On a New Proof of the Faber-Manteuffel Theorem

Petr Tichý

joint work with

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Outline

- Introduction
- Formulation of the problem
- The Faber-Manteuffel theorem
- 4 The ideas of a new proof

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Krylov subspace methods Basis

Methods based on projection onto the Krylov subspaces

$$\mathcal{K}_j(\mathbf{A}, v) \equiv \operatorname{span}(v, \mathbf{A}v, \dots, \mathbf{A}^{j-1}v)$$
 $j = 1, 2, \dots$

$$\mathbf{A} \in \mathbb{R}^{n \times n}$$
, $v \in \mathbb{R}^n$.

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- The trivial choice $v, \mathbf{A}v, \dots, \mathbf{A}^{j-1}v$ is computationally infeasible (recall the Power Method).
- For numerical stability: Well conditioned basis.
- For computational efficiency: Short recurrence.

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- For computational efficiency: Short recurrence.
- Best of both worlds:
 Orthogonal basis computed by short recurrence.

with short recurrences

CG (1952), MINRES, SYMMLQ (1975)

• based on three-term recurrences

$$r_{j+1} = \gamma_j \mathbf{A} r_j - \alpha_j r_j - \beta_j r_{j-1} ,$$

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- generate orthogonal (or A-orthogonal) Krylov subspace basis,
- optimal in the sense that they minimize some error norm:

$$\begin{split} &\|x-x_j\|_{\mathbf{A}} \text{ in CG,} \\ &\|x-x_j\|_{\mathbf{A}^T\mathbf{A}} = \|r_j\| \text{ in MINRES,} \\ &\|x-x_j\| \text{ in SYMMLQ -here } x_j \in x_0 + \mathbf{A}\mathcal{K}_j(\mathbf{A},r_0). \end{split}$$

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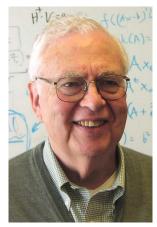
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An important assumption on A:
 A is symmetric (MINRES, SYMMLQ) & pos. definite (CG).

Gene Golub



G. H. Golub, 1932-2007

- By the end of the 1970s it was unknown if such methods existed also for general unsymmetric A.
- Gatlinburg VIII (now Householder VIII) held in Oxford from July 5 to 11, 1981.
- "A prize of \$500 has been offered by Gene Golub for the construction of a 3-term conjugate gradient like descent method for non-symmetric real matrices or a proof that there can be no such method".

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$$x_{j+1} = x_j + \alpha_j p_j$$
, $j = 0, 1, \dots$,

 p_j is a direction vector, α_j is a scalar (to be determined),

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$$\alpha_j = \frac{\langle x - x_j, p_j \rangle_{\mathbf{B}}}{\langle p_j, p_j \rangle_{\mathbf{B}}} \quad \text{and} \quad \langle p_j, p_i \rangle_{\mathbf{B}} = 0.$$

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• p_0, \ldots, p_j has to be a **B**-orthogonal basis of $\mathcal{K}_{j+1}(\mathbf{A}, r_0)$.

Faber and Manteuffel, 1984

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NECESSARY AND SUFFICIENT CONDITIONS FOR THE EXISTENCE OF A CONJUGATE GRADIENT METHOD*

VANCE FABER† AND THOMAS MANTEUFFEL†

Abstract. We characterize the class CG(s) of matrices A for which the linear system $A\mathbf{x} = \mathbf{b}$ can be solved by an s-term conjugate gradient method. We show that, except for a few anomalies, the class CG(s) consists of matrices A for which conjugate gradient methods are already known. These matrices are the Hermitian matrices, $A^* = A$, and the matrices of the form $A = e^{i\theta}(dI + B)$, with $B^* = -B$.

- Faber and Manteuffel gave the answer in 1984:
 For a general matrix A there exists no short recurrence for generating orthogonal Krylov subspace bases.
- What are the details of this statement?

