf{A}v_j - \sum_{i=1}^j h_{i,j} v_i, \quad h_{i,j} = \frac{\langle \mathbf{A}v_j, v_i \rangle}{\langle v_i, v_i \rangle},$$

$$j = 0, \dots, d-1.$$

Formulation of the problem

The Arnoldi algorithm - matrix representation

In matrix notation:

$$\begin{aligned} v_1 &= v, \\ \mathbf{A} \underbrace{[v_1, \dots, v_{d-1}]}_{\equiv \mathbf{V}_{d-1}} &= \underbrace{[v_1, \dots, v_d]}_{\equiv \mathbf{V}_d} \underbrace{\begin{bmatrix} h_{1,1} & \cdots & h_{1,d-1} \\ 1 & \ddots & \vdots \\ & \ddots & h_{d-1,d-1} \\ & & & 1 \end{bmatrix}}_{\equiv \mathbf{H}_{d,d-1}}, \end{aligned}$$

$\mathbf{V}_d^* \mathbf{V}_d$ is diagonal, $d = \dim \mathcal{K}_n(\mathbf{A}, v)$.

$$(s+2)\text{-term recurrence:} \quad v_{j+1} = \mathbf{A} v_j - \sum_{\mathbf{i}=\mathbf{j}-s}^j h_{i,j} v_i.$$

Formulation of the problem

Optimal short recurrences (Definition - Liesen, Strakoš 2008)

A admits an optimal $(s + 2)$ -term recurrence, if

- for any v , $\mathbf{H}_{d,d-1}$ is at most $(s + 2)$ -band Hessenberg, and
- for at least one v , $\mathbf{H}_{d,d-1}$ is $(s + 2)$ -band Hessenberg.

$$\mathbf{A} \mathbf{V}_{d-1} = \mathbf{V}_d$$

Sufficient and necessary conditions on **A**?

The Faber-Manteuffel theorem

Definition. If $\mathbf{A}^* = p_s(\mathbf{A})$, where p_s is a polynomial of the smallest possible degree s , \mathbf{A} is called $\text{normal}(s)$.

Theorem

[Faber, Manteuffel 1984], [Liesen, Strakoš 2008]

Given nonsingular \mathbf{A} and nonnegative s , $s + 2 < d_{\min}(\mathbf{A})$.

\mathbf{A} admits an optimal $(s + 2)$ -term recurrence

if and only if

\mathbf{A} is $\text{normal}(s)$.

- **Sufficiency** is straightforward, **necessity** *is not*. Key words from the proof of necessity in [Faber, Manteuffel 1984] include: “continuous function” (analysis), “closed set of smaller dimension” (topology), “wedge product” (multilinear algebra).

A new proof of the Faber-Manteuffel theorem

- Motivated by the paper [Liesen, Strakoš 2008] which contains a completely reworked theory of short recurrences for generating orthogonal Krylov subspace bases.

“It is unknown if a simpler proof of the necessity part can be found. In view of the fundamental nature of the Faber-Manteuffel Theorem, such proof would be a welcome addition to the existing literature. It would lead to a better understanding of the theorem by enlightening some (possibly unexpected) relationships, and it would also be more suitable for classroom teaching.”

- In [Faber, Liesen, T. 2008] we give two new proofs of the Faber-Manteuffel theorem that use more elementary tools.

Extension of $\mathbf{A} \mathbf{V}_{d-1} = \mathbf{V}_d \mathbf{H}_{d,d-1}$

Matrix representation of \mathbf{A} in \mathbf{V}_d

Since $\mathcal{K}_d(\mathbf{A}, v)$ is invariant, $\mathbf{A}v_d \in \mathcal{K}_d(\mathbf{A}, v)$ and

$$\mathbf{A}v_d = \sum_{i=1}^d h_{i,d} v_i.$$

$$\mathbf{A} \mathbf{V}_d = \mathbf{V}_d \begin{bmatrix} \overbrace{\quad s+1 \quad} & & & & & & \\ \bullet & \dots & \bullet & & & & \bullet \\ & \ddots & & \ddots & & & \vdots \\ & & \ddots & & \ddots & & \\ & & & \ddots & & \ddots & \\ & & & & \ddots & & \vdots \\ & & & & & \ddots & \bullet \\ & & & & & & \vdots \\ & & & & & & \bullet \\ & & & & & & \bullet \end{bmatrix}$$

$\underbrace{\hspace{10em}}_{d-1}$

Faber-Manteuffel Theorem – Summary

Generating an orthogonal basis of $\mathcal{K}_d(\mathbf{A}, v)$ via Arnoldi-type recurrence

Arnoldi-type recurrence
($s + 2$)-term



\mathbf{A} is normal(s)
 $\mathbf{A}^* = p(\mathbf{A})$



the only interesting case
is $s = 1$,
collinear eigenvalues

- When is \mathbf{A} normal(s)?
- \mathbf{A} is normal and
 - [Faber, Manteuffel 1984],
[Khavinson, Świątek 2003]
[Liesen, Strakoš 2008]
 - 1. $s = 1$ if and only if the eigenvalues of \mathbf{A} lie on a line in \mathbb{C} .
 - 2. For $s > 1$, \mathbf{A} has at most $3s - 2$ different eigenvalues.
- All classes of “interesting” matrices are known.

