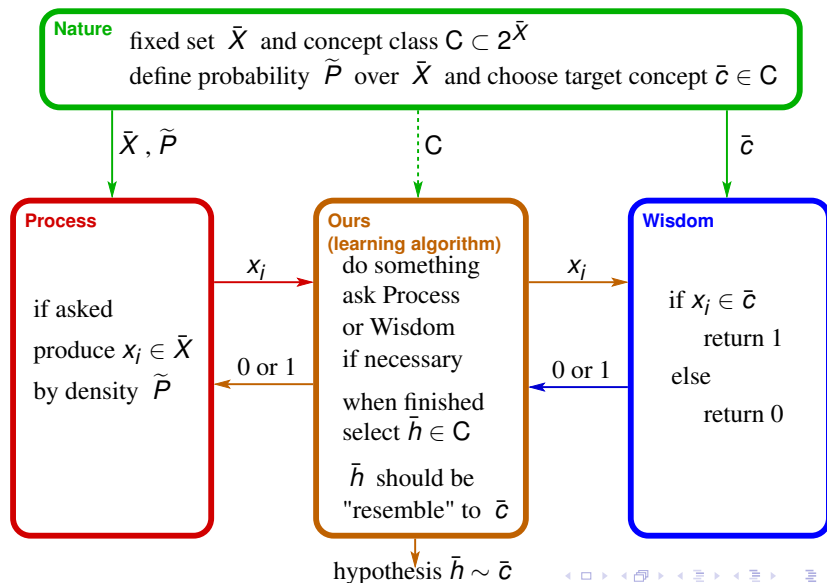


PAC learning model

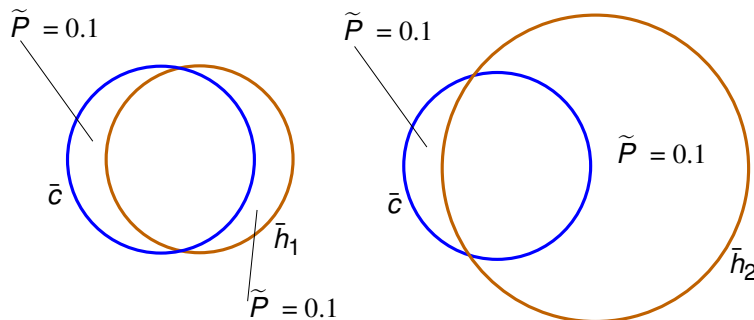
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" \bar{h} should be resemble to \bar{c} "



Definition

- 1 $e_{\bar{P}}(\bar{h}, \bar{c}) \stackrel{\text{def}}{=} \bar{P}(\bar{c} \triangle \bar{h}) \left(= (\bar{c} - \bar{h}) \cup (\bar{h} - \bar{c}) \right)$
- 2 \bar{h} is consistent if and only if $\{x_i, \dots, x_m\} \cap (\bar{c} \triangle \bar{h}) = \emptyset$

Definition (sample space)

Let $\widetilde{x} \stackrel{\text{def}}{=} \{x_1, \dots, x_m\}$, $x_i \in \bar{X}$, $i \in \{1, \dots, m\}$, $\vec{z} \in \{-1, +1\}^m$ and let $\bar{c} \subset \bar{X}$. Then the ordered tuple

$$(\widetilde{x}, \vec{z})$$

is a m -SAMPLE OF CONCEPT \bar{c} if and only if

$$(\forall i \in \{1, \dots, m\}) ((x_i \in \bar{c}) \Leftrightarrow (\vec{z}_i = 1)).$$

For concept class C define SAMPLE SPACE OF CONCEPT CLASS as

$$\bar{S}_C \stackrel{\text{def}}{=} \bigcup_{m \geq 1} \left\{ \bigcup_{\bar{c} \in C} \left\{ (\widetilde{x}, \vec{z}) \mid (\widetilde{x}, \vec{z}) \text{ is a } m\text{-sample of concept } \bar{c} \right\} \right\}.$$