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B-inner product, Input and Notation

Without loss of generality, $\mathbf{B} = \mathbf{I}$. Otherwise change the basis:

$$\langle x, y \rangle_{\mathbf{B}} = \langle \mathbf{B}^{1/2} x, \mathbf{B}^{1/2} y \rangle, \quad \hat{\mathbf{A}} \equiv \mathbf{B}^{1/2} \mathbf{A} \mathbf{B}^{-1/2}, \quad \hat{\mathbf{v}} \equiv \mathbf{B}^{1/2} v.$$

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Notation:

- ullet $d_{\min}(\mathbf{A})$... the degree of the minimal polynomial of \mathbf{A} .
- $d = d(\mathbf{A}, v)$... the grade of v with respect to \mathbf{A} , the smallest d s.t. $\mathcal{K}_d(\mathbf{A}, v)$ is invariant under mult. with \mathbf{A} .

Our Goal

• Generate a basis v_1, \ldots, v_d of $\mathcal{K}_d(\mathbf{A}, v)$ s.t.

1.
$$\operatorname{span}\{v_1, \dots, v_j\} = \mathcal{K}_j(A, v), \text{ for } j = 1, \dots, d,$$

2.
$$\langle v_i, v_j \rangle = 0$$
, for $i \neq j$, $i, j = 1, \dots, d$.

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 - 2. $\langle v_i, v_j \rangle = 0$, for $i \neq j$, $i, j = 1, \dots, d$.

Arnoldi's method:

Standard way for generating the orthogonal basis (no normalization for convenience): $v_1 \equiv v$,

$$\mathbf{v_{j+1}} = \mathbf{A}v_j - \sum_{i=1}^j h_{i,j} \mathbf{v_i}, \qquad h_{i,j} = \frac{\langle \mathbf{A}v_j, v_i \rangle}{\langle v_i, v_i \rangle},$$

$$j = 0, \dots, d - 1.$$

Arnoldi's method - matrix formulation

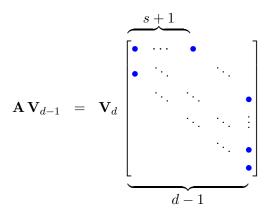
In matrix notation:

$$v_1 = v$$
,
$$\mathbf{A} \underbrace{\begin{bmatrix} v_1, \dots, v_{d-1} \end{bmatrix}}_{\equiv \mathbf{V}_{d-1}} = \underbrace{\begin{bmatrix} v_1, \dots, v_d \end{bmatrix}}_{\equiv \mathbf{V}_d} \underbrace{\begin{bmatrix} h_{1,1} & \cdots & h_{1,d-1} \\ 1 & \ddots & \vdots \\ & \ddots & h_{d-1,d-1} \end{bmatrix}}_{\equiv \mathbf{H}_{d,d-1}},$$

$$\mathbf{V}_d^* \mathbf{V}_d$$
 is diagonal, $d = \dim \mathcal{K}_n(\mathbf{A}, v)$.

Optimal short recurrences (Definition - Liesen and Strakoš, 2008)

- A admits an optimal (s+2)-term recurrence, if
 - ullet for any v, $\mathbf{H}_{d,d-1}$ is at most (s+2)-band Hessenberg, and
 - for at least one v, $\mathbf{H}_{d,d-1}$ is (s+2)-band Hessenberg.



Basic question

What are sufficient and necessary conditions for ${\bf A}$ to admit an optimal (s+2)-term recurrence?

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Example of sufficiency: If $\mathbf{A}^* = \mathbf{A}$, then s = 1 and \mathbf{A} admits an optimal 3-term recurrence.

Definition. If

$$\mathbf{A}^* = p_s(\mathbf{A}),$$

where p_s is a polynomial of the smallest possible degree s, A is called normal(s).

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Theorem. [Faber and Manteuffel, 1984], [Liesen and Strakoš, 2008]

Let ${\bf A}$ be a nonsingular matrix with minimal polynomial degree $d_{\min}({\bf A})$. Let s be a nonnegative integer, $s+2 < d_{\min}({\bf A})$:

 ${\bf A}$ admits an optimal (s+2)-term recurrence if and only if

 \mathbf{A} is normal(s).