When is \mathbf{A} orthogonally reducible

to $(s + 2)$ -band Hessenberg form?

The matrix representation of the Arnoldi algorithm can be extended by one column to

$$\mathbf{A} \mathbf{V}_d = \mathbf{V}_d \mathbf{H}_d$$

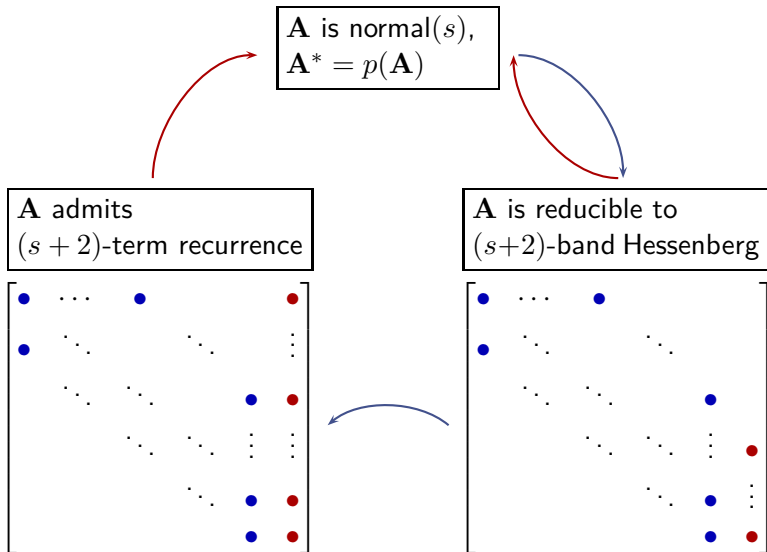
where $\mathbf{H}_d \in \mathbb{C}^{d \times d}$ is unreduced upper Hessenberg matrix.

We say that \mathbf{A} is **orthogonally reducible** to $(s + 2)$ -band Hessenberg form if \mathbf{H}_d is $(s + 2)$ -band Hessenberg matrix for each starting vector v_1 .

What are **necessary and sufficient conditions** on \mathbf{A} to be **orthogonally reducible** to $(s + 2)$ -band Hessenberg form?

When is \mathbf{A} orthogonally reducible

to $(s + 2)$ -band Hessenberg form?



When is \mathbf{A} orthogonally reducible

to $(s + 2)$ -band Hessenberg form?

Theorem

[Liesen, Strakoš 2008]

Let s be a nonnegative integer, $s + 2 < d_{\min}(\mathbf{A})$. Then the following three assertions are equivalent:

1. \mathbf{A} admits an optimal $(s + 2)$ -term recurrence.
2. \mathbf{A} is normal(s).
3. \mathbf{A} is orthogonally reducible to $(s + 2)$ -band Hessenberg form.

- $1 \iff 2$: [Faber, Manteuffel 1984].
- $2 \iff 3$: a simple proof in [Faber, Liesen, T. 2009].
- The subtle difference between 1. and 3. \rightarrow source of confusions [Voevodin, Tyrtysnikov 1981], [Liesen, Saylor 2005].

The role of the matrix \mathbf{B}

Faber-Manteuffel theorem

Let $\mathbf{B} \in \mathbb{C}^{n \times n}$ be a Hermitian positive definite (HPD), defining the **B-inner product**, $\langle x, y \rangle_{\mathbf{B}} \equiv y^* \mathbf{B} x$.

B-normal(s) matrices: there exists a polynomial p_s of the smallest possible degree s such that

$$\mathbf{A}^+ \equiv \mathbf{B}^{-1} \mathbf{A}^* \mathbf{B} = p_s(\mathbf{A}),$$

where \mathbf{A}^+ the **B-adjoint of \mathbf{A}** .

Theorem

[Faber, Manteuffel 1984], [Liesen, Strakoš 2008]

For \mathbf{A} , \mathbf{B} as above, and an integer $s \geq 0$ with $s + 2 < d_{\min}(\mathbf{A})$:

\mathbf{A} admits for the given \mathbf{B} an optimal $(s + 2)$ -term recurrence if and only if \mathbf{A} is B-normal(s).