Definition $((\epsilon, \delta)$ -learning algorithm)

- 1 $((\epsilon, \delta)$ -LEARNING ALGORITHM is each mapping $\widetilde{A}^* : \bar{S}_C \rightarrow C$ such that for all $\bar{c} \in C$, $\epsilon, \delta \in (0, 1)$ and \widetilde{P} on \bar{X} , the probability of the set

$$\left\{ \bar{x} \mid \left(\bar{x}, \bar{z} \right) \text{ is } m\text{-sample of } \bar{c} \text{ and } e_{\widetilde{P}} \left(\bar{c}, \widetilde{A}^* \left(\left(\bar{x}, \bar{z} \right) \right) \right) \geq \epsilon \right\}$$

is smaller than the number δ .

- 2 If such a learning algorithm exists we say that C is UNIFORMLY LEARNABLE .

Theorem

Let us assume that C is a concept class over finite set \bar{X} and $H = C$. Then, for each learning algorithm A^ requiring*

$$\frac{1}{\epsilon} \ln \left(\frac{|C|}{\delta} \right)$$

queries and producing for the given concept $\bar{c} \in C$ a consistent hypothesis it holds that

$$\text{Prob}_{\tilde{P}} \left(e_{\tilde{P}} \left(\bar{c}, \widetilde{A^*} \left(\left(\bar{x}, \bar{z} \right) \right) \right) \geq \epsilon \right) < \delta.$$

Definition

For any arbitrary $R \subset 2^{\bar{X}}$ and for any arbitrary probability density \tilde{P} defined on \bar{X} and for an arbitrary $\epsilon > 0$ let us define $R_{\tilde{P}, \epsilon} \stackrel{\text{def}}{=} \{\bar{r} \in R \mid \text{Prob}_{\tilde{P}}(\bar{r}) > \epsilon\}$. Then, we will call $\bar{T}_{\tilde{P}, \epsilon} \subset \bar{X}$ ϵ -TRANSVERSAL R just if

$$\left(\forall \bar{r} \in R_{\tilde{P}, \epsilon} \right) \left(\bar{r} \cap \bar{T}_{\tilde{P}, \epsilon} \neq \emptyset \right)$$

Example

$\bar{X} = \langle 0, 1 \rangle^n$, \tilde{P} uniform on \bar{X} , $R = \{\bar{b} \subset \bar{X} \mid \bar{b} \text{ is a ball}\}$, $\epsilon = \frac{1}{2}$.
Then $\bar{T} = \{k\epsilon \mid k = 0, \dots, \frac{1}{\epsilon}\}^n$ is a $\left(\frac{\pi^{\frac{n}{2}} (\sqrt{n}\epsilon)^n}{\Gamma(\frac{n}{2}+1)2^n} \right)$ -transversal of R .

...if hypotheses \bar{h} produced by an algorithm is consistent and has $e_{\bar{p}}(\bar{h}, \bar{c}) > \epsilon$, then $\{x_i, \dots, x_m\}$ can't be ϵ -transversal of the system $R \stackrel{\text{def}}{=} \{\bar{h} \triangle \bar{c} \mid \bar{h} \in H\}$...

Definition

For each $m \geq 1$, $\epsilon > 0$ let

$$\bar{Q}_{m,\epsilon} \stackrel{\text{def}}{=} \left\{ \bar{x} \in \bar{X}^m \mid \bar{x} \text{ do not form } \epsilon\text{-transversal of } R \right\}$$

and (assume that $\bar{x}, \bar{y} \in \bar{X}^m$)

$$\bar{J}_{\epsilon}^{2m} \stackrel{\text{def}}{=} \left\{ \bar{x}\bar{y} \in \bar{X}^{2m} \mid \left(\exists \bar{r} \in R_{\bar{p},\epsilon} \right) \left(\bar{x} \cap \bar{r} = \emptyset \text{ and } |\bar{y} \cap \bar{r}| \geq \frac{\epsilon m}{2} \right) \right\}.$$

... the probability of the set $\bar{Q}_{m,\epsilon}$ is a probability of producing consistent hypothesis with error $e_{\bar{p}}(\bar{h}, \bar{c}) > \epsilon$...