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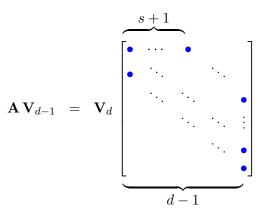
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Sufficiency is rather straightforward, necessity is not. Key
words from the proof of necessity in (Faber and Manteuffel,
1984) include: "continuous function" (analysis), "closed set of
smaller dimension" (topology), "wedge product" (multilinear
algebra).

Why is necessity so hard?

Optimal (s+2)-term recurrence:

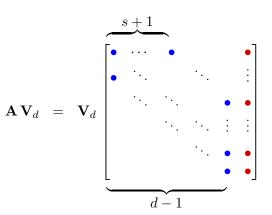


Prove something about the linear operator A, without complete knowledge of the structure of its matrix representation.

Why is necessity so hard?

Since $\mathcal{K}_d(\mathbf{A}, v)$ is invariant, $\mathbf{A}v_d \in \mathcal{K}_d(\mathbf{A}, v)$ and

$$\mathbf{A}v_d = \sum_{i=1}^d h_{id} \, \mathbf{v_i}.$$



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V. Faber, J. Liesen and P. Tichý., 2008

The Faber-Manteuffel Theorem for Linear Operators

 Motivated by the paper [J. Liesen and Z. Strakoš, 2008] which contains a completely reworked theory of short recurrences for generating orthogonal Krylov subspace bases.

"It is unknown if a simpler proof of the necessity part can be found. In view of the fundamental nature of the Faber-Manteuffel
Theorem, such proof would be a welcome addition to the existing
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- We give two new proofs of the Faber-Manteuffel theorem that use more elementary tools,
- first proof improved version of the Faber-Manteuffel proof,
- second proof completely new proof based on orthogonal transformations of upper Hessenberg matrices.

V. Faber, J. Liesen and P. Tichý, 2008

(for simplicity, we omit indices by \mathbf{V}_d and $\mathbf{H}_{d,d}$)

Let ${\bf A}$ admit an optimal (s+2)-term recurrence

$$\mathbf{A}\,\mathbf{V} = \mathbf{V}\,\mathbf{H}, \quad \mathbf{V}^*\mathbf{V} = \mathbf{I}\,.$$

Up to the last column, \mathbf{H} is (s+2)-band Hessenberg.

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$$\mathbf{A} \underbrace{\left(\mathbf{VG}\right)}_{\mathbf{W}} \; = \; \underbrace{\left(\mathbf{VG}\right)}_{\mathbf{W}} \underbrace{\left(\mathbf{G}^*\mathbf{HG}\right)}_{\widetilde{\mathbf{H}}} \; .$$

 ${f W}$ is unitary.

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$$\mathbf{A} \ \underbrace{(\mathbf{VG})}_{\mathbf{W}} \ = \ \underbrace{(\mathbf{VG})}_{\mathbf{W}} \ \underbrace{(\mathbf{G}^*\mathbf{HG})}_{\widetilde{\mathbf{H}}} \ .$$

 ${\bf W}$ is unitary. If ${\bf G}$ is chosen such that $\widetilde{\bf H}$ is again unreduced upper Hessenberg matrix, then

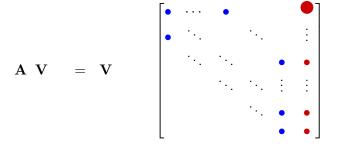
$$\mathbf{A}\mathbf{W} = \mathbf{W}\tilde{\mathbf{H}}.$$

represents the result of Arnoldi's method applied to $\bf A$ and w_1 . Up to the last column, $\widetilde{\bf H}$ has to be (s+2)-band Hessenberg.

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Proof by contradiction. Let $\mathbf A$ admit an optimal (s+2)-term recurrence and $\mathbf A$ not be normal(s).