The role of the matrix \mathbf{B} : Examples

The only interesting case: \mathbf{B} -normal(1) matrices

- If \mathbf{A} is diagonalizable and the eigenvalues are collinear, then there exists an HPD \mathbf{B} such that \mathbf{A} is \mathbf{B} -normal(1).
[Liesen, Strakoš 2008] \rightarrow complete parametrization of all \mathbf{B} 's.
- Find a **preconditioner** \mathbf{P} so that \mathbf{PA} is \mathbf{B} -normal(1) for some \mathbf{B} , e.g. [Concus, Golub 1976], [Widlund 1978], [Eisenstat 1983], [Bramble, Pasciak 1988], [Stoll, Wathen 2008].
- Saddle point matrix:

$$\mathbf{A} = \begin{bmatrix} A_1 & A_2^T \\ -A_2 & A_3 \end{bmatrix}, \quad \mathbf{B}_\gamma = \begin{bmatrix} A_1 - \gamma I_m & A_2^T \\ A_2 & \gamma I_k - A_3 \end{bmatrix}$$

where $A_1 = A_1^T > 0$, $A_3 = A_3^T \geq 0$, A_2 full rank.

This matrix satisfies $\mathbf{B}_\gamma^{-1} \mathbf{A}^T \mathbf{B}_\gamma = \mathbf{A}$.

How to choose γ such that \mathbf{B}_γ is positive definite?

[Fischer et al. 1998], [Benzi, Simoncini 2006], [Liesen, Parlett 2007].

Other types of recurrences

The existence of an optimal Krylov subspace method with short recurrences

For which \mathbf{A} is it possible to generate an **orthogonal basis** of the Krylov subspace using short recurrences?

- We can use a different kind of recurrences than Arnoldi-like.
- For (shifted) unitary matrices: **Isometric Arnoldi** process
[Gragg 1982; Jagels, Reichel 1994].
- Generalized by [Barth, Manteuffel 2000] to **(ℓ, m) -recursion**.
A sufficient condition: \mathbf{A}^* is a low degree rational func. of \mathbf{A} .
Practical use: matrices with concyclic eigenvalues [Liesen 2007].
- [Barth, Manteuffel 2000]: **Short multiple recursion** for \mathbf{A} such that $\Delta \equiv \mathbf{A}^* q_m(\mathbf{A}) - p_\ell(\mathbf{A})$ has low rank.
- [Beckermann, Reichel 2008]: GMRES-like algorithm with short recurrences for \mathbf{A} such that $\Delta \equiv \mathbf{A}^* - \mathbf{A}$ is of low rank.
Application: Path following methods.

Conclusions

- We characterized matrices for which it is possible to generate an orthogonal basis of Krylov subspaces via short recurrences.
- We presented ideas of a new proof of the Faber-Manteuffel theorem and studied its consequences.
- Practical case: If the eigenvalues of \mathbf{A} are collinear or concyclic, then there exists an HPD matrix \mathbf{B} such that \mathbf{A} admits short recurrences for generating a \mathbf{B} -orthogonal basis.
- Examples: Find a preconditioner \mathbf{P} so that short recurrences exist for \mathbf{PA} , saddle point matrices.

An interesting case to study:

- Short multiple recursion for \mathbf{A} such that $\mathbf{A}^* q_m(\mathbf{A}) - p_\ell(\mathbf{A})$ has low rank. Practical cases? Algorithmic realizations?

Related papers

- V. Faber and T. Manteuffel, [Necessary and sufficient conditions for the existence of a conjugate gradient method, SIAM J. Numer. Anal., 21 (1984), pp. 352–362.]
- T. Barth and T. Manteuffel, [Multiple recursion conjugate gradient algorithms. I. Sufficient conditions, SIAM J. Matrix Anal. Appl., 21 (2000), pp. 768–796.]
- J. Liesen and Z. Strakoš, [On optimal short recurrences for generating orthogonal Krylov subspace bases, SIAM Review, 50, 2008, pp. 485–503].
- J. Liesen, [When is the adjoint of a matrix a low degree rational function in the matrix? SIAM J. Matrix Anal. Appl., 2007 , 29 , 1171–1180].
- V. Faber, J. Liesen and P. Tichý, [The Faber-Manteuffel Theorem for Linear Operators, SIAM J. Numer. Anal., 46 (2008), pp. 1323–1337.]
- V. Faber, J. Liesen, and P. Tichý, [On orthogonal reduction to Hessenberg form with small bandwidth, Numer. Algorithms, 51 (2009), pp. 133–142.]

Thank you for your attention!