Definition

- 1 The class H is well-behaved if the sets $\bar{Q}_{m,\epsilon}$ and $\bar{J}_{\epsilon}^{2^m}$ are measurable for any probability \tilde{P} , any $m \geq 1$, $\epsilon > 0$, and any system of sets $R \stackrel{\text{def}}{=} \{\bar{h} \triangle \bar{c} \mid \bar{h} \in H\}$, where \bar{c} is an arbitrary Borelian set.
- 2 The class $H \subset 2^{\bar{X}}$ is universally separable, if there exists a countable subset T of the class H such that for all $\bar{h} \in H$ there exists a sequence $\{\bar{h}_i\}_1^{\infty}$ of sets from T such that

$$(\forall x \in \bar{X}) (\exists n \geq 1) ((\forall i \geq n) (x \in \bar{h}_i \text{ if and only if } x \in \bar{h})).$$

Theorem

If H is universally separable, then H is well-behaved.

Lemma

Let $R \neq \emptyset$ be a concept class and \tilde{P} be probability on \bar{X} for which $\bar{Q}_{m,\epsilon}$ and \bar{J}_ϵ^{2m} are measurable for all $m \geq 1, \epsilon > 0$. Then for each $\epsilon > 0$ and $m \geq \frac{2}{\epsilon}$

$$\text{Prob}_{\tilde{P}^m}(\bar{Q}_{m,\epsilon}) < 2\text{Prob}_{\tilde{P}^{2m}}(\bar{J}_\epsilon^{2m}) \leq 2\Pi_R(2m)2^{-\frac{\epsilon m}{2}}.$$

Lemma

Let $m \geq 1, \bar{c} \subset \bar{X}, H \subset 2^{\bar{X}}$ and $R \stackrel{\text{def}}{=} \{\bar{h} \triangle \bar{c} \mid \bar{h} \in H\}$. Then $\Pi_R(m) = \Pi_H(m)$.

Lemma

Let $d = \text{VC}_{\dim}(H) \in \mathbb{N}, \epsilon, \delta \in (0, 1), \Gamma \stackrel{\text{def}}{=} 2\Pi_H(2m)2^{-\frac{\epsilon m}{2}}$. Then

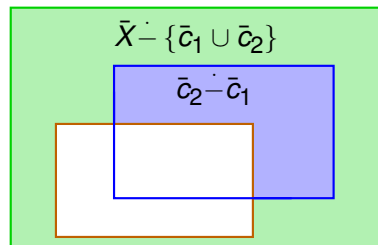
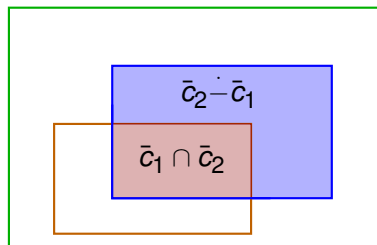
$$m \geq \max \left\{ \frac{4}{\epsilon} \log_2 \left(\frac{2}{\delta} \right), \frac{8d}{\epsilon} \log_2 \left(\frac{12.611}{\epsilon} \right) \right\} \Rightarrow \Gamma \leq \delta.$$

Definition

Concept class C is called nontrivial iff

$$(\exists \bar{c}_1, \bar{c}_2 \in C) (\bar{c}_1 \neq \bar{c}_2 \text{ and } (\bar{c}_1 \cap \bar{c}_2 \neq \emptyset \text{ or } \bar{c}_1 \cup \bar{c}_2 \neq \bar{X})) .$$

Concept class is called trivial in other cases.