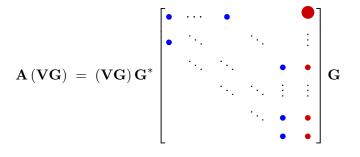
Then there exists a starting vector v such that $h_{1,d} \neq 0$.



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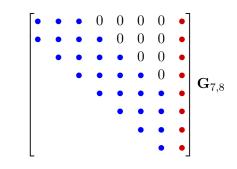
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Find unitary G (a product of Givens rotations) such that H is unreduced upper Hessenberg, but \widetilde{H} is not (s+2)-band (up to the last column) - **contradiction**.

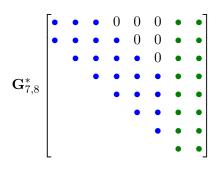
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Let v be a starting vector such that $h_{1,8} \neq 0$.



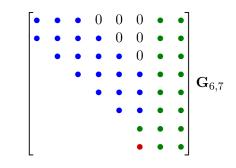
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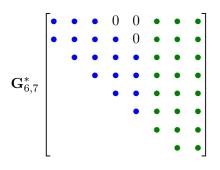
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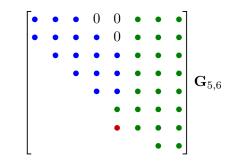
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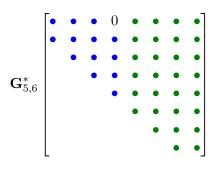
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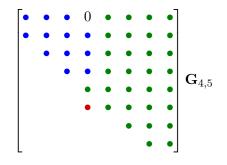
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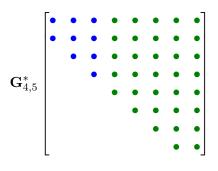
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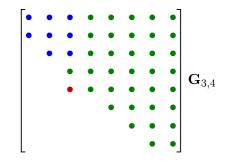
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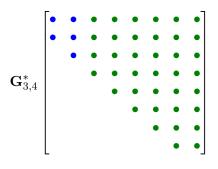
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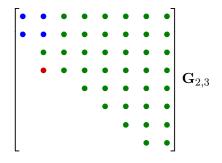
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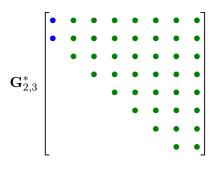
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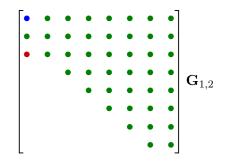
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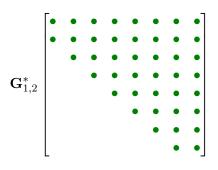
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Let v be a starting vector such that $h_{1,8} \neq 0$.

Choose Givens rotation $G_{7,8}$.

$$G \equiv G_{7,8}G_{6,7}\dots G_{1,2}, \qquad \widetilde{H} \equiv G^* H G.$$

We proved: It is possible to choose $G_{7.8}$ such that

$$h_{1,8} \neq 0 \quad \Longrightarrow \quad \tilde{h}_{1,7} \neq 0 \ \ \text{or} \ \ \tilde{h}_{2,7} \neq 0 \, .$$

Generating of orthogonal basis of $\mathcal{K}_d(\mathbf{A},v)$ via short recurrences

Arnoldi-type recurrence



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 - 1. s=1 if and only if the eigenvalues of \mathbf{A} lie on a line in \mathbb{C} .
 - 2. If the eigenvalues of \mathbf{A} are *not* on a line, then $d_{\min}(\mathbf{A}) \leq 3s 2$.

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- All classes of "interesting" matrices are known.

Related papers

- J. Liesen and Z. Strakoš, [On optimal short recurrences for generating orthogonal Krylov subspace bases, to appear in SIAM Review, 2008].
 Completely reworked theory of short recurrences for generating orthogonal Krylov subspace bases
- V. Faber, J. Liesen and P. Tichý, [The Faber-Manteuffel Theorem for Linear Operators, SIAM J. Numer. Anal., 2008, 46, 1323-1337].
 New proofs of the fundamental theorem of Faber and Manteuffel

More details can be found at

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Thank you for your attention!