Two cases of minimal content of nontrivial concept class.
(colored sets are nonempty)

Theorem (main result of PAC theory)

Let C be a nontrivial, well-behaved class. Then:

① If $\text{VC}_{\dim}(C) = d < +\infty$. Then

① for any $0 < \epsilon < \frac{1}{2}$ there is no (ϵ, δ) -learning algorithm with number of queries less than

$$\max \left(\frac{1-\epsilon}{\epsilon} \ln \left(\frac{1}{\delta} \right), d(1 - 2(\epsilon(1-\delta) + \delta)) \right). \quad (1)$$

② for arbitrary $0 < \epsilon < 1$, any learning algorithm using **at least**

$$\max \left(\frac{4}{\epsilon} \log_2 \left(\frac{2}{\delta} \right), \frac{8d}{\epsilon} \log_2 \left(\frac{12.611}{\epsilon} \right) \right) \quad (2)$$

queries and returning a consistent hypothesis is an (ϵ, δ) -learning algorithm.

② C is uniformly learnable **if and only if** $\text{VC}_{\dim}(C) < +\infty$.

Sketch of the proof:

- 1
 - 1
 - $\frac{1-\epsilon}{\epsilon} \ln \left(\frac{1}{\delta} \right)$: (c&c) Any nontrivial concept class can be reduced to one of the cases discussed above. For uniform probability we get a contradiction.
 - $d(1 - 2(\epsilon(1 - \delta) + \delta))$: (c&c) Reduce \bar{X} to d -element subset with uniform probability. Then use the "matrix" $\mathbf{Z}_{\bar{c}, \bar{h}} \stackrel{\text{def}}{=} \mathbf{e}_{\bar{P}}(\bar{c}, \bar{h})$ to show, that $m > d(1 - 2(\epsilon(1 - \delta) + \delta))$ imply that $(\exists \bar{h}^*)$ contradicts (ϵ, δ) -property ... "broadly speaking".
 - 2 See previous slides.
- 2
 - \Leftarrow (construction) Use Zermelo's well-ordering theorem to well-order \bar{H} . Let algorithm get m -sample of \bar{c} and return the first hypothesis consistent with \bar{c} . The statement follows from 1)-2).
 - \Rightarrow (by contradiction) For any $d \in \mathbb{N}$ we carry out steps 1)-1)-(second term). Choose (ϵ, δ) such that $(1 - 2(\epsilon(1 - \delta) + \delta)) > 0$. Hence m can't be upper-bounded.

Definition (discrete delta rule)

Let $(\vec{x}_1, y_1), \dots, (\vec{x}_t, y_t)$ be a given sequence of tuples in $\mathbb{R}^n \times \{-1, +1\}$, $t \geq 1$. Further, let vector's sequence $\{\vec{w}_i\}_1^\infty$ satisfy the following recursive formulas

- 1 put $\vec{w}_1 \stackrel{\text{def}}{=} \vec{0}$, $k = 1$
- 2 let $k = k + 1$ and $\bar{J} \stackrel{\text{def}}{=} \{j \in \{1, \dots, t\} \mid \widetilde{\text{sgn}}(\langle \vec{w}_k | \vec{x}_j \rangle) \neq y_j\}$
 - 1 if $\bar{J} = \emptyset$ put $\vec{w}_{k+1} = \vec{w}_k$ and STOP,
 - 2 else let $j_k \in \bar{J}$ be arbitrary. Then put

$$\vec{w}_{k+1} \stackrel{\text{def}}{=} \vec{w}_k + y_{j_k} \vec{x}_{j_k}$$

and REPEAT step 2).

Then we say that $\{\vec{w}_i\}_1^\infty$ is DELTA SEQUENCE of $(\vec{x}_1, y_1), \dots, (\vec{x}_t, y_t)$.

Theorem (delta rule convergence)

Let $\{\vec{\mathbf{w}}_i\}_1^\infty$ is a delta sequence and let there exists a vector $\hat{\vec{\mathbf{w}}}$ such that for all indexes $i \in \{1, \dots, t\}$ holds $\widetilde{\text{sgn}}(\langle \hat{\vec{\mathbf{w}}} | \vec{\mathbf{x}}_i \rangle) = y_i$. Further let

$$\alpha \stackrel{\text{def}}{=} \max_{i \in \{1, \dots, t\}} \left\{ \|\vec{\mathbf{x}}_i\|^2 \right\} \quad \text{and} \quad \beta \stackrel{\text{def}}{=} \min_{i \in \{1, \dots, t\}} \left\{ \left| \langle \hat{\vec{\mathbf{w}}} | \vec{\mathbf{x}}_i \rangle \right| \right\} > 0 .$$

Then there exists an natural number $z > 0$ satisfying $\vec{\mathbf{w}}_{z+1} = \vec{\mathbf{w}}_z$ and z can be estimated as

$$z \leq \frac{\alpha \|\hat{\vec{\mathbf{w}}}\|^2}{\beta^2} + 1 .$$

Theorem (delta rule complexity)

There exists a linearly separable dichotomy of the $\{-1, +1\}^n$ such that any integer linear separator (\vec{w}, t) of this dichotomy satisfies estimation

$$2^{\frac{n-2}{2}} \leq \sum_{k=1}^n |\vec{w}_k| + |t|.$$

Definition (Mangasarian LP)

Let $\bar{A} \stackrel{\text{def}}{=} \{\vec{a}_1, \dots, \vec{a}_i\}$ and $\bar{B} \stackrel{\text{def}}{=} \{\vec{b}_1, \dots, \vec{b}_j\}$ be a finite subsets of the \mathbb{R}^n . Then MANGASARIAN LINEAR PROBLEM is defined as the problem find vectors $\vec{y} \in \mathbb{R}^i$, $\vec{z} \in \mathbb{R}^j$, $\vec{w} \in \mathbb{R}^n$ and $t \in \mathbb{R}$ that minimizes

$$\sum_{\alpha=1}^i \vec{y}_{\alpha} + \sum_{\beta=1}^j \vec{z}_{\beta}$$

subject to

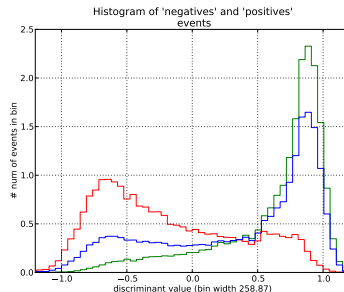
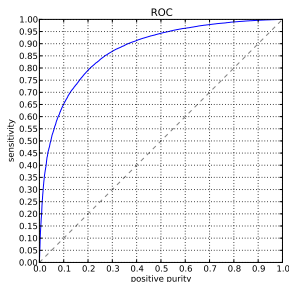
$$\begin{aligned} \vec{y}_{\alpha} + \left\langle \vec{w} \mid \vec{a}_{\alpha} \right\rangle - t &\geq 1 && \text{for } \alpha \in \{1, \dots, i\} \\ \vec{z}_{\beta} - \left\langle \vec{w} \mid \vec{b}_{\beta} \right\rangle + t &\geq 1 && \text{for } \beta \in \{1, \dots, j\} \\ \vec{y}_{\alpha} &\geq 0 && \text{for } \alpha \in \{1, \dots, i\} \\ \vec{z}_{\beta} &\geq 0 && \text{for } \beta \in \{1, \dots, j\} . \end{aligned}$$

Theorem

Let $\bar{A} \stackrel{\text{def}}{=} \{\vec{a}_1, \dots, \vec{a}_i\}$ and $\bar{B} \stackrel{\text{def}}{=} \{\vec{b}_1, \dots, \vec{b}_j\}$ be a finite subsets of the \mathbb{R}^n . Then

- 1 There exists a linear separator of the sets \bar{A} and \bar{B} if and only if the optimal value of the corresponding Mangasarian LP is zero.
- 2 If the optimal value of the corresponding Mangasarian LP is zero and $(\vec{y}^*, \vec{z}^*, \vec{w}^*, t^*)$ is optimal solution, then (\vec{w}^*, t^*) is linear separator of the sets \bar{A} and \bar{B} .

... a real case ...



The necessary condition on consistency of hypotheses produced is unsatisfied, PAC model isn't applicable. We have to use

Probably Approximately Optimal (PAO) model

\mathcal{PAC